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Efficient Implementation of Minimal Polynomial and Reduced Rank Extrapolation Methods

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EFFICIENT IMPLEMENTATION OF MINIMAL POLYNOMIAL AND REDUCED RANK EXTRAPOLATION METHODS[†]

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ABSTRACT

The minimal polynomial extrapolation (MPE) and reduced rank extrapolation (RRE) are two very effective techniques that have been used in accelerating the convergence of vector sequences, such as those that are obtained from iterative solution of linear and nonlinear systems of equations. Their definitions involve some linear least squares problems, and this causes difficulties in their numerical implementation. In this work timewise efficient and numerically stable implementations for MPE and RRE are developed. A computer program written in FORTRAN 77 is also appended and applied to some model problems.

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1. INTRODUCTION

The minimal polynomial extrapolation (MPE) of Cabay and Jackson [2] and the reduced rank extrapolation (RRE) of Eddy [3] and Mešina [9] are two methods used in accelerating the convergence of a large class of vector sequences. In particular, they are employed for accelerating the convergence of fixed point iterative techniques for linear or nonlinear systems of equations, such as those that arise in the discrete solution of continuum problems.

A unified treatment of these and other extrapolation methods has been given in the survey paper of Smith, Ford, and Sidi [19], where some numerical testing for them is also provided. Detailed convergence analyses for MPE and RRE have been presented in Sidi [12], Sidi and Bridger [16], and Sidi [13], and we shall mention some of the results that follow from these analyses later in this work. Also, both MPE and RRE are very closely related to some well known Krylov subspace methods when they are applied to linearly generated vector sequences, and this subject is explored in detail in [13]. In fact, MPE and RRE are equivalent to the Arnoldi method and generalized conjugate residuals (GCR), respectively, when they are all applied to linear systems of equations starting with the same initial approximation. For the method of Arnoldi see Saad [10], and for GCR see Eisenstat, Elman, and Schultz [4]. We also mention that the conjugate gradient type method of Axelsson [1], the method of Young and Jea [22] that has been called ORTHODIR, and the recent generalized minimal residual method (GMRES) of Saad and Schultz [11] are all equivalent to GCR, and are used in solving linear equations. Recursion relations that exist amongst various approximations that are obtained from both methods are discussed in the paper by Ford and Sidi [6], where the existence of an interesting four-term lozenge recursion is shown. MPE and RRE have been employed successfully in Sidi and Celestina [17] in accelerating the convergence of some finite difference solution techniques in large scale computational fluid dynamics problems. Finally, the application of MPE and RRE and other vector extrapolation methods to the iterative solution of consistent singular linear systems has been considered in Sidi [15], where this approach is shown to be sound theoretically, and precise convergence analyses are also provided.

The definitions of MPE and RRE involve the solution of a linear least squares problem, the number of equations in this problem being equal to the dimension of the vectors in the given sequence. Since, in general, this dimension may be very large, as it is, for example, in three-dimensional computational fluid dynamics problems, the matrix of the least squares problem may be

very large. Thus, if standard linear least squares packages are used, the time and core memory requirements in the implementation of MPE and RRE may become prohibitive. To circumvent this problem, the solution of the linear least squares problem was achieved in [17] by solving the corresponding normal equations that is much less costly than using least squares packages. This approach proves to be quite efficient when the amount of extrapolation is not very large. When the amount of extrapolation is increased, however, the accuracy decreases, as the normal equations become very ill conditioned.

In the present work we propose new implementations for MPE and RRE, which are very inexpensive as far as both time and core memory requirements are concerned, and are stable numerically as the amount of extrapolation is increased. These implementations are also quite interesting mathematically, as they allow one to compute exactly (or estimate) the accuracy achieved in the extrapolation process without actually computing the residuals at each stage. This can be employed to further reduce the cost of implementation.

