

Formal Requirement Elicitation with FRET

Anastasia Mavridou KBR at NASA Ames Research Center

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Invited talk at DTU

NASA's center in Silicon Valley

Aerial image of NASA Ames Research Center Credits: NASA

robust software engineering technical area

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Popular repositories			
ikos Public	fret Public		
Static analyzer for C/C++ based on the theory of Abstract Interpretation.	A framework for the elicitation, specification, formalization and understanding of requirements.		
● C++ ☆ 1.9k 왕 166	⊖ JavaScript 🛱 212 😵 38		
CoCoSim	mesa		
Automated Analysis Framework for Simulink/Stateflow models.	Actor-based Runtime Verification Tool		
● MATLAB 🏠 29 😵 8	🕒 Scala 🛛 🏠 7		
homebrew-core Public	AdaStress.jl Public		
Homebrew formulae from NASA - Software Verification and Validation	Reinforcement learning framework to find and analyze the likeliest failures of a system under test.		
● Ruby ☆ 4	● Julia 🛱 3		

github.com/NASA-SW-VnV

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how developers write requirements

- Exceeding sensor limits shall latch an autopilot pullup when the pilot is not in control (not standby) and the system is supported without failures (not apfail).
- The autopilot shall change states from TRANSITION to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from TRANSITION to NOMINAL when the system is supported and sensor data is good.
- The autopilot shall change states from NOMINAL to MANEUVER when the sensor data is not good.
- The autopilot shall change states from NOMINAL to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from MANEUVER to STANDBY when the pilot is in control (standby) and sensor data is good.

how developers write requirements

10 Lockheed Martin Cyber-Physical System Challenge, component FSM:

Every time these conditions hold or only when they become true?

- Exceeding sensor limits shall latch an autopilot pullup when the pilot is not in control (not standby) and the system is supported without failures (not apfail).
- The autopilot shall change states from TRANSITION to STANDBY when the pilot is in control (standby).
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how developers write requirements

10 Lockheed Martin Cyber-Physical System Challenge, component FSM:

• Exceeding sensor limits shall latch an autopilot pullup when the system is supported without failures (not apfail).

• • •

Instantly, or within a time limit?

- The autopilot shall change states from TRANSITION to STANDBY when the pilot is in control (standby).
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- The autopilot shall change states from NOMINAL to MANEUVER when the sensor data is not good.
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- The autopilot shall change states from MANEUVER to STANDBY when the pilot is in control (standby) and sensor data is good.

what analysis tools understand

```
var autopilot: bool = (not standby) and supported and (not
   apfail);
var pre_autopilot: bool = false -> pre autopilot;
var pre_limits: bool = = false -> pre limits;
guarantee "FSM-001v2" S((((((autopilot and pre_autopilot and
   pre_limits) and (pre ( not (autopilot and pre_autopilot and
   pre_limits)))) or ((autopilot and pre_autopilot and
  pre_limits) and FTP)) => (pullup)) and FTP), ((((autopilot
   and pre_autopilot and pre_limits) and (pre ( not (autopilot
   and pre_autopilot and pre_limits))) or ((autopilot and
   pre_autopilot and pre_limits) and FTP)) => (pullup)));
```

FRET bridges the gap

- Captures requirements in structured natural language with unambiguous semantics.
- Explains formal semantics in various forms.
- Formalizes requirements in a compositional (hence extensible) manner.
- Checks realizability of requirements compositionally.
- Connects with analysis tools:
 - Exports formalizations in SMV language.
 - Exports Lustre code.
 - Exports specifications for runtime monitoring.



welcome to FRET



AOS AOS-R2U2-2 after new_waypoint_targeted, the vehicle shall, within 3 seconds, satisfy new_heading_achieved

github.com/NASA-SW-VnV/fret

Team: Andreas Katis, Anastasia Mavridou, Tom Pressburger, Johann Schumann, Khanh Trinh

Alumni: David Bushnell, Dimitra Giannakopoulou, Nija Shi

Interns: Milan Bhandari, Tanja DeJong, Kelly Ho, George Karamanolis, David Kooi, Jessica Phelan, Julian Rhein, Daniel Riley, Gricel Vazquez

And many other collaborators..

capturing, explaining, and formalizing requirements

let's speak FRETish

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S	Comments When in cruising mode, the whenever altitude_hold is Requirement Description A requirement follows the sentence of with "*". For information on a field for SCOPE CONDITIONS	etructure displayed below, whe rmat, click on its correspondin COMPONENT* SHALL*	ide ss indicated				
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In cruising mode, the autopilot shall always satisfy if altitude_hold then maintain_altitude



compositional generation of LTL formulas

In cruising mode, the autopilot shall always satisfy if altitude_hold then maintain_altitude

scope in: [LEFT, RIGHT) \rightarrow [FiM, LiM)

