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# STAR ADAPTATION FOR TWO ALGORITHMS USED ON SERIAL COMPUTERS

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to consider when adapting	an algorithm for ST	'AR are di	iscussed. Progr	ram listings of the
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# STAR ADAPTATION FOR TWO ALGORITHMS USED ON SERIAL COMPUTERS

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#### SUMMARY

Two representative algorithms used on a serial computer and presently executed on the Control Data Corporation 6000 computer were adapted to execute efficiently on the Control Data STAR-100 computer. Gaussian elimination for the solution of simultaneous linear equations and the Gauss-Legendre quadrature formula for the approximation of an integral are the two algorithms discussed. This paper describes how the programs were adapted for STAR and why these adaptations were necessary to obtain an efficient STAR program. Some points to consider when adapting an algorithm for STAR are discussed. Program listings of the 6000 version coded in 6000 FORTRAN, the adapted STAR version coded in 6000 FORTRAN, and the STAR version coded in STAR FORTRAN are presented in the appendices.

#### INTRODUCTION

Many algorithms which are presently used on the Control Data Corporation 6000 computer and executed in a serial mode are suitable for the Control Data STAR-100 computer. However, if these algorithms were converted line-by-line to the STAR coding, it is not likely that they would make an efficient STAR program and might actually produce code which would run inefficiently on a vector computer such as STAR. The 6000 code will need to be adapted for STAR to produce an efficient STAR code. This paper discusses two algorithms of this nature. A comparison of the 6000 coding and STAR coding of the identical algorithm is made for two different algorithms: one for the solution of simultaneous linear equations using Gaussian elimination with partial pivoting and the other for numerical evaluation of an integral using the Gauss-Legendre quadrature formula. The paper discusses how the 6000 program was adapted for STAR, the reasons for the adaptations, and some factors which should be considered when adapting a program. FORTRAN codings of the algorithms are presented in the appendices. The STAR codings use the FORTRAN language defined in reference 1.

#### AIDS FOR ADAPTING AN ALGORITHM

When beginning to adapt an algorithm for STAR, it is not enough to look at just segments in the 6000 coding, but it is necessary to look at the entire algorithm to get the total picture. Some questions to pose are: What is the final result?, What is needed at various steps in the algorithm?, and What is computed independent of other steps and what is repeated?

It will be helpful to review a few definitions and terms which are important to remember when formulating a program for the STAR computer. (For a comprehensive discussion of STAR architecture and hardware instructions, see ref. 2.)

(1) A vector is a set of elements stored in contiguous locations in memory. (Array and vector are used interchangeably in this discussion.)

(2) Vector timing is the time required by the central processing unit to process a vector. It is obtained by the equation T = S + l/c, where

Т	time in clocks (1 clock represents 40 nanoseconds)
S	startup time, different for each vector instruction or macro
l	length of vector (number of elements in the vector)
с	constant depending on the type of instruction and whether the vector elements are each stored in 32- or 64-bit words

(3) A page is a block of storage which contains data or instructions. A program is made up of one or more pages. (See ref. 3.)

(4) A program's working set is the smallest set of pages which must be in central memory for the program to operate efficiently.

(5) A page fault occurs when a program references a page which is not contained in central memory; that is, it is not in the program's present working set.

(6) Paging is the process of bringing a page into central memory or releasing a page.

When adapting the 6000 code to the STAR code, the following factors are important:

(1) Use vector instructions. This requirement may mean reordering steps in an algorithm so that elements in contiguous locations can be operated on or it may mean rearranging storage.

(2) Use long vectors. If a choice is available whether to use many short vectors or a few long vectors, use the long vectors unless the overhead to create the long vectors

is too great. This procedure reduces the effect of the startup time associated with each vector instruction.

(3) Avoid or use sparingly coding which will generate costly vector instructions and macros, that is, costly as compared with some of the faster vector instructions. Refer to the most current timings available. Examples of such instructions and macros are divide, transmit indexed list, and dot product.

(4) Avoid unnecessary paging problems by creating a reasonable working set for the program. When working with a particular array, perform all the operations possible with this array before working with another array. This factor will be more critical with long arrays, but it is a good habit to form and should help reduce page faults.

(5) Investigate the feasibility of creating more answers than are really needed. Because of the high result rate of vector instructions, it may be advantageous to use an approach which generates a larger number of results than are needed in order to avoid scalar computation. This idea, however, should be used cautiously. (See the discussion in ref. 4 on parallel algorithms for tri-diagonal equation solvers.)

# ADAPTATION OF AN ALGORITHM FOR THE SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS

A Langley Research Center 6000 library subroutine GELIM (see listing in appendix A) uses Gaussian elimination with partial pivoting to obtain the solution of the set of simultaneous linear equations, AX = B, where A is the square matrix of coefficients of order n, X is a vector of unknowns of length n, and B is a constant matrix of order  $n \times r$  where r is the number of right-hand sides. The matrix A is factored into a lower unit triangular matrix and an upper triangular matrix. (For numerical details of the algorithm, see any numerical analysis text, such as ref. 5.) The subroutine also contains an option for the evaluation of the determinant of matrix A.

The version of GELIM on the Langley Research Center 6000 library performs Gaussian elimination as generally defined by operating on one row of A at a time. At the kth step, column k is searched for the largest element; then row k and the row which contained the largest element are interchanged. The pivot element is then used to obtain zeroes in all positions in its column below the diagonal. This procedure requires multiplying, a row of A by the appropriate scalar and subtracting this product from another row of A.

Since these row modifications are the usual way the steps in Gaussian elimination are thought of being performed, it would seem normal that for STAR the matrix would be stored by rows. This method of storage means that elements in one row of a matrix would be stored in contiguous locations so a row of the matrix could be a vector and the steps would be performed by using vector instructions.

In the present version of STAR FORTRAN, two-dimensional arrays are stored columnwise and an optional storage arrangement by rows is not available. This column storage means that elements in one row of matrix A are not stored in contiguous locations and modifying the matrix a row at a time would mean that few vector instructions could be used. This fact makes it desirable to see whether Gaussian elimination can be performed by modifying matrix A by using columns so that vector instructions can be used. As will be shown below, Gaussian elimination can be performed by operating on columns of the matrix.

An option to store matrices by rows may be available in a later version of STAR FORTRAN, but when the 6000 code was modified to perform Gaussian elimination by columns, many advantages for the column storage over the row storage appeared; therefore, column storage is recommended. The following section will show how Gaussian elimination can be performed by using vector instructions when the matrix is stored by columns and will identify these advantages. The section entitled "Row Storage" shows the sequence of steps performed in the row storage which is identical to the present 6000 algorithm.

#### Column Storage

Gaussian elimination can be performed by modifying one column of the matrix at a time. This is done by a reordering of the operations from the usual row operations. Accomplishing the triangularization of matrix A by performing the work on columns makes efficient use of STAR and does the identical arithmetic normally done when performing the work on rows.

The kth step of the triangularization can be performed as shown below, where n is the number of equations and r is the number of right-hand sides. All references to the kth column refer to column entries below the diagonal.

(1) Divide the kth column of A by the  $a_{kk}$  element and store in the kth column. This is a vector divided by a scalar:

$$a_{ik} = \frac{a_{ik}}{a_{kk}}$$
 (i = k + 1, . . ., n)

(2) Multiply the kth column by the  $a_{kj}$  element and subtract this result from the jth column. This is a vector multiplied by a scalar and then a vector subtract:

$$a_{ij} = a_{ij} - a_{ik}a_{kj}$$
 (i = k + 1, . . ., n)

(3) Repeat the sequence of vector instructions in step (2) for the columns of A (j = k + 1, ..., n).

(4) Multiply the kth column of A by the  $b_{kj}$  element of B and subtract this result from the jth column of B. This is a vector multiplied by a scalar and then a vector subtract:

$$b_{ij} = b_{ij} - a_{ik}b_{kj}$$
 (i = k + 1, . . ., n)

(5) Repeat the vector instructions in step (4) for the remaining right-hand sides of B (j = 1, 2, ..., r).

(6) Repeat steps (1) to (5) until the triangularization of matrix A is complete (k = 1, 2, ..., n - 1).

Often a subroutine for the solution of simultaneous equations requires the user to append the right-hand sides to the original matrix. The column storage arrangement allows the right-hand sides to be done as separate vector instructions without having to append the right-hand sides to the original matrix. This arrangement is less cumbersome for the user since the right-hand side can be a separate array.

GELIM uses partial pivoting which means that rows will need to be interchanged at times. At the kth step, column k is searched for the largest element and the row containing the largest element and the kth row are interchanged. The column search to find the largest element can easily make use of vector instructions, but the row interchange presents a problem. Elements of rows will need to be interchanged and none of them will be stored in contiguous locations. Not only are they not contiguously stored, but for a very large matrix a column could use one or more pages. For a large matrix it is unlikely that a program will be allowed a working set large enough to contain the entire matrix. This situation would mean that the row elements to be interchanged would be on separate pages and would have to be brought into and then out of core only to reference two elements on the page. Then when the column modifications are performed, the same pages will need to be brought back into core again. This procedure would be extremely inefficient.