The plan of this paper is as follows: In Section 2 we briefly review the definitions of MPE and RRE. In Section 3 we consider the application of MPE and RRE to vector sequences that are generated by iterative solution of linear systems as this provides the motivation for different modes of usage of the methods. We devote Sections 4-6 to the development of the new implementations of MPE and RRE and the description of the mathematical features of these implementations. In Section 4 we give the details of the new implementations. One of the crucial ingredients of these implementations is the efficient solution of the least squares problems by use of QR factorization. In Section 5 we show how, in these new implementations, the l_2-norms of the residuals can be computed exactly for linear systems (or estimated for nonlinear systems) without doing extra vector computations. This enables us to assess the accuracy of the extrapolation without actually carrying it out, and can be used to reduce the amount of computation drastically. In Section 6 we discuss the operation counts and the storage requirements for the new implementations. In Section 7 we discuss some practical matters concerning the efficient use of MPE or RRE or any other vector extrapolation methods. Finally, in Section 8 we give some numerical results obtained by applying MPE and RRE through their new implementations to certain model problems. A computer program written in FORTRAN 77 that implements MPE and RRE is provided in the appendix.

2. REVIEW OF MPE AND RRE

Let $x_0, x_1, x_2, \ldots$, be a given sequence of N-dimensional column vectors, and denote its limit or antilimit by s. The vectors x_j are assumed to be complex, in general. Define

$$u_i = \Delta x_i = x_{i+1} - x_i \quad \text{and} \quad w_i = \Delta u_i = \Delta^2 x_i, \quad i = 0, 1, 2, \ldots . \tag{2.1}$$

Define the $N \times (j+1)$ matrices $U_j^{(n)}$ and $W_j^{(n)}$ by

$$U_j^{(n)} = [u_n \mid u_{n+1} \mid \cdots \mid u_{n+j}] \tag{2.2}$$

and

$$W_j^{(n)} = [w_n \mid w_{n+1} \mid \cdots \mid w_{n+j}] . \tag{2.3}$$

2.1 Definition of MPE

For MPE the approximation $s_{n,k}$ to s, the desired limit or antilimit, is defined by

$$s_{n,k} = \sum_{j=0}^{k} \gamma_j x_{n+j} , \tag{2.4}$$

where the γ_j are determined as follows:

(i) Use the least squares method to solve the overdetermined and, in general, inconsistent linear system

$$U_{k-1}^{(n)} c = -u_{n+k} , \tag{2.5}$$

where $c = (c_0, c_1, \ldots, c_{k-1})^T$.

(ii) Set $c_k = 1$, and compute the γ_j by

$$\gamma_j = \frac{c_j}{\sum_{i=0}^{k} c_i}, \quad 0 \le j \le k, \tag{2.6}$$

assuming that $\sum_{i=0}^{k} c_i \neq 0$. When this condition is not satisfied $s_{n,k}$ does not exist.

2.2 Definition of RRE

For RRE the approximation $s_{n,k}$ to s, the desired limit or antilimit, is defined by

$$s_{n,k} = x_n + \sum_{i=0}^{k-1} \xi_i u_{n+i} , \tag{2.7}$$

where the ξ_i are determined by solving the overdetermined and, in general, inconsistent linear system

$$W_{k-1}^{(n)}\xi = -u_n \,, \tag{2.8}$$

with $\xi = (\xi_0, \xi_1, \ldots, \xi_{k-1})^T$, using the least squares method. Since a least squares solution to (2.8) always exists, $s_{n,k}$ always exists. In particular, $s_{n,k}$ exists uniquely when the matrix $W_{k-1}^{(n)}$ has full rank, i.e., rank $(W_{k-1}^{(n)}) = k$, or, equivalently, when the vectors $w_n, w_{n+1}, \ldots, w_{n+k-1}$ are linearly independent. It can easily be shown that rank $(W_{k-1}^{(n)}) = k$, thus $s_{n,k}$ exists uniquely, when rank $(U_k^{(n)}) = k+1$.