FiM = MODE and (FTP or previous (not MODE))

LiM = not MODE and previous MODE



compositional generation of LTL formulas

In cruising mode, the autopilot shall always satisfy if altitude_hold then maintain_altitude

scope in: [LEFT, RIGHT) \rightarrow [FiM, LiM)

FiM = MODE and (FTP or previous (not MODE))

LiM = not MODE and previous MODE

timing always: BASEFORM \rightarrow RES

historically (RIGHT implies previous (BASEFORM since inclusive required LEFT))

```
scope: in, condition: null, timing: always, response: satisfaction
historically (LiM implies previous (RES since inclusive required FiM))
```

```
optimize historically (MODE implies RES)
```

translate to SMV (H (MODE \rightarrow RES))

instantiate (H (cruising \rightarrow (altitude_hold \rightarrow maintain_altitude)))

related papers

Automated Formalization of Structured Natural Language Requirements

Dimitra Giannakopoulou^{a,*}, Thomas Pressburger^a, Anastasia Mavridou^b, Johann Schumann^b

^aNASA Ames Research Center, CA, USA ^bKBR, NASA Ames Research Center, CA, USA

Abstract

The use of structured natural languages to capture requirements provides a reasonable trade-off between ambiguous natural language and unintuitive formal notations. There are two major challenges in making structured natural language amenable to formal analysis: 1) formalizing requirements as formulas that can be processed by analysis tools and 2) ensuring that the formulas conform to the semantics of the structured natural language. FRETISH is a structured natural language that incorporates features from existing research and from NASA applications. Even though FRETISH is quite expressive, its

Dimitra Giannakopoulou, Thomas Pressburger, Anastasia Mavridou, Johann Schumann (2021). <u>Automated formalization of structured</u> <u>natural language requirements</u>, Information and Software Technology (IST) Journal, 137, 106590, Special Section on REFSQ'20, 2021.

Dimitra Giannakopoulou, Thomas Pressburger, Anastasia Mavridou, Johann Schumann. <u>Generation of Formal Requirements from</u> <u>Structured Natural Language</u>, REFSQ 2020.

checking realizability of requirements

- The autopilot shall change states from TRANSITION to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from TRANSITION to NOMINAL when the system is supported and sensor data is good.

- The autopilot shall change states from **TRANSITION** to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from **TRANSITION** to NOMINAL when the system is supported and sensor data is good.
- Input state: **TRANSITION**

- The autopilot shall change states from **TRANSITION** to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from **TRANSITION** to NOMINAL when the system is supported and sensor data is good.
- Input state: TRANSITION
- Condition 1: standby
- Condition 2: supported & good_sensor_data

- The autopilot shall change states from **TRANSITION** to STANDBY when the pilot is in control (standby).
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- The autopilot shall change states from TRANSITION to STANDBY when the pilot is in control (standby).
- The autopilot shall change states from TRANSITION to NOMINAL when the system is supported and sensor data is good.
- Input state: TRANSITION
- Condition 1: standby
- Condition 2: supported & good_sensor_data
- Output state 1: STANDBY
- Output state 2: NOMINAL

why realizability?

- Defining requirements is a challenging, error prone task
- Realizability checking >> consistency checking
- We want to ensure requirement consistency for all inputs
- And we want to do it efficiently

An AG contract is realizable if there exists a system implementation that satisfies the contract guarantees for all assumption-complying stimuli provided by the environment.

We proposed a novel approach for compositional realizability checking.

compositional realizability

Partial AG contracts:



checking realizability within FRET

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0	FSM001	FSM shall always satisfy (limits & !standby & !apfail & supported) => pullup				
	FSM002	FSM shall always satisfy (standby & state = ap_transition_state) => STATE = ap_standby_state				
	FSM003	FSM shall always satisfy (state = ap_transition_state & good & supported) => STATE = ap_nominal_state				
	FSM004	FSM shall always satisfy (! good & state = ap_nominal_state) => STATE = ap_maneuver_state				
	FSM005	FSM shall always satisfy (state=ap_nominal_state & standby) => STATE = ap_standby_state				
	FSM006	FSM shall always satisfy (state = ap_maneuver_state & standby & good) => STATE = ap_standby_state				
	FSM007	FSM shall always satisfy (state = ap_maneuver_state & supported & good) => STATE = ap_transition_state				
	FSM008	FSM shall always satisfy (state = ap_standby_state & !standby) => STATE = ap_transition_state				
	FSM009	FSM shall always satisfy (state = ap_standby_state & apfail)=> STATE = ap_maneuver_state				
	FSM010	FSM shall always satisfy (senstate = sen_nominal_state & limits) => SENSTATE = sen_fault_state				
		Rows per page: 10 - 1-10 of 13 < >				

related papers

From Partial to Global Assume-Guarantee Contracts: Compositional Realizability Analysis in FRET