A form of indexing could be set up to achieve row interchange, but it would need one of the more time-consuming instructions (transmit indexed list) and the referencing across page boundaries would still be present. No vector instruction can help perform the interchange efficiently; thus, a scalar interchange will be just as efficient. The possible paging problem can be alleviated by not interchanging an entire row at one time as it is performed on the present 6000 version. At the kth step when the largest element in column k is found, interchange the two elements only in column k. When working with the jth column, both column k and column j are needed. Both of the elements which need interchanging are in the jth column which will be in the program's working set at that time. Therefore, before beginning the operations on the jth column, interchange the two elements. In this way, the page or pages containing that column will need to be brought into core only once per step of the algorithm.

The subroutine also computes the value of the determinant of matrix A which is equal to the product of the elements of the diagonal of the triangular matrix. In the 6000 version, the product is computed after the triangularization is completed. For the STAR, this computation presents a similar situation as the rows interchange; the diagonal elements may be on separate pages and are not stored contiguously. Therefore, vector instructions cannot be utilized. The product will be scalar multiplication, but after the kth step has been completed, the partial product can be formed by using the diagonal element of column k while column k is still in the program's working set. This grouping of row interchange, triangularization of the matrix, and evaluation of the determinant should create an efficient working set for the program with a minimum of paging.

The remaining task of the subroutine is to perform the back substitution. Back substitution generally uses the dot product of a row and the solution vector. This method is used on the 6000 version, but presents problems for STAR column storage since it uses the costly dot product macro and references elements of one array by rows and references elements of the other array by columns. This problem can be eliminated by reordering the steps needed to perform back substitution and all the work can be performed on columns.

The steps to find the kth unknown by back substitution by using vector instructions but not using the dot product are as follows:

(1) A scalar divide is always necessary,  $b_{kj} = \frac{b_{kj}}{a_{kk}}$ . This step obtains the kth unknown for the jth right-hand side and stores the unknown in the right-hand side vector.

(2) Multiply column k of A by the kth unknown obtained in step (1) and subtract this result from the jth column of B. This is a vector multiplied by a scalar and then a vector subtract:

$$b_{ij} = b_{ij} - a_{ik}b_{kj}$$
 (i = 1, 2, . . ., k - 1)

(3) Steps (1) and (2) are repeated for all the right-hand sides (columns of B) (j = 1, 2, ..., r).

(4) Steps (1), (2), and (3) are repeated for all the unknowns (k = n, ..., 2, 1). When k = 1, step (2) is omitted.

#### Row Storage

The way of performing the steps in Gaussian elimination as commonly seen in texts is by operating with rows. The elements in one row would be stored contiguously and the triangularization of matrix A would make efficient use of vector instructions.

To make the triangularization most efficient, the right-hand sides must be appended to the matrix A for row storage. The length of the longest vector used would be n + r - 1 where n is the number of unknowns and r is the number of right-hand sides. If a separate array is used for the right-hand side, two identical vector instructions would be needed for each operation, one of length n and one of length r. This method would be inefficient because the startup times would be multiplied by a factor of 2 and if r is 1, it would mean using vector instructions of length 1. Appending the right-hand side is a disadvantage in that it is awkward for the user.

Step k of the triangularization when the matrix is stored by rows is as follows:

(1) Perform a scalar divide,  $a_{ik}/a_{kk}$  and store in  $a_{ik}$ .

(2) Multiply row k of A by  $a_{ik}$  and subtract from the ith row. This is a vector multiplied by a scalar and then a vector subtract:

 $a_{ij} = a_{ij} - a_{ik}a_{kj}$  (j = k + 1, . . ., n + r)

(3) Repeat steps (1) and (2) for all rows (i = k + 1, ..., n).

(4) Repeat steps (1), (2), and (3) for all columns until the triangularization is complete (k = 1, 2, ..., n - 1).

The row interchanges necessary for partial pivoting are very easy when the matrix is stored by rows, but the column search will be scalar operations. In addition, there will be no way to avoid possible paging problems for a large matrix during the column search. The elements on each row of the column may be on separate pages, but the search has to be completed before any row operations can begin.

The determinant evaluation could be performed the same way as in the column storage, after a step in the Gaussian elimination is completed. The pages for a large matrix would not have to be brought in only to get one element for the evaluation of the determinant, but the information would be used while the page was still in memory.

When performing the back substitution, if the matrix is stored by rows, there is no way of using the scheme devised to eliminate the need for the dot product macro. The back substitution references a row of the matrix and a column of the right-hand side. To be able to use vector instructions, a vector would need to be created which would contain the elements in the column. This method would be expensive and inefficient, and it is likely that the vector code would be no better than the scalar code. Either would be inefficient here, because of the referencing of a row and a column.

Table I summarizes the comparison of the two storage arrangements at the various steps in the algorithm. The column storage is more advantageous for STAR than the row storage arrangement.

# TABLE I.- COMPARISON OF ROWWISE AND COLUMNWISE STORAGE OF MATRIX

Step	Rowwise	Columnwise
Column search	Scalar: no way to avoid possible paging problems	Vector
Row interchange	Vector: easy	Scalar: can avoid possible paging problems
Triangularization	Vector: usual way it is done	Vector: with steps reordered
Back substitution	Dot product macro: costly, referencing row and col- umn, possibly not vectorizable	Vectors: columns only, no dot product, more efficient
Determinant evaluation	Scalar	Scalar
Treatment of right-hand sides	Must be appended to matrix for efficient vector use	Can be separate variable and still use vectors efficiently

## Flow Chart

Subroutine GELIM was adapted for use on the STAR computer by using the same numerical method that was used on the 6000 computer. The matrix is stored by columns. By reordering computational steps, vector instructions can be used and a reasonable working set for the program established. Figure 1 shows a flow chart of the 6000 and STAR versions of the algorithm.

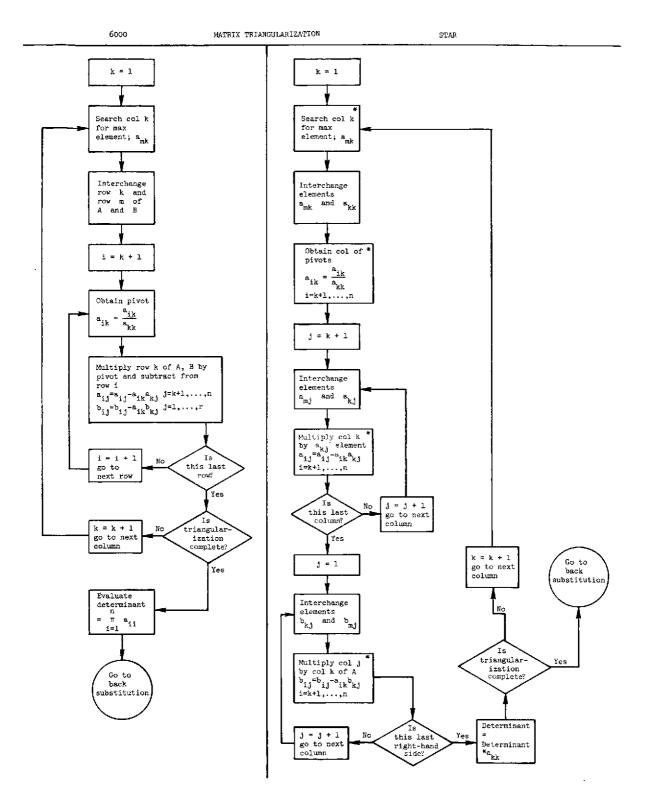


Figure 1.- Flow chart of 6000 and STAR version of segments of Gaussian elimination. An asterisk denotes use of vector instructions.

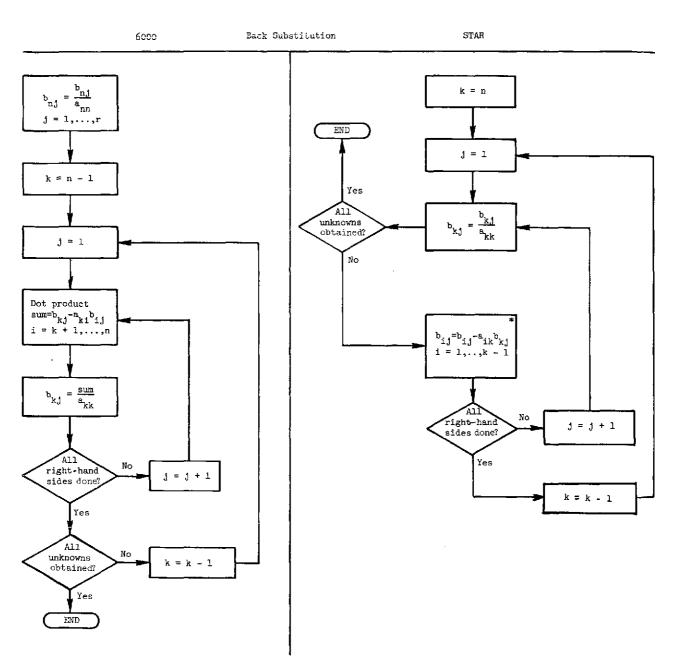


Figure 1.- Concluded.

Appendix A contains coding of the versions of the algorithm. The 6000 version and the STAR version coded in 6000 FORTRAN and in STAR FORTRAN are included. The calling sequence of the subroutine remains the same and no additional storage was necessary.

#### Additional Comment

An additional interesting fact was noted when the STAR version of subroutine GELIM was executed on the 6000 computer. The STAR version ran faster than the 6000 version; as a result, there was about a 10-percent decrease in execution time.