There exists an equivalent formulation of RRE that seems to be more suitable for computer implementation. It also has the advantage of unifying most of the algorithmic aspects of MPE and RRE. In this formulation $s_{n,k}$ is of the form given in (2.4); only this time the γ_j are obtained by the least squares solution of the overdetermined and, in general, inconsistent linear system

$$U_k^{(n)}\gamma = 0 \,, \tag{2.9}$$

where $\gamma = (\gamma_0, \gamma_1, \ldots, \gamma_k)^T$, subject to the constraint

$$\sum_{j=0}^{k} \gamma_j = 1 \,. \tag{2.10}$$

(Note that the γ_j in MPE satisfy (2.10) automatically, as can easily be seen from (2.6).)

Remarks:

(1) It is important to realize that the γ_j in $s_{n,k} = \sum_{j=0}^{k} \gamma_j x_{n+j}$ depend on both n and k.

(2) In most applications, N, the dimension of the vectors x_i, is much larger than k, so that the matrices $U_j^{(n)}$ have many more rows than columns. Therefore, there is great need to reduce the amount of numerical work with the columns of the matrices $U_j^{(n)}$.

3. APPLICATION OF MPE AND RRE TO LINEAR SYSTEMS

Consider the linear nonsingular N-dimensional linear system

$$x = Ax + b \,, \tag{3.1}$$

where A is an $N \times N$ matrix and b is an N-dimensional column vector. Pick an initial vector x_0, and generate the vectors $x_1, x_2, \ldots$, by the iterative scheme

$$x_{i+1} = Ax_i + b, \quad i = 0, 1, \ldots \,. \tag{3.2}$$

The solution s of (3.1) is now the limit of the sequence $x_0, x_1, x_2, \ldots$, when the latter converges, otherwise, s is the antilimit.

Let k_0 be the degree of the minimal polynomial of the matrix A with respect to the vector $x_n - s$. Then the following statements are true.

(i) s_{n,k_0} is uniquely defined both for MPE and RRE, and

$$s_{n,k_0} = s \ . \tag{3.3}$$

Also the linear systems in (2.5), (2.8), and (2.9) are consistent for $k = k_0$, even though they may be overdetermined. This is a consequence of the fact that the vectors u_{n+j}, $0 \le j \le k_0 - 1$, are linearly independent, and u_{n+k_0} lies in their span. (See [13, Section 2.2].)

(ii) For $k < k_0$, $s_{n,k}$ is uniquely defined for RRE. For MPE, however, $s_{n,k}$ may fail to exist when $k < k_0$. When the matrix $C = I - A$ has positive definite hermitian part, $s_{n,k}$ exists uniquely for MPE also for $k < k_0$. (See [13, Section 2.2].) More generally, $s_{n,k}$ exists uniquely for MPE also for $k < k_0$, if the eigenvalues of C all lie on one side of a straight line through the origin in the complex plane, or, equivalently, if they all lie in an open sector $S = \{\mu: |\arg \mu - \theta| < \pi/2\}$, for some θ, $-\pi < \theta \le \pi$. This result can be proved exactly as Theorem 2.2 in [13] with C there replaced by $e^{-i\theta}C$.

(iii) When the Arnoldi method and GCR are used in solving the linear system $Cx = b$, where $C = I - A$, with x_n as the initial vector, they become equivalent to MPE and RRE, respectively. Specifically, the approximations obtained from the Arnoldi method and GCR are exactly $s_{n,1}, s_{n,2}, \ldots$, that are produced by MPE and RRE, respectively. (See [13, Section 2.3].)

(iv) If the *distinct nonzero* eigenvalues of A are denoted λ_j, $j = 1,2,\ldots$, and are ordered such that

$$|\lambda_1| \ge |\lambda_2| \ge |\lambda_3| \ge \cdots \ , \tag{3.4}$$

then, provided

$$|\lambda_k| > |\lambda_{k+1}| \ , \tag{3.5}$$

and A is diagonalizable, we have

$$s_{n,k} - s = O(|\lambda_{k+1}|^n) \quad \text{as } n \to \infty \ , \tag{3.6}$$

both for MPE and RRE. (The coefficient of $|\lambda_{k+1}|^n$ on the right hand side of (3.6) becomes large when the largest eigenvalues $\lambda_1, \lambda_2, \ldots$, are close to 1.) In view of the fact that