Anastasia Mavridou¹, Andreas Katis¹, Dimitra Giannakopoulou², David Kooi³, Thomas Pressburger², and Michael W. Whalen⁴

¹ KBR, NASA Ames Research Center, CA, USA ² NASA Ames Research Center, CA, USA {anastasia.mavridou, andreas.katis, dimitra.giannakopoulou, tom.pressburger}@nasa.gov ³ University of California, Santa Cruz, CA, USA dkooi@ucsc.edu ⁴ University of Minnesota, MN, USA whalen@cs.umn.edu

Abstract. Realizability checking refers to the formal procedure that aims to determine whether an implementation exists, always complying to a set of requirements, regardless of the stimuli provided by the system's environment. Such a check is essential to ensure that the specification does not allow behavior that can force the system to violate safety constraints. In this paper, we present an approach that decomposes realizability checking into smaller, more tractable problems. More

Capture, Analyze, Diagnose: Realizability Checking of Requirements in FRET

Andreas Katis¹^[0000-0001-7013-1100], Anastasia Mavridou¹, Dimitra Giannakopoulou^{2*}, Thomas Pressburger², and Johann Schumann¹

¹ Employed by KBR; NASA Ames Research Center, CA, USA ² NASA Ames Research Center, CA, USA

Abstract. Requirements formalization has become increasingly popular in industrial settings as an effort to disambiguate designs and optimize development time and costs for critical system components. Formal requirements elicitation also enables the employment of analysis tools to prove important properties, such as consistency and realizability. In this paper, we present the realizability analysis framework that we developed as part of the Formal Requirements Elicitation Tool (FRET). Our

Andreas Katis, Anastasia Mavridou, Dimitra Giannakopoulou, Thomas Pressburger, Johann Schumann. <u>Capture, Analyze, Diagnose:</u> <u>Realizability Checking of Requirements in FRET</u>, CAV 2022.

Anastasia Mavridou, Andreas Katis, Dimitra Giannakopoulou, David Kooi, Thomas Pressburger, Michael W. Whalen. <u>From Partial to Global</u> Assume-Guarantee Contracts: Compositional Realizability Analysis in FRET, FM 2021.

connection with analysis tools

generation of Simulink monitors

FRETish:

when in roll_hold_mode autopilot shall immediately satisfy if roll_angle > 3 then roll_hold_reference = 3

Lustre specification:

-- AP-003c-v3 requirement in CoCoSpec guarantee H((roll_hold and (FTP or (pre (not roll_hold)))) => abs(roll_angle) > 30 => roll_hold_reference = 30 * sign(roll_angle))

Simulink monitor

model checking Simulink models



model checking PLC code

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□ □ □ PLCverif project □ SCL file □ Verification case 8 ✓ □ OnOff > ▷ src-gen	Verification case						
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CPC_FB_ONOFF.scl		Checking transition between Manual Mode and Auto Mode	When (instance.MMoSt & instance.AuAuMoR) the CPC_FB_OnOff shall eventually satisfy instance.AuMoSt &	MoSt & instance.AuAuMoR) the CPC_FB_OnOff shall eventually satisfy instance.AuMoSt & PLC_ENI	END Fretish requirement:		
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model checking PLC code

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FRET_Requirement.vc3	Source files Here the scope of the verific Source files: Language frontend: SI Entry block: CI Selection and configuration Backend: Ni Algorithm: IC Advanced settings	Rationale and Comments ^ Rationale If the ONOFF object is in "Manual mode" and the condition "Auto Auto Mode Request" is TRUE, the ONOFF object shall eventually be in "Auto Mode" at the end of the PLC cycle Comments Checking transition between Manual Mode and Auto Mode	 Mode Input Output Internal Undefined CPC_DB_VERSION.Baseline_version CPC_GLOBAL_VARS.First_Cycle CPC_GLOBAL_VARS.UNICOS_Counter1 CPC_GLOBAL_VARS.UNICOS_LiveCounter CPC_GLOBAL_VARS.UNICOS_TimeSmooth instance.Al instance.Al_old instance.AlB 	> (X (F :AuMoSt and
E Outline 🗁 🗇 There is no active editor that provides an outline.		Requirement Description A requirement follows the sentence structure displayed below, where fields are optional unless indicated with "*". For information on a field format, click on its corresponding bubble. Scope CONDITIONS COMPONENT* SHALL* TIMING RESPONSES* ()	 > instance.AlBW > instance.Alinc > instance.AlSt > instance.AlUnAck > instance.AuUnAck_old > instance.AuAlAck > instance.AuAlAckR_old > instance.AuAuMoR > instance.AuAuMoR_old > instance.AuAuMoR_old > instance.AuAuMoR_old > instance.AuAuMoR_old > instance.AuAuMoR_old 	
		When (instance.MMoSt & instance.AuAuMoR) the CPC_FB_OnOff shall eventually satisfy instance.AuMoSt & PLC_END SEMANTICS	 instance.AulhFoMoSt instance.AulhMMo instance.AulhMMoSt instance.AuMoSt instance.AuMoSt_aux instance.AuMoSt_old instance.AuMPW/ 	