#### ADAPTATION OF A NUMERICAL INTEGRATION ALGORITHM

Subroutine GLEGEN (see listing in appendix B) uses the Gauss-Legendre quadrature formula to evaluate simultaneously an array of integrals

$$\int_{a}^{b} f_{k}(x) dx \qquad (k = 1, 2, ..., N)$$

where N is the number of functions. This subroutine is an example of the manner in which vectors rather than single variables can be used on STAR. As in the preceding example, the sequence of instructions is not the same as in the 6000 version. The subroutine subdivides the integration interval (a, b) into NQ panels and the Gauss-Legendre quadrature formula is applied to each panel by using a 3- or 10-point formula. The resulting integral  $I_k$  is

$$I_{k} = \Delta \sum_{j=1}^{NQ} \sum_{i=1}^{NP} r_{i}f_{k}(p_{i}) \qquad (k = 1, 2, ..., N)$$

where

 $f_k$ functionNQnumber of quadratures or panelsNPnumber of points per quadrature $p_i = l_j + \Delta x_i$  (where  $l_j$  is the lower limit of integration for quadrature j) $r_i$ weights for the point formula used

abscissas for the point formula used

$$\Delta = \frac{b - a}{NQ}$$

x,

(For a complete discussion of the Gauss-Legendre quadrature formula, see any numerical analysis text, such as ref. 5.)

The 6000 version of the subroutine obtains the first point in a quadrature, evaluates the function at that point, multiplies the function by the appropriate weight, and adds the product to a sum. It continues in the same manner until all the points have been evaluated and used in that quadrature; then the process is repeated for the remaining quadratures until the integral has been evaluated over the desired range. The sum is then multiplied by the appropriate delta to obtain the value of the integral.

Looking at this sequence of instructions shows that much of this work is actually computed independently. The STAR version takes advantage of this. All the information required to compute the points is known so all the points for all the quadratures can be computed to form a long array. This array can be passed to the subroutine to evaluate the functions and a savings can be obtained by using vector instructions to evaluate the functions. This routine will now return an array of function values. All the weights are known so an array can be formed which contains the weights for all the quadratures. The remaining task is to multiply the weights by the functions and sum the results. This computation can use the dot product macro; even though it is costly compared with most vector instructions, it is less expensive than a vector multiply followed by scalar adds. This formulation requires the use of the dot product only once. A scalar multiply of the sum by the appropriate delta completes the integral evaluation.

The weights and abscissas for each quadrature are identical; therefore, the length of these arrays in the 6000 version is equal to 10 for the 10-point formula. A version of the subroutine using short vectors could have been formed. The STAR version could have used vector instructions and computed results using one quadrature at a time with the same array of weights and array length as in the 6000 version. Computing the results one quadrature at a time means that the startup time associated with each vector instruction and the startup times used in the function evaluation also would be multiplied by the number of quadratures used. Depending upon the nature of the function, this amount of time could be significant.

Vector transmits could be used to create a long vector of weights which is really a repetition of the short vector. The long vector is

$$(\mathbf{r}_1, \mathbf{r}_2, \ldots, \mathbf{r}_m, \mathbf{r}_1, \mathbf{r}_2, \ldots, \mathbf{r}_m, \ldots, \mathbf{r}_1, \mathbf{r}_2, \ldots, \mathbf{r}_m)$$

where m depends upon the quadrature formula and q is the number of quadratures. The cost of setting up this vector seems to be minimal when compared with the multiplication of the startup times which would have prevailed in the short-vector version.

A problem of this type (excluding the function evaluation) would probably not produce a significant paging problem. All the arrays are one dimensional and will most likely not be extremely long. However, consideration should still be given to the working set which is created. Currently, there is a tendency to initialize variables at the beginning of the program. The vector of weights could be created at the beginning of the program, but they are not used until after the functions are evaluated; therefore, the vector is not created until just before it is used. Forming the vector of weights and then immediately using them results in less paging; therefore, a better working set is created than would have been created by initializing the vector.

Subroutine GLEGEN was adapted for the STAR computer by using the same numerical method as was used on the 6000. An efficient STAR routine was obtained by using long vectors and reordering the sequence of instructions. Figure 2 shows a flow chart of the 6000 and STAR versions.

Appendix B contains the coding of the two versions of the algorithm. The 6000 version and the STAR version in 6000 FORTRAN and STAR FORTRAN are included. Notice that the calling sequence of the STAR version reflects the use of long arrays in that more working space is needed.

#### CONCLUDING REMARKS

Two algorithms used on the Control Data Corporation 6000 computer were adapted for use on the STAR computer. This adaptation required a rethinking of the flow of the entire problem. Array variables are used where the 6000 code used a single-value variable and the steps in the algorithm are not computed in the same order for the two versions. This reordering of steps and changes in variable assignments allow vector instructions to be used and a reasonable working set can be established for the program; thereby efficient use of the STAR computer and some implied improvements for the Control Data Corporation 6000 computer are made.

Langley Research Center,

National Aeronautics and Space Administration, Hampton, Va., February 20, 1974.

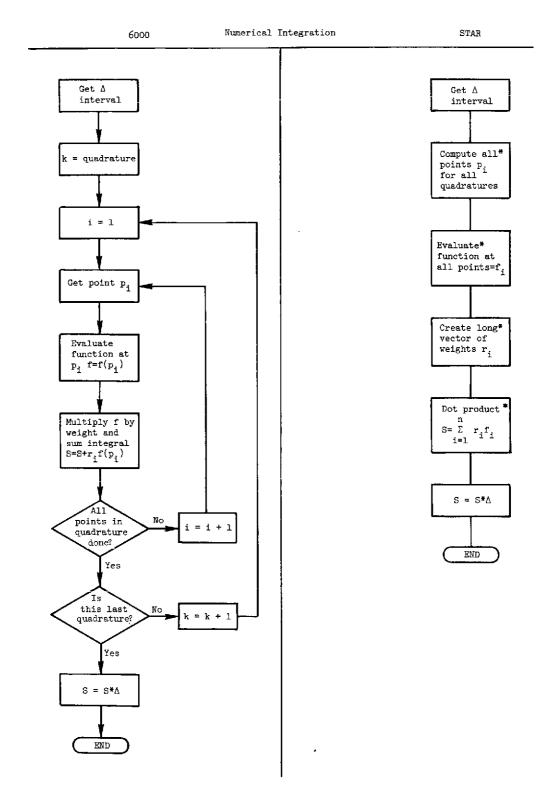


Figure 2.- Flow chart of 6000 and STAR versions of Gauss-Legendre quadrature formula for numerical integration. An asterisk denotes use of vector instructions.

APPENDIX A



# FORTRAN CODING OF 6000 AND STAR VERSIONS OF ALGORITHM FOR THE

# SOLUTION OF A SET OF SIMULTANEOUS EQUATIONS

### 6000 Version in 6000 FORTRAN

	PRESERVE AND AND A REPORT AND	GELIMOUI
	SUBROUTINE GELIM(A,N,B,NRHS,MAXN,IPIVUT,10P,DETERM,ISCALE)	GELIMOU2
	DIMENSION A (MAXN, MAXN), B (MAXN, NR HS), IPIVOT (MAXN)	GELINU03
3	A(MAXN, MAXN) = SQUARE MATRIX OF COEFFICIENTS(A IS DESTROYED)	GELIMOOD
2	N = NUMBER OF ROWS AND COLUMNS IN A	
5	B(MAXN, NRHS) = MATRIX OF CONSTANTS(REPLACED BY SOLUTION MATRIX)	GELIMOUS
3	NRHS = NUMBER OF COLUMNS IN B	GLLINUOU
5	MAXN = MAXIMUM NUMBER OF ROWS AND COLUMNS IN A	GELIMOU7
0	IPIVOT(MAXN) = RECORD OF POW INTERCHANGES	GELIMUOS
2	IOP + IOP=1, EVALUATE DETERMINANT. IOP=0, SKIP DETERMINANT EVALUATI	ON 009
5	DETERM - GIVES VALUE OF DETERMINAT(DETERM=(10**100)**ISCALE*DETER	M) 010
()	ISCALE = SCALE FACTOR COMPUTED BY SUBPOUTINE IN DETERM EVALUATION	GELIMUII
5	TO KEEP DETERM WITHIN THE FLUATING POINT WORD SIZE OF TH	E 012
2	COMPUTER	GELIMOI3
	5IGN=1.0	GELIM020
	00 35 I=1,N	GELIMO30
35	IPIVOT(I) = I	GELIMU40
-	$I SC \Delta L E = 0$	GELIMODO
	R1=10.0++100	GELIMQOO
	R 2=1.0/R 1	GELIM070
	DETERM=1.0	GELIMOdic
10	NMI = N-1	GEL IMÖYÖ
	IF(NM1.EQ.0) GU TO 1000	GELIMIJO
* * * * *	*FIND LARGEST PEMAINING TEPM IN I-TH CULUMN FOR PIVOT	GELIMIIG
	00 200 I=1.NM1	GELIM120
	816=0.	GELIM130
	00 50 K=I+N	GELIMI40
	TERM=ABS(A(K,I))	GELIMISU
	IF(TEPM-BIG)50.50.30	GELIMIOO
2.2	BIG=TERM	GELIM170
30	L=K	GELIMINO
	-	GEL1M190
50	CONTINUE IPIVOT(I)=L	GELIM200
		GELIM210
	IF(BIG)80,60,80	GEL 1MZZO
50	DETERM=0.0	GELIM225
	RETUPN	GELIM230
80	1F(1-L)90,120,90	GELIM240
	*SWAP ROWS OF A AND B. SET IPIVOT(J)=K.	GELIM250
90	SIGN=-SIGN	OLE INZ 20
	DO 100 J=1.N	GELIM270
	TEMP1=A(1,J)	GELIM280
	$\Lambda(I,J) = \Lambda(L,J)$	
100	∧ (L + J) = T EMP1	GELIM290
	DO 101 J=1+NPHS	GELIMINO
	TEMP2=R(I.J)	GELIMSIG
	$\mathbf{P}\left(1\cdot\mathbf{J}\right)=\mathbf{P}\left(1\cdot\mathbf{J}\right)$	GELIM320
101	B(L+J)=TEMP2	GELIM330
120	C ONTINUE	GELIM340
* * * * *	*STORE PIVOT IN A(I,J). MULTIPLY A AND B BY PIVOT.	GELIM350
	[P1=I+L	GEL 1M300
	00 31 JT=EP1+N	GELIM370
	A(I1,I) = A(I1,I) / A(I,I)	GELIM330
	X3=A(II,I)	GELIM3∀0
	00 32 K=IP1+N	GELIM400