$x_n - s = O(|\lambda_1|^n)$ as $n \to \infty$, we conclude that MPE and RRE are both true acceleration methods. Under the same conditions, if $s_{n,k} - s$ is precisely $O(|\lambda_{k+1}|^n)$ as $n \to \infty$, then, the γ_j for MPE and RRE are such that

$$P^{n,k}(\lambda) = \sum_{j=0}^{k} \gamma_j \lambda^j = \prod_{i=1}^{k} \frac{\lambda - \lambda_i}{1 - \lambda_i} + O(|\lambda_{k+1}/\lambda_k|^n) \quad \text{as } n \to \infty, \tag{3.7}$$

i.e., for fixed k and for all sufficiently large n, the polynomial $P^{(n,k)}(\lambda)$ has precisely k zeros that tend to $\lambda_1, \lambda_2, \ldots, \lambda_k$. Furthermore, if we denote the zero of $P^{(n,k)}(\lambda)$ that tends to λ_j by $\lambda_j(n)$, then

$$\lambda_j(n) - \lambda_j = O(|\lambda_{k+1}/\lambda_j|^n) \quad \text{as } n \to \infty, \ 1 \le j \le k. \tag{3.8}$$

The proofs of (3.6) and (3.7) have been given in [12, Sections 3 and 4]. The proof of (3.8) will be published in the future. In case the matrix A in (3.2) is normal, the right hand sides of (3.7) and (3.8) can be replaced by $O(|\lambda_{k+1}/\lambda_k|^{2n})$ and $O(|\lambda_{k+1}/\lambda_j|^{2n})$, respectively. (The result in (3.6) remains the same, however.) This implies that when A is normal the rates of converge of $P^{(n,k)}(\lambda)$ and its zeros $\lambda_j(n)$ are twice those that can be achieved otherwise. These results follow from the corresponding results of Sidi [14].

For the most general case in which the matrix A is not diagonalizable, the results in (3.6)-(3.8) need to be modified considerably. For a complete treatment of this case see [16, Sections 2,3, and 5], where modifications of (3.6) and (3.7) are given. The modification of (3.8) will be published in the future.

A direct consequence of the result given in (3.6) is that better accuracy may be obtained if extrapolation is preceded by a number of fixed point iterations. This has indeed been observed numerically both for linear and nonlinear problems. We shall comment on this again in Section 7.

(v) Let us denote $C = I - A$. Then s is the solution to $Cx = b$. Denote by π_k the set of all polynomials $Q_k(\lambda)$ of degree at most k that satisfy $Q_k(0) = 1$. Consider now $s_{n,k}$ as obtained by applying MPE or RRE to the vector sequence $x_0, x_1, \ldots$. Then

$$\|r(s_{n,k})\| \le \left(\min_{Q_k \in \pi_k} \|Q_k(C)\| \right) \|r(x_n)\| \quad \text{for RRE,} \tag{3.9}$$

where $r(x) = Ax + b - x = b - Cx = -C(x - s)$ is the residual for x, and $\|\cdot\|$ is the l_2 vector norm,

or the matrix norm induced by it. (In fact, $\|r(s_{n,k})\|$, $k=0,1,2,\ldots$, is a monotonically decreasing sequence for RRE.) Similarly, if $C_h=\frac{1}{2}(C+C^*)$, the hermitian part of C, is positive definite, then

$$\|C_h^{1/2}(s_{n,k}-s)\| \le \beta(\min_{Q_k \in \pi_k} \|Q_k(C)\|)\|C_h^{1/2}(x_n-s)\| \quad \text{for MPE,} \tag{3.10}$$

where

$$\beta=\begin{cases} L & \text{if } C \text{ is normal} \\ L\sqrt{\text{cond}(C_h)} & \text{otherwise,} \end{cases} \tag{3.11}$$