runtime monitoring



examples of case studies/projects that use FRET



galois

High Assurance Rigorous Digital Engineering for Nuclear Safety (HARDENS)

Joe Kiniry, Galois (kiniry@galois.com)

May 2022

Theme: Driving FM to Practice

Keywords: digital engineering, model-based engineering, software engineering, hardware engineering, safety engineering, requirements engineering, formal verification, rigorous runtime verification, Cryptol, SAW, ACSL, SysML, FRET, RISC-V

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related papers

Zsófia Ádám, Ignacio D. Lopez-Miguel, Anastasia Mavridou, Thomas Pressburger, Marcin Bęś, Enrique Blanco Viñuela, Andreas Katis, Jean-Charles Tournier, Khanh V. Trinh, Borja Fernandez Adiego. <u>From Natural Language Requirements to the Verification of Programmable Logic</u> <u>Controllers: Integrating FRET into PLCverif</u>, NFM 2023.

Joseph Kiniry, Alexanders Bakst, Simon Hansen, Michal Podhradsky, and Andrew Bivin. <u>The HARDENS Final Report</u>, Galois Inc Technical Report.

Thomas Pressburger, Andreas Katis, Aaron Dutle, Anastasia Mavridou. <u>Authoring, Analyzing, and Monitoring Requirements for a Lift-Plus-Cruise</u> <u>Aircraft</u>, REFSQ 2023.

Ivan Perez, Anastasia Mavridou, Tom Pressburger, Alwyn Goodloe, Dimitra Giannakopoulou. <u>Automated Translation of Natural Language</u> <u>Requirements to Runtime Monitors</u>, TACAS 2022.

Hamza Bourbouh, Marie Farrell, Anastasia Mavridou, Irfan Sljivo, Guillaume Brat, Louise A. Dennis, Michael Fisher. <u>Integrating Formal</u> <u>Verification and Assurance: An Inspection Rover Case Study</u>, NFM 2021.

Anastasia Mavridou, Hamza Bourbouh, Dimitra Giannakopoulou, Tom Pressburger, Pierre-Loic Garoche, Johann Schumann. <u>The Ten Lockheed</u> <u>Martin Cyber-Physical Challenges: Formalized, Analyzed, and Explained</u>, RE 2020, Industry track.

Anastasia Mavridou, Hamza Bourbouh, Pierre Loic Garoche, Dimitra Giannakopoulou, Thomas Pressburger, Johann Schumann. Bridging the Gap Between Requirements and Simulink Model Analysis, REFSQ 2020, Poster Paper.

Full list: https://github.com/NASA-SW-VnV/fret/blob/master/PUBLICATIONS.md



FRET: <u>https://github.com/NASA-SW-VnV/fret</u>

CoCoSim: <u>https://github.com/NASA-SW-VnV/CoCoSim</u>

Ogma: <u>https://github.com/nasa/ogma</u>

PLCverif: <u>https://gitlab.com/plcverif-oss</u>

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Zsófia Ádám, Alexanders Bakst, Swee Balachandran, Milan Bhandari, Marcin Bęś, Enrique Blanco Viñuela, Geoffrey Biggs, David Bushnell, Maxime Artaud, Hamza Bourbouh, Guillaume Brat, Esther Conrad, Louise A. Dennis, Tanja DeJong, Michael Dille, Aaron Dutle, Marie Farrell, Borja Fernandez Adiego, Michael Fisher, Pierre-Loic Garoche, Dimitra Giannakopoulou, Alwyn Goodloe, Simon Hansen, Kelly Ho, Michael Jeronimo, George Karamanolis, Andreas Katis, Joseph Kiniry, David Kooi, Ignacio D. Lopez-Miguel, Carlos Mao de Ferro, Patrick J. Martin, Francisco Martins, Amalaye Oyake, Ivan Perez, Jessica Phelan, Tom Pressburger, Julian Rhein, Daniel Riley, Johann Schumann, Nija Shi, Irfan Sljivo, Laura Titolo, Jean-Charles Tournier, Khanh V. Trinh, Gricel Vazquez, Tim Wang, Michael W. Whalen, Alexander Will.