# APPENDIX A – Continued

32	A(II,K)=A(II,K)-X3*A(J,K)	GÉLIM410
	00 33 K=1,NRHS	GEL1M420
33	B(II.K)=B(II.K)-X3*B(I.K)	GELIM430
31	CONTINUE	GELIM440
200	CONTINUE	GEL IM450
200		GELIM460
-	TF(JOP+EQ+0)60 TO 221	GELIM470
		-
	*SCALE THE DETERMINANT	GELIM450
2		GELIM490
	DO 122 [=1,N	GEL1M500
	PIVUTI=A(I,I)	GELIM510
1035	IF(ABS(DETERM)-R1)1030,1010,1010	GELIM520
1010	DETERM=DETERM/R1	GELIM530
	ISCALE= ISCALE+1	GELIM540
	IF(ABS(DETERM)-R1)1060,1020,1020	GELIM550
1620	DETERM=DETERM/R1	GELIMSUU
1020		GELIM5/0
	ISCALE=ISCALE+1	
	GU TO 1060	GELIM580
	IF(ABS(DETERM)-R2)1040,1040,1060	GELIM590
1040	DETERM=DETERM*R1	GELIMõJü
	I SCALE=I SCALE+1	GELIMOLU
	TF(ABS(DETERM)-R2)1050,1050,1060	GELIM620
1050	DETERM=DETERM*R1	GELIM636
_	ISCALE=ISCALE-1	GELIMo40
1.060	IF(ABS(PIVOTI)-R1)1090,1070,1070	GELIM650
	PIVOTI=PIVOTI/R1	GEL IM660
2010	I SCALE=I SCALE+1	GEL 1M676
	IF(ABS(PIVDT))-R1)320,1080,1080	GELIM600
1080	PIVOTI=PIVOTI/R1	GELIM690
1000	I SCALE= I SCALE+1	GEL IM700
	GO TO 320	GEL 1M710
1000	IF(ABS(PIVOTI)-R2)2000,2000,320	GELIM720
	PIVOTI=PIVOTI*R1	GELIM73Ú
7000	I SCALE=I SCALE+1	GELIM740
		GEL1M75u
	IF(ABS(PIVOTI)-P2)2010,2010,320	
2010	PIVOTI=PIVOTI*R1	GELIM760
	I SCALE=I SCALE=1	GELIM770
320	DETERM=DETERM*PIVOTI	GELIM780
122	CONTINUE	GEL IN790
С		GEL IM800
****	*PERFORM BACK SUBSTITUTION	GELIM810
221	CONTINUE	GEL IMB20
	D0 57 IC=1+NRHS	GELIN830
	B(N, TC) = B(N, TC) / A(N, N)	GELIM840
57	CONTINUE	GELIM850
2.	00 61 KK±1.NM1	GELIM860
	I=N-KK	
		GELIM870
		GEL IM880
	DO 61 J=1,NRHS	GEL IM900
	SUM=B(I.J)	GELIM910
	00 62 K=I1.N	GELIM920
	SUM=SUM-A(I,K)*B(K,J)	GEL IM930
61	B(I,J)=SUM/A(I,I)	GELIM940
	DETERM=DETERM=SIGN	GELIM950
	RETURN	GEL IM960
1000	IF(A(1,11,E0.0) GO TO 60	GELIM970
	DO 1500 J≠1,NRHS	GEL IM975
1500	8(1,J)=B(1,J)/A(1,1)	GELIM980
	DETERM=A(1,1)	GELIM990
	RETURN	GEL IM995
	END	GELM1000

	SUBROUTINE GELIMIA, N, B, NRHS, MAXN, IPIVOT, IOP, DETERM, ISCALE)	GIL60010
	DIMENSION A(MAXN,MAXN),B(MAXN,NRHS),IPIVOT(MAXN)	G1L60020
5	A(MAXN, MAXN) = SQUARE MATRIX OF COEFFICIENTS(A IS DESTROYED)	GIL60030
	N = NUMBER OF ROWS AND COLUMNS IN A	G1L60040
00000	BIMAXN, NRHS) = MATRIX OF CONSTANTS (REPLACED BY SOLUTION MATRIX)	GIL60050
2	NRHS = NUMBER OF COLUMNS IN 8	G1L60060
ž	MAXN = MAXIMUM NUMBER OF ROWS AND COLUMNS IN A	GIL 60070
ž	IPIVOT(MAXN) = RECORD OF ROW INTERCHANGES	GIL60080
-	IOP - IOP=1,EVALUATE DETERMINANT. IOP=0,SKIP DETERMINANT EVALUAT	
C	10P - 10P=1, EVALUATE DETERMINANT, 10P=0, SATE DETERMINANT EVALUATE	EBMC1140100
C	DETERM - GIVES VALUE OF DETERMINATIDETERM=(10*+100)++ISCALE+DET	
0 0 0	ISCALE = SCALE FACTOR COMPUTED BY SUBROUTINE IN DETERM EVALUATION	IN GILOUILU
2	TO KEEP DETERM WITHIN THE FLOATING POINT WORD SIZE OF	HEGILOUIZU
0	COMPUTER	G1L60130
	1 D≠ O	GIL60140
	S[GN=1.0	GIL60150
	DD 3 I=1.N	GIL60160
3	IPIVOT(I)=I	GIL60170
	I SCALE=0	GIL60180
	R1=10.0**100	GIL60190
	$R_{2}=1.0/R_{1}$	GIL60200
	DETERM=1.0	G1L60210
10		GIL60220
		G1L60230
	IF (NM1.E0.0) GO TO 600	GIL60240
	NPI=N+1	GL[60250
	<b>*FIND LARGEST REMAINING TERM IN I-TH COLUMN FOR PIVOT</b>	+
	DO 200 I=1,NM1	G1L60260
	B I G=0.	G1L60270
	00 50 K=I+N	GIL60280
	TERM=ABS(A(K.I))	G1L60290
	IF(TERM-BIG)50,50,30	G1L60300
30	BIG=TERM	G1L60310
	L=K	GIL60320
50	CONTINUE	G <b>[l60330</b>
	IPIVOT(I)≠L	GIL60340
	1F(8[G]80+60+80	GIL60350
60	DETERM=0.0	GIL60360
	RETURN	GIL60370
80	IF(I-L)90,120,90	GIL60380
	INTERCHANGE ELEMENTS IN COLUMN I ONLY	
90	SIGN=-SIGN	GIL60400
90	TEMP1 = A(1,1)	GIL60410
	$A(I \bullet I) = A(I \bullet I)$	GIL60420
		G1L60430
	A(L, 1)=TEMP1	G1L60440
120	CONTINUE	GIL60450
####	*STORE PIVOT IN A(I, J). MULTIPLY A AND B BY PIVOT.	
	fP1=I+1	GIL60460
5***1	+OBTAIN COLUMN OF PIVOTS	
	DO 128 K=IP1+N	GIL60470
128	3 A(K+I)=A(K+I)/A(I+I)	GIL60480
0***4	**INTERCHANGE ELEMENTS ONLY IN COLUMN J	
	DO 132 J≠1+NRHS	GIL60490
	TEMP1= B(1+J)	GIL60500
	B(I+J) = B(L+J)	GIL60510
	B(L,J) = TEMPL	GIL60520
	00 130 K=IP1+N	GIL60530
120	) B(K,J)=B(K,J) - A(K,J)*B(I,J)	GIL60540
	2 CONTINUE	GIL60550
1 J C 1 ± ± ± *	**INTERCHANGE ELEMENTS ONLY IN COLUMN J	
	DO140 J=IP1.N	GIL60560
	TEMP1 = A(I,J)	GIL60570
	$A(I \bullet J) = A(I \bullet J)$	GIL60580
	A11+J-A1417	GIL60590
	A(L,J)=TEMP1	GIL60600
	D0135 K±IP1+N	9100000