with $L=\|C_h^{-1/2}C\,C_h^{-1/2}\| \ge 1$. Note that both $\|r(x)\|$ and $\|C_h^{1/2}x\|$ are true norms for x. Two types of bounds for $\min_{Q_k \in \pi_k} \|Q_k(C)\|$, in case C_h is positive definite, are given in [13, Section 4], and these can be used to derive upper bounds for $\|r(s_{n,k})\|$ and $\|C_h^{1/2}(s_{n,k}-s)\|$ for fixed n and increasing k. For details see [13]. These bounds are employed in [17] to justify the use of the extrapolation strategy that has been called "cycling" in [19] and all subsequent publications.

Finally, analogous and almost identical results exist for the case in which the system in (3.1) is singular but consistent, so that it has an infinity of solutions. In this case the limit or antilimit depends on x_0 in a very specific manner. For details, see [15].

Remark: The various Krylov subspace methods like the Arnoldi method and GCR and others can be applied only to linear systems. Acceleration methods such as MPE and RRE, however, can be applied to nonlinear systems as well as linear ones. The reason for this is that, unlike the Krylov subspace methods, MPE and RRE are defined exclusively in terms of the given vector sequence, which may be generated, for example, by an iterative method. Whether the vector sequence is generated linearly or nonlinearly is irrelevant to the definitions of MPE and RRE and other vector extrapolation methods. This is a very important property of vector extrapolation methods.

4. IMPLEMENTATION OF MPE AND RRE

4.1 General Considerations

As we have seen in Section 2, both MPE and RRE entail linear least squares problems in their definitions. There is, therefore, an immediate need for the efficient solution of these problems. We

propose to solve these problems by applying the QR factorization to the matrices $U_k^{(n)}$.

To keep the notation simple we shall set $n = 0$ everywhere, and denote the matrices $U_j^{(0)}$ by U_j. This amounts to simply renaming x_n and calling it x_0.

We assume that the vectors $u_0, u_1, \ldots, u_k$ are linearly independent so that the $N \times (k+1)$ matrix U_k is of full rank $k+1$. The case in which $u_0, u_1, \ldots, u_k$ are linearly dependent will be discussed later in this section. We recall that for the linear system in (3.1) this assumption is valid when $k < k_0$, where k_0 is the degree of the minimal polynomial of the matrix A with respect to the vector $x_0 - s$. Therefore, there is a unique $N \times (k+1)$ matrix Q_k,

$$Q_k = [q_0 \mid q_1 \mid \cdots \mid q_k] , \tag{4.1}$$

whose columns q_i satisfy

$$(q_i, q_j) = q_i^* q_j = \delta_{ij} , \tag{4.2}$$

and a unique $(k+1) \times (k+1)$ upper triangular matrix R_k,

$$R_k = \begin{bmatrix} r_{00} & r_{01} & r_{02} & \cdots & r_{0k} \\ & r_{11} & r_{12} & \cdots & r_{1k} \\ & & r_{22} & \cdots & r_{2k} \\ & O & & \ddots & \vdots \\ & & & & r_{kk} \end{bmatrix} , \tag{4.3}$$

with $r_{ii} > 0$, $i = 0, 1, \ldots, k$, such that

$$U_k = Q_k R_k . \tag{4.4}$$

This QR factorization amounts to orthonormalizing the vectors $u_0, u_1, u_2, \ldots$, in this order. It is important to retain this order, as this enables us to form the QR factorization of U_{k+1} by appending one additional column to Q_k to obtain Q_{k+1}, and a corresponding column to R_k to obtain R_{k+1}. Needless to say, this results in considerable savings in computing time.

QR factorization can be performed in different ways. The simplest way is the Gram-Schmidt (GS) process for orthonormalization of $u_0, u_1, u_2, \ldots$. This process is very unstable, however, in the sense that the computed vectors $q_0, q_1, q_2, \ldots$, are very far from being orthogonal. The modified Gram-Schmidt (MGS) process, on the other hand, seems to be quite stable, and is the one that we have

preferred. We recall that MGS is entirely equivalent to GS mathematically, and requires the same number of arithmetic operations as GS. The two methods are different numerically, however. For details, see, e.g., Golub and Van Loan [7, pp. 218-219].