135	A {K, J}=A{K, J}~A{K,I}*A{I, J}	GIL60610
140	CONTINUE	GIL60620
	IF (10P.EQ.0) GO TO 200	GIL60630
0	· · · · · · · · · · · · · · · · · · ·	GIL60640
-	*SCALE THE DETERMINANT	G1L60650
C		GIL60660
	PIVOTI=A(I.I)	GIL60670
	IF(ABS(DETERM)-R1)1030,1010,1010	GIL60680
1010	DETERM=DETERM/R1	GIL60690
	I SCALE= I SCALE+1	GIL60700
	IF(ABS(DETERM)-R1)1060,1020,1020	GIL60710
1 02 0	DETERM=DETERM/RL	G1L60720
	I SCALE=I SCALE+1	GIL60730 GIL60740
1020	GO TO 1060 If[Abs(determ)-r2)1040+1040+1060	G1L60750
		GIL60760
1040	DETERM=DETERM=R1 ISCALE=ISCALE-1	G1L60770
	IF(ABS(DETERM)-R2)1050,1050,1060	GIL60780
1.050	DETERM=DETERM=R1	GIL60790
1000	I SCALE=I SCALE=1	GIL60800
1.060	IF(ABS(PIVOTI)-R1)1090,1070,1070	GIL60810
	PIVOTI=PIVOTI/R1	G1L60820
1010	ISCALE=ISCALE+1	GIL60830
	IF(A8S(PIVOTI)-R1)320+1080+1080	GIL60840
1080	PIVOTI=PIVOTI/R1	GIL60850
	I SCALE= I SCAL E+1	GIL60860
	GO TO 320	GIL60870
1090	[F(ABS(PIVOTI)-R2)2000,2000,320	GIL60880
	PIVOTI=PIVOTI*R1'	GIL60890
	ISCALE=ISCALE-1	GIL60900
	IF(ABS(PIVOTI)-R2)2010,2010,320	GIL60910
2010	PIVOTI=PIVOTI*R1	GIL60920
	I SCALE=I SCALE+1	GIL60930
****	*OBTAIN PARTIAL PRODUCT FOR DETERMINANT	
320	DETERM=DETERM*PIVOTI	GIL60940
	IF (I.NE.NMI) GO TO 200	GIL60950
	IF (ID.E0.1) GO TO 200	GIL60960
	PIVOTI=A(N.N)	GIL60970
	1 D= 1	GIL60980
	GD TO 1005	GIL60990
200	CONTINUE	G1L61000
	40-00-004 04-04 CHR (717104	GIL61010
	*PERFORM BACK SUBSTITUTION CONTINUE	GIL61020
540	DO 450 KK=1.N	GIL61030 GIL61040
	K=NP1-KK	GIL61050
[ ****	*SCALAR DIVIDE	01201000
	DO 430 J=1.NRHS	G1L61060
	B(K,J)=B(K,J)/A(K,K)	G1L61070
	IF (K.EQ.1) GO TO 430	GIL61080
	KM1=K-1	G1L61090
C ****	COLUMN K BY UNKNOWN AND SUBTRACT	
	DO 420 I=1.KM1	GIL61100
	B(I,J) = B(I,J) - A(I,K) + B(K,J)	G1161110
	CONTINUE	GIL61120
450	CONTINUE	GIL61130
	DETERM=DETERM*SIGN	GIL61140
	RETURN	GIL61150
600	IF (A(1,1).E0.0) GD TO 60	GIL61160
	DO 800 J=1, NRHS	G[L61170
800	B(1+J) = B(1+J)/A(1+1)	GIL61180
	DETERM=A(1+1) Return	GIL61190
	END	GIL61200
		GIL61210

## STAR Version in STAR FORTRAN

	SUBROUTINE GELIM(A.N.B.NRHS.MAXN, IPIVOT, IDP, DETERM, ISCALE)	GLISOOLO
	DIMENSION A(MAXN,MAXN),B(MAXN,NRHS),IPIVOT(MAXN)	GLI \$0020
2	A{MAXN,MAXN) = SOUARE MATRIX OF COEFFICIENTS(A IS DESTROYED)	GLI \$0030
0	N = NUMBER OF ROWS AND COLUMNS IN A	GLI 50040
	B(MAXN, NRHS) = MATRIX OF CONSTANTS(REPLACED BY SOLUTION MATRIX)	GLI \$0050
÷.	NRHS = NUMBER OF COLUMNS IN B	GLI \$0060
Ĩ.	MAXN = MAXIMUM NUMBER OF ROWS AND COLUMNS IN A	GLIS0070
-	IPIVOT(MAXN) = RECORD OF ROW INTERCHANGES	GL150080
ž	IOP - IDP=1, EVALUATE DETERMINANT. IOP=0, SKIP DETERMINANT EVALUAT	
0000000000	DETERM - GIVES VALUE OF DETERMINANT. IDP-D, SKIP DETERMINANT EVALUAT	
-	DETERM - GIVES VALUE OF DETERMINATOREREMINATORIA	KHGLISUIUU
ž	ISCALE = SCALE FACTOR COMPUTED BY SUBROUTINE IN DETERM EVALUATIO	N GLISUIIU
5	TO KEEP DETERM WITHIN THE FLOATING POINT WORD SIZE OF T	
2	COMPUTER	GLIS0130
	I D=O	GLI\$0140
	SIGN=1.0	GEISO150
	00 3 I=1+N	GLIS0160
3	IPIVDT(I)=I	GLI \$0170
5	I SCALE=0	GLISO180
6	R1=10.0**100	GLIS0190
	R2=1.0/R1	GLI SU200
	DETERM=1.0	GLIS0210
	NM1=N+1	GL150220
	IF {NML.EQ.0} GO TO 600	GLIS0230
	NP1=N+1	GLIS0240
		GLIS0240
	*FIND LARGEST REMAINING TERM IN I-TH COLUMN FOR PIVOT	
	DO 200 I=1.NM1	GLISO260
5		GLISO2/0
C NEE	D VECTOR ABSOLUTE , MAXIMUM BECAUSE YOU ARE SEARCHING COLUMN	GLISO280
Ċ.		GLISU290
	BIG=0.	GL150300
	DO 50 K=I+N	GLISO310
	TERM=ABS(A(K+I))	GL I SO320
	IF(TERM-BIG)50,50,30	GL1S0330
30	BIG≠TERM	GLI \$0340
	L=K	GLI \$0350
50	CONTINUE	GL I SO 360
	IPIVOT(I)=L	GL1S0370
	IF(BIG)80.60.80	GLI 50380
60	DETERM=0.0	GL1 \$0390
	RETURN	GL1S0400
80	IFII-L190,120,90	GLIS0410
90	SIGN=-SIGN	GLI \$0430
• •	TINTERCHANGE ELEMENTS IN COLUMN I ONLY	00130430
<b>, ***</b> *		C1 1 C0 64 0
	TEMP1= A(I,I)	GLIS0440
	A(I,I) = A(L,I)	GLIS0450
	A(L,I)=TEMP1	GLI \$0460
120		GLIS0470
C ****	*STORE PIVOT IN A(I,J). MULTIPLY A AND B BY PIVOT.	GLI S0480
	IP1=I+1	GLI S0490
C ****	*OBTAIN COLUMN OF PIVOTS	
	A(IP1:N+I)= A(IP1:N+I)/A(I+I)	GLIS0500
:****	*INTERCHANGE ELEMENTS ONLY IN COLUMN J	
-	DOI31 J=1+NRHS	GLIS0510
	TEMPL= B(I.J)	GLI S0520
	B(1+J)=B(L+J)	GL I S 0 5 3 0
	B(L.J)=TEMP1	GLIS0540
131	B(IP1:N,J) = B(IP1:N,J) - A(IP1:N,I) * B(I,J)	GLI S0550
	*INTERCHANGE ELEMENTS ONLY IN COLUMN J	
5 + + + + +	DO134 J=IPL.N	GLI 50560
	$TEMP1= A(I_*J)$	
		GLIS0570
	A(I,J) = A(L,J)	GLISO580
	A(L + J) = TEMP1	GL I S0590

	C1 160(00
134 A(IP1:N.J)= A(IP1:N.J) -A(IP1:N.I)* A(I.J)	GLISO600
IF (IOP.EQ.0) GD TO 200	GLIS0610
	GLIS0620
C++++SCALE THE DETERMINANT	GLI50630
	GLIS0640
C C++++SCALE THE DETERMINANT PIVOTI=A(I,I) 1005 IF(ABS(DETERM)-R1)1030,1010,1010	GLIS0650
1010 DETERM=DETERM/R1	GL150670
I SCALE= I SCALE+1	GLIS0680
[F(ABS(DETERM)-R1)1060,1020,1020	GLIS0690
1020 DETERM=DETERM/R1	GL150700
I SCALE= I SCALE+1	GLIS0710
GO TO 1060	GLIS0720
<pre>ISCALE=ISCALE+1 IF(ABS(DETERM)-R1)1060,1020,1020 1020 DETERM=DETERM/R1 ISCALE=ISCALE+1 G0 T0 1060 1030 IF(ABS(DETERM)-R2)1040,1040,1060 1040 DETERM=DETERM#R1</pre>	GLI \$0730
1040 DETERM=DETERM*R1	
I SCALE= I SCALE-1	GL1S0750
IF(ABS(DETERM)-R2)1050,1050,1060	GL150760
1050 DETERM#DETERM#R1	GLI 50770
I SCALE=I SCALE-1	GLISO780
1060 IF(ABS(PIVOTI)-R1)1090,1070,1070	GL I S0790
1070 PIVOTI=PIVOTI/R1	GLI SUBOO
I SCALE= I SCALE+ 1	GLI SO810
IF(ABS(PIVOTI)-R1)320,1080,1080	GLIS0820
1080 PIVOTI=PIVOTI/R1	GLISOB30
t SCALE=1SCALE+1	GLIS0840
GO TO 320	GL I \$0850
1090 IF(ABS(P1VDTI)-R2)2000,2000,320	GL I S0860
2000 Ptvoti=Pivoti*R1	GLISU870
I SC ALE=I SCALE-1	GLI SO880
IF(ABS(PIVOTI)-R2)2010,2010,320	GL I S0890
2010 PIVOTI=PIVOTI*RI	GLI 50900
I SCALE= I SCALE-1	GLI 50910
S#####OBTAIN PARTIAL PRODUCT FOR DETERMINANT	
320 DETERM=DETERM*PIVOTI	GLIS0920
IF (I.NE.NML) GO TO 200	GL150930
IF (ID.EQ.1) GO TO 200	GL1S0940
PIVOTI=A(N+N)	GL150950
I D= 1	GLIS0960
GO TO 1005	GL1S0970
200 CONTINUE	GE1 50980
	GLIS0990
C****PERFORM BACK SUBSTITUTION	GLISIOUO
340 CONTINUE	GLISI010
DO 450 KK=1.N	GLIS1020
K=NP1-KK	GLIS1030
C*****SCALAR DIVIDE	-
DO 430 J=1,NRHS	GLI 51040
B(K+J)=B(K+J)/A(K+K)	GLIS1050
IF (K.EQ.1) GO TO 430	GLIS1060
KM1=K-1	GLIS1070
3*****COLUMN K BY UNKNOWN AND SUBTRACT	
B(1:KM1,J)= B(1:KM1,J) -A(1:KM1,K)+ B(K,J)	GLIS1080
430 CONTINUE	GLI S1090
450 CONTINUE	GLIS1100
DETERM=DETERM*SIGN	GLIS1110
RETURN	GLIS1120
600 IF (A(1,1).EQ.0) GO TO 60	GLIS1130
B(1,1:NRHS)= B(1,1:NRHS)/A(1,1)	GLIS1140
DETERM=A(1.1)	GLIS1150
RETURN	GLIS1160
END	GLIS1170

#### APPENDIX B

# FORTRAN CODING OF 6000 AND STAR VERSIONS OF ALGORITHM

## FOR NUMERICAL INTEGRATION EVALUATION

## 6000 Version in 6000 FORTRAN

- <u>.</u>	SUBROUTINE GLEGEN	(SUM+1CODE+A+B+FX+F0FX+NFX+NQ+NPQ) ************************************	GLEGN	1
- 0 * - 0 *	*****		GLEGN	2 3
				_
*	PURPUSEI	TO COMPUTE THE INTEGRALS, F(I) OF X * DX FROM *		4
 			GLEGN	5
-			GLEGN	6
2*			GLEGN	7
C*	USE 2	CALL GLEGEN (SUM, ICODE, A, B, FX, FOFX, NFX, NO, NPO)*		8
C *			GLEGN	9
	PARAMETERS:		GLEGN	
0*	SUM	THE ARRAY FOR THE VALUE(S) OF THE INTEGRAL(S) *		
<u>0</u> *			GLEGN	
0*	ICODE	AN INTEGER TEST CODE RETURNED BY THE ROUTINE *		
0*			GLEGN	-
<u>^</u> *			GLEGN	-
<u>]</u> *		= 2. NPQ NOT PROPERLY SPECIFIED (NOT 3 OR 10) *	GLEGN	16
<u>°</u> *		= 3, NO NOT PROPERLY SPECIFIED *	GLEGN	17
C.*		*	GLEGN	18
0*		THIS PARAMETER SHOULD BE TESTED UPON RETURN *	GLEGN	19
C *			GLEGN	20
C *		*	GLEGN	21
0*	A	LOWER LIMIT OF INTEGRATION *	GLEGN	22
-0*		*	GLEGN	23
Ç*	8	UPPER LIMIT OF INTEGRATION *	GLEGN	24
÷.		*	GLEGN	25
€*	FX	THE NAME OF A USER SUPPLIED SUBROUTINE WITH *		
C*		ARGUMENTS X AND FOFX TO EVALUATE THE FUNCTIONS*	GLEGN	27
C*		IT MUST BE DECLARED AS EXTERNAL. *	GLEGN	28
:*			GLEGN	
::*	FOFX	THE ARRAY TO STORE THE VALUES OF THE FUNCTIONS*	GLEGN	30
0*		*	GLEGN	31
0*	NEX	AN INTEGER, THE NO. OF FUNCTIONS TO INTEGRATE *	GLEGN	32
C+		*	GLEGN	33
<b>C</b> *	NQ	AN INTEGER, THE NUMBER OF QUADRATURES *	GLEGN	34
0*			GLEGN	
0*	NPQ	AN INTEGER. THE NO. OF POINTS PER QUADRATURE *	GLEGN	36
0*		IT MUST BE 3 OR 10 *	GLEGN	37
<u>:</u> *			GLEGN	
5*	REQUIRED ROUTINES:		GLEGN	
÷.				-
Ċ*	SOURCE/IMPLEMENTER:	* COMPUTER SCIENCES CORP./ G. W. HAIGLER *	GLEGN	41
÷\$		*	GLEGN	42
Č*	DATE RELEASED:		GLEGN	
0*		*	GLEGN	
C+	LATEST REVISION:		GLEGN	
-				

# APPENDIX B – Continued

+	* * * * * * * * * * * * * * * * * * * *		
5		GLEGN	
	DIMENSION U1(3),U2(10),R1(3),R2(10),U(13),R(13),SUM(1),FOFX(1)	GLEGN	
<u>,</u>		GLEGN	
	EQUIVALENCE (U1(1),u(1)),(U2(1),U(4)),(R1(1),R(1)),(R2(1),R(4))	GLEGN	
5		GLEGN	
	DATA U1/.1127016653792595887298334620742/.U2/.013046735741414		
	1.06/4683166555071602952158504882833023029353764255628305091		
	2557443 <b>7</b> 16949081671669 <b>76</b> 9706462483970478414951293253168334		
	3493+ • 986 953264258586/ • R 1/ • 277777777777778 • • 44444444444444444. • • 27777	7GLEGN	55
	4777777778/.R2/.033335672154344.074725674575291109543181257991.		
	5.1346333596549981477621123573771477621123573771346333596549	9GLEGN	57
(	68109543181257991074725674575291033335672154344/	GLEGN	58
C		GLEGN	59
	ICODE = 0	GLEGN	60
	IF(NFX.LE.O)ICODE = 1	GLEGN	61
	IF(NO.LE.O)ICODE = 3	GLEGN	62
	IF(NPQ.EQ.3)GO TO 5	GLEGN	63
	$IF(NPQ \cdot NE \cdot LO)ICODE = 2$	GLEGN	64
5	IF(ICODE.GT.O)RETURN	GLEGN	65
	£=1	GLEGN	66
	IF (NPQ .EQ. 3) J=O	GLEGN	67
	DELT= (B-A)/FLDAT(NQ)	GLEGN	68
	DO 10 I = $1.NFX$	GLEGN	69
10	SUM(1) = 0.0	GLEGN	70
	00 80 K=1+NQ	GLEGN	71
	XI = K → 1	GLEGN	72
	FF = A + XI*OELT	GLEGN	73
	DO 80 L=1,NPQ	GLEGN	74
	UU=U(J+L)*DELT+FF	GLEGN	75
	CALL FX(UU+FOFX)	GLEGN	76
	FACT=R(J+L)	GLEGN	
	DD 80 J8=1.NFX	GLEGN	
80	SUM(JB) =FOFX(JB)+FACT+SUM(JB)	GLEGN	
. +	D3 100 I I=1, NFX	GLEGN	
100	SUM(II) = SUM(II)*DELT	GLEGN	
	RETURN	GLEGN	
	END	GLEGN	- +

# STAR Version in 6000 FORTRAN

		NS(SUM,ICODE,A,B,FX,FOFX,NFX,NQ,NPQ,UU,FF,UD,MAXI)GLENS001
0.# 0.#	*** ** ** ** * * * * * * * * * * * * * *	**************************************
5= 5=	PURPOSE:	TO COMPUTE THE INTEGRALS, F(I) OF X * DX FROM * GLENSOOS
C*	FORT DEL.	A TO B USING THE GAUSS-LEGENDRE QUADRATURE * GLENSOO5
- 0 *		FORMULA. + GLENSOD6
ű÷		* GLENSOO7
	USE :	CALL GLEGEN (SUM, ICODE, A, B, FX, FOFX, NFX, NQ, NPQ) * GLENSOOB
	032.	CALL GLEGEN (SOMITCODE, AFBIFATORAINEATHEINEAT
C*	PARAMETERS:	* GLENSOLO
2*	SUM	THE ARRAY FOR THE VALUE(S) OF THE INTEGRAL(S) * GLENSO11
	304	THE ARRAT FOR THE VALUESSI OF THE INTEGRALIST + GLENSOII + GLENSOII
	ICODE	AN INTEGER TEST CODE RETURNED BY THE ROUTINE * GLENSO12
*	TCODE	= 0, NORMAL RETURN + GLENSOIS
0= 0+		
0.+ C.+		
3. 10. 10.		
0.+ 0.+		+ GLENSOLB
0+ 0+		THIS PARAMETER SHOULD BE TESTED UPON RETURN + GLENSO19
0.4 C*		BY THE CALLING PROGRAM + GLENSO20
U #		* GLENS021

22

# APPENDIX B - Continued

_C*	А	LOREN EINIT OF INTEDNATION	GLENS022
C *			GLENS023
- <b>C</b> *	В		GLENS024
<b>C</b> *			GLENS025
5*	FX		GLENS026
C*		ARGUMENTS X AND FOFX TO EVALUATE THE FUNCTIONS*	
C*		IT NOOT DE DECEMBED NO ENTENNIOT	GLENS028
C *			GLENS029
	FOFX	A TWO DIMENSIONED ARRAY TO STORE VALUES OF THE	
		FUNCTIONS, FOFX(MAXI,NFX) IN CALLING PROGRAM	GLENS031
C *			GLENS032
C *	NFX	AN INTEGER, THE NO. OF FUNCTIONS TO INTEGRATE *	
3*			GLENS034
C *	NO		GLENS035
C *			GLENS036
C *	NPO		GLENS037
C *			GLENS038
C.*			GLENS039
C *	UU	ARRAY TO STORE ALL POINTS AT WHICH FUNCTION IS	
0*		EVALUATED ,ALSO STORES WEIGHTS	GLENS041
C*	FF	ARRAY OF WORKING SPACE, CONTAINS LOWER LIMIT OF	
C.*		INTEGRATION FOR EACH QUADRATURE	GLENS043
C*	UD	ARRAY OF WORKING SPACE, ABSCISSAS* DELT	GLENS044
C \$	MAXI	MAXIMUM ROW DIMENSIONS OF FOFX IN CALLING PROG.	
51			GLE6S046
C#	AFAUTAER AAUTHER-		GLE6S047
	REQUIRED ROUTINES:		GLE6S048
	CONDER (THOLENENTED -		GLE6S049
	SUURCE/IMPLEMENTER:		GLE6S050
0*	ONTE DELEACED-		GLE6S051
	DATE RELEASED:	NOV. 14.1972 *	GLE6\$052
~ *		*	C1 C4 C05 3
0* **	LATEST BENISTON.		GLE6S053
Č*		NOV. 15,1972 *	GLE6S054
C* C*			GLE65054 GLE65055
Č*	* * * * * * * * * * * * * * * * * * * *	NOV. 15,1972 * ***********************************	GLE65054 GLE65055 GLE65056
C* C*	**************************************	NOV. 15.1972 ** **********************************	GLE6S054 GLE6S055 GLE6S056 GLE6S057
C* C* C	**************************************	NOV. 15,1972 * ***********************************	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058
C* C*	**************************************	NOV. 15,1972 ** **********************************	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058 GLE6S059
C* C* C	**************************************	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058 GLE6S059 GLE6S060
C* C* C	**************************************	<pre>NOV. 15.1972 ** **********************************</pre>	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061
C* C* C	**************************************	NOV. 15,1972 ** **********************************	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 , GLE6S062
C* C* C	**************************************	<pre>NOV. 15.1972 ** **********************************</pre>	GLE6S054 GLE6S055 GLE6S056 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 • GLE6S062 8GLE6S063
C* C* C	**************************************	<pre>NOV. 15.1972 ** **********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S062 8GLE6S063 4GLE6S064
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S061 GLE6S063 4GLE6S064 7GLE6S065 GLE6S065
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S061 GLE6S063 4GLE6S064 7GLE6S065 GLE6S065
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S061 GLE6S063 4GLE6S064 7GLE6S065 GLE6S065
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S059 GLE6S060 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S065 GLE6S065 GLE6S066 9GLE6S067
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S059 GLE6S060 GLE6S061 GLE6S063 4GLE6S063 4GLE6S064 7GLE6S065 GLE6S066 9GLE6S067 GLE6S068
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 * ***********************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S059 GLE6S060 GLE6S061 ,GLE6S062 BGLE6S063 4GLE6S065 GLE6S065 GLE6S066 9GLE6S067 GLE6S068 GLE6S069
C* C* C	<pre>************************************</pre>	<pre>N0v. 15,1972 ************************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S059 GLE6S060 GLE6S061 ,GLE6S062 BGLE6S063 4GLE6S065 GLE6S065 GLE6S066 9GLE6S067 GLE6S069 GLE6S069 GLE6S070
C* C* C	<pre>************************************</pre>	<pre>NOV. 15.1972 ************************************</pre>	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S062 BGLE6S063 4GLE6S065 GLE6S065 GLE6S066 9GLE6S066 9GLE6S067 GLE6S069 GLE6S070 GLE6S071
C* C* C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 ,GLE6S062 BGLE6S063 4GLE6S064 7GLE6S065 GLE6S066 9GLE6S066 9GLE6S067 GLE6S069 GLE6S070 GLE6S072
C* C* C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 ,GLE6S062 BGLE6S063 4GLE6S064 7GLE6S065 GLE6S066 GLE6S066 GLE6S067 GLE6S069 GLE6S070 GLE6S072 GLE6S073
C* C* C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S065 GLE6S065 GLE6S066 9GLE6S068 GLE6S069 GLE6S069 GLE6S070 GLE6S072 GLE6S073 GLE6S074
C* C* C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S065 GLE6S065 GLE6S066 9GLE6S066 GLE6S068 GLE6S069 GLE6S069 GLE6S070 GLE6S071 GLE6S073 GLE6S074 GLE6S075
C* C* C	<pre>DIMENSION U1(3).U 1 FOFX(MAXI.1 ). EQUIVALENCE (U1(1) DATA U1/.11270166 1.067468316655507. 25.57443716949081 3493.986953264258 477777778/.R2/.03 5.134633359654998. 68.10954318125799 ICODE = 0 NT = TOTAL NUMBER 0 NT=NQ*NPQ IF(NFX.LE.0)ICODE IF(NQ.LE.0)ICODE IF(NPQ.E0.3)GO TO IF(NPQ.E0.3)GO TO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NP</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S065 GLE6S066 9GLE6S066 GLE6S070 GLE6S071 GLE6S073 GLE6S074 GLE6S075 GLE6S076
C* C* C	<pre>DIMENSION U1(3).U 1 FOFX(MAXI.1 ). EQUIVALENCE (U1(1) DATA U1/.11270166 1.067468316655507. 25.57443716949081 3493.986953264258 477777778/.R2/.03 5.134633359654998. 68.10954318125799 ICODE = 0 NT = TOTAL NUMBER 0 NT=NQ*NPQ IF(NFX.LE.0)ICODE IF(NQ.LE.0)ICODE IF(NPQ.E0.3)GO TO IF(NPQ.E0.3)GO TO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NPQ.E0.3)FO IF(NP</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S060 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S065 GLE6S066 9GLE6S066 GLE6S070 GLE6S071 GLE6S073 GLE6S074 GLE6S075 GLE6S076 GLE6S077
C C C C C C C C C C C C C C C C C C C	<pre>DIMENSION U1(3).U DIMENSION U1(3).U DIMENSION U1(3).U DATA U1/.11270166 .067468316655507. 2557443716949081 3493986953264258 4777777778/.R2/.03 5.134633359654998. 6810954318125799 ICODE = 0 NI = TOTAL. NUMBER 0 NT=NQ*NPQ IF(NFX.LE.0)ICODE IF(NO.LE.0)ICODE IF(NO.LE.0)ICODE IF(NO.LE.0)ICODE IF(NO.NE.10)ICOD 5 IF(ICODE.GT.0)RET J=3 IF (NPQ .E0. 3 DELT= (B-A)/FLOAT</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S060 GLE6S061 GLE6S062 8GLE6S063 4GLE6S064 7GLE6S064 7GLE6S065 GLE6S066 GLE6S069 GLE6S070 GLE6S071 GLE6S074 GLE6S075 GLE6S077 GLE6S078
C C C C C C C C C C C C C C C C C C C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S060 GLE6S061 GLE6S063 4GLE6S063 4GLE6S064 7GLE6S065 GLE6S066 GLE6S067 GLE6S070 GLE6S071 GLE6S073 GLE6S075 GLE6S076 GLE6S078 GLE6S078 GLE6S078 GLE6S079
C C C C C C C C C C C C C C C C C C C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S057 GLE6S058 GLE6S060 GLE6S060 GLE6S061 GLE6S063 4GLE6S063 4GLE6S064 7GLE6S065 GLE6S066 GLE6S067 GLE6S070 GLE6S071 GLE6S073 GLE6S075 GLE6S076 GLE6S078 GLE6S078 GLE6S078 GLE6S079
C C C C C C C C C C C C C C C C C C C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 GLE6S062 BGLE6S063 4GLE6S065 GLE6S065 GLE6S066 9GLE6S066 9GLE6S066 GLE6S070 GLE6S071 GLE6S072 GLE6S073 GLE6S075 GLE6S076 GLE6S077 GLE6S077 GLE6S078 GLE6S079 GLE6S079 GLE6S079 GLE6S081 GLE6S081 GLE6S082
C C C C C C C C C C C C C C C C C C C	<pre>************************************</pre>	NOV. 15.1972 ************************************	GLE6S054 GLE6S055 GLE6S057 GLE6S058 GLE6S059 GLE6S060 GLE6S061 , GLE6S062 BGLE6S063 GLE6S064 7GLE6S065 GLE6S066 9GLE6S066 GLE6S067 GLE6S070 GLE6S072 GLE6S074 GLE6S075 GLE6S077 GLE6S077 GLE6S077 GLE6S078 GLE6S078 GLE6S078 GLE6S079 GLE6S081

# APPENDIX B - Continued

C*****COMPUTE FIRST PDINT IN EACH QUADRATURE				
	DO 30 K=1.NO	GLE6S084		
	XI =K-1	GLE65085		
	30 FF{K} =A + XI*DELT	GLE65086		
5	CREATE VECTOR OF ALL POINTS FOR ALL QUADRATURES	GLE65087		
	D0 40 K=1+NQ	GLE6S088		
	K1 = (K-1) * NPQ	GLE 6508 9		
	DO 35 I=1.NPQ	GLE6S090		
	35 UU(KI+I) = FF(K) + UD(I)	GLE6S091		
	40 CONTINUE	GLE6S092		
0	EVALUATE FUNCTION AT ALL POINTS	GLE6\$093		
0	NOTE CHANGE IN CALLING SEQUENCE FOR FUNCTION EVALUATION SUBROUTINE	GLE6S094		
	CALL FXS(UU,FOFX,MAXI,NT)	GLE65095		
0	CREATE VECTOR OF WEIGHTS	GLE6SD96		
	DO 50 K=1.NQ	GLE6S097		
	K1 =1K-1)*NPQ	GLE6S098		
	DO 45 I= 1.NPQ	GLE6S099		
	45 UU(K1+I) = R(J+I)	GLE6S100		
	50 CONTINUE	GLE65101		
	DO 80 I=1.NFX	GLE6S102		
	SUM(I)=0.0	GLE65103		
	D0 70 K=1+NT	GLE6S104		
	70 SUM(I) = SUM(I) + UU(K) * FOFX(K,I)	GLE6S105		
	SUM(I) = DELT * SUM(I)	GLE6S106		
	80 CONTINUE	GLE65107		
	RETURN	GLE6S108		
	END	GLE6S109		

# STAR Version in STAR FORTRAN

		S{SUM, ICODE, A, B, FX, FOFX, NFX, NQ, NPQ, UU, FF, UD, MAXI *****	
· ; ≈≠≠≈≠≈≠≈≠≈≠≠≈ ; ¢	* * * * * * *		GLE6S002
L* C* PURPOSE:		TO COMPUTE THE INTEGRALS, F(1) OF X # DX FROM *	
C* PUKPUSE: C*			GLE65005
0* 0*			GLE65006
0.* 5.*			GLE65007
C* USE:		CALL GLEGEN (SUM.ICODE, A.B. FX.FOFX.NFX.NO.NPQ)*	
C+ 036-			GLE65009
C* PARAMETERS:			GLE65010
S≠ FARAMETERS+	SUM	THE ARRAY FOR THE VALUE(S) OF THE INTEGRAL(S) *	
C *	304		GLE6S012
	CODE		GLE6S013
- C*	0000		GLE6S014
C #			GLE6S015
C.*		•••••••••••••••••••••••••••••••••••••••	GLE6S016
C.*			GLE65017
 C*			GLE6S018
 *		THIS PARAMETER SHOULD BE TESTED UPON RETURN *	GLE6S019
_ _ *		BY THE CALLING PROGRAM *	GLE6S020
C *			GLE6S021
C.+	4	LOWER LIMIT OF INTEGRATION *	GLE6S022
C.*		*	GLE6S023
C #	8	UPPER LIMIT OF INTEGRATION *	GLE6S024
C #		*	GLE6S025
) C * C * C * C *	FX	THE NAME OF A USER SUPPLIED SUBROUTINE WITH *	GLE6S026
0 <b>*</b>		ARGUMENTS & AND FOFX TO EVALUATE THE FUNCTIONS*	GLE6S027
- : * : *		IT MUST BE DECLARED AS EXTERNAL. *	GLE65028
C *		*	GLE6S029
	FOFX	A TWO DIMENSIONED ARRAY TO STORE VALUES OF THE	GLE6S030
		FUNCTIONS, FOFX(MAXI,NFX) IN CALLING PROGRAM	GLE6S031

# APPENDIX B - Continued

C.*	±	GLE6S032
C* NFX	AN INTEGER. THE NO. OF FUNCTIONS TO INTEGRATE *	
		GLE6S034
C* NO		GLE6S035
0.00 C *	RA INTEGENT THE NORDER OF EDROHMFORED	GLE6S036
C* NPQ	AN INTEGER. THE NO. OF POINTS PER QUADRATURE *	
		GLE65038
	ARRAY TO STORE ALL POINTS AT WHICH FUNCTION IS	GLE6S039
0.4		GLE6S040
	ARRAY OF WORKING SPACE, CONTAINS LOWER LIMIT OF	GLE65041
 C*	INTEGRATION FOR EACH QUADRATURE	GLE6\$042
C* UD	ARRAY OF WORKING SPACE, ABSCISSAS* DELT	GLE6S043
Č*		GLE6S044
C* MAXI	MAXIMUM ROW DIMENSIONS OF FOFX IN CALLING PROG.	GLE65045
C*		GLENS046
C* REQUIRED ROUTINES:	NONE	GLENSO47
C*		GLENS048
C* SOURCE/IMPLEMENTER:		GLENS049
C*		GLENS050
C* DATE RELEASED:		GLENS051
C.≉		GLENS052
C* LATEST REVISION:	NOV. 15.1972 *	GLENS053
-	*****	
C		GLENS055
	2(10),R1(3),R2(10),U(13),R(13),SUM(1),	GLENS056
	UU(1 ),UD(1 ),FF(1 )	GLENS057
<b>.</b>		GLENS058
	}+U(1)}+(U2(1)+U(4)}+(R1(1)+R(1))+(R2(1)+R(4))	GLENS059
		GLENS060
DATA U17+11270166	15379259+ • 5+ • 887298334620742/ •U2/ •013046735741414 • 160295215850488+ • 283302302935376+ • 4255628305091	OLENSUOL
1.00(408310033007+	• 16029 521 5 8504 8 8 • • 2 8 5502 502 95 5 7 8 • • • 2 5 5 8 50 9 1 6 • • 71669 769 7064 6 24 • • 8 39 704 784 149 512 • • 9 3 25 316 8 33 4	401ENS062
201020000000000000000000000000000000000	1586/+R1/•277777777777778+•4444444444444444444444	TGLENSOGA
477777778/.92/.03	3335672154344,.074725674575291109543181257991.	GLENS065
	+147762112357377++147762112357377++1346333596549	
	1+.074725674575291033335672154344/	GLENS067
5		GLENS068
ICODE = 0		GLENS069
	F POINTS AT WHICH FUNCTION WILL BE EVALUATED	GLENS070
NT=NO*NPO		GLENS071
IF(NFX.LE.O) ICODE	= 1	GLENS072
IF(NO.LE.O)ICODE	= 3	GLENS073
IF(NPQ.EQ.3)GO TO	1.5	GLENS074
IF(NPQ.NE.10)ICOD	1E = 2	GLENS075
5 IF(ICODE.GT.O)RET	URN	GLENS076
J=3		GLENS077
	0=L (	GLENS078
DELT= (B-A)/FLOAT	(NO)	GLENS079
J1 = J+1		GLENS080
NJ= J +NPQ		GLENS081
	BE ADDED TO FIRST POINT IN EACH QUADRATURE	GLENS082
UD(1:NPQ) = U(J1:		GLENS083
*****COMPUTE FIRST POI	NI IN EACH QUADKATURE	GLENS084
00 30 K=1.NO		GLENSO85
XI = K - 1	T	GLENSO86
30 FF(K) =A + XI*DEL		GLENSO87 Glenso88
	ALL POINTS FOR ALL QUADRATURES	GLENSO89
DO 40 K=1+NQ Kl =(K-1)*NPQ		GLENS089
		•
NK= K + NPQ UU{K1+1:NK} = FF{	K) + UD(1+NPO)	GLENSO91
40 CONTINUE	NI T UUTT+HLAI	GLENS092 GLENS093
C EVALUATE FUNCTION AT		GLENS094
C CARTONIC LONGITON MI	n.e , 0.1112	JUL 113074

#### APPENDIX B - Concluded

```
NOTE CHANGE IN CALLING SEQUENCE FOR FUNCTION EVALUATION SUBROUTINE
                                                                            GLENS095
0
      CALL FXS(UU,FOFX,MAXI,NT)
                                                                            GLENS096
                                                                            GLENS097
   CREATE VECTOR OF WEIGHTS
3
      00 50 K=1,NO
                                                                            GLENS098
                                                                            GLENS099
      K1 = (K-1) = NPQ
                                                                            GLENS100
      NK = K + NPQ
                                                                            GLENS101
      UU(K1+1:NK) = R(J1:NJ)
                                                                            GLENS102
   50 CONTINUE
   THESE STATEMENTS CAN BE REPLACED BY DOT PRODUCT FUNCTION IF AVAILABLEGLENSIO3
C
C
    SUCH AS
                                                                            GLENS104
5
5
5 **
      SUM(I) = DOTP (UULI:NT) ,FOFX(1:NT, I)
                                                                            GLENS105
                                                                            GLENS106
                                                                            GLENS107
      DO 80 I=1+NFX
                                                                            GLENS108
                                                                            GLENS109
      SUM(1)=0.0
      DO 70 K=1+NT
                                                                            GLENS110
   70 SUM(I)= SUM(I) +UU(K)* FOFX(K+I)
                                                                            GLENS111
      SUM(I) = DELT * SUM(I)
                                                                            GLENS112
   80 CONTINUE
                                                                            GLENS113
C. **
                                                                            GLENS114
      RETURN
                                                                            GLENS115
      END
                                                                            GLENS116
```

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