For the sake of completeness we describe MGS for the case in which the vectors $u_0, u_1, u_2, \ldots$, are introduced one by one and in this order.

Algorithm MGS

Step 1. read u_0, and compute the scalar r_{00} and the vector q_0 according to
$r_{00} = (u_0, u_0)^{1/2}$ and $q_0 = u_0 / r_{00}$.

Step 2. for $k = 1,2,\ldots,$ do
read u_k, and set $u_k^{(0)} = u_k$
for $j=0$ to $k-1$ do

$r_{jk} = (q_j, u_k^{(j)})$

$u_k^{(j+1)} = u_k^{(j)} - r_{jk} q_j$
end
compute r_{kk} and q_k according to
$r_{kk} = (u_k^{(k)}, u_k^{(k)})^{1/2}$ and $q_k = u_k^{(k)} / r_{kk}$
end

(Here (y,z) stands for the Euclidean inner product $y^* z$, as before.)

It is easy to see that, when implementing MGS on a computer, $u_k^{(0)}, u_k^{(1)}, \ldots, u_k^{(k)}$, and q_k can all be made to occupy the same storage locations. As we shall see in the next paragraph, the computation of $s_{0,k}$ can be based on the q_j without the need to save either the x_j or the u_j. We can thus let u_k occupy the same storage locations as the $u_k^{(i)}$.

QR factorization can also be achieved by using Householder transformations. Although the computed matrices Q_k produced in this approach are closer to unitary than those produced by MGS when the l_2 condition number of U_k is large, the amount of computing in this approach is about twice that required by MGS. We shall elaborate on this further in Section 7.

We now recall from the definitions of MPE and RRE, that the approximations $s_{0,k}$ for both methods can be expressed in the form

$$s_{0,k} = \sum_{j=0}^{k} \gamma_j x_j \quad \text{with} \quad \sum_{j=0}^{k} \gamma_j = 1 \, . \tag{4.5}$$

Assuming that $\gamma_0, \gamma_1, \ldots, \gamma_k$ have been determined, let us compute $\xi_0, \xi_1, \ldots, \xi_{k-1}$ from

$$\xi_0 = 1-\gamma_0 \quad \text{and} \quad \xi_j = \xi_{j-1}-\gamma_j\ , \quad 1 \le j \le k-1\ . \tag{4.6}$$

Then, we can reexpress $s_{0,k}$ in the form

$$s_{0,k} = x_0 + \sum_{i=0}^{k-1} \xi_i u_i = x_0 + U_{k-1}\xi\ , \tag{4.7}$$

where $\xi = (\xi_0, \xi_1, \ldots, \xi_{k-1})^T$. Substituting now $U_{k-1} = Q_{k-1}\ R_{k-1}$ in (4.7), we obtain

$$s_{0,k} = x_0 + Q_{k-1}(R_{k-1}\xi) = x_0 + \sum_{j=0}^{k-1} \eta_j q_j\ , \tag{4.8}$$

where

$$\eta_j = (j+1)\text{st component of the column vector } R_{k-1}\xi\ , \quad j = 0,1,\ldots,k-1\ . \tag{4.9}$$

This approach to the computation of $s_{0,k}$ is very advantageous, as it enables us to overwrite $x_1, x_2, \ldots$, and $u_0, u_1, \ldots$, and thus saves a lot of storage.

4.2 Determination of the γ_j When rank $(U_k) = k+1$

The only thing that remains to be done now is to determine the γ_j, and this requires separate treatments for MPE and RRE.

4.2.1 Determination of the γ_j for MPE

As mentioned in Section 2, in order to determine the γ_j for $s_{0,k}$ in MPE we first solve the overdetermined system

$$U_{k-1}c = -u_k \tag{4.10}$$

by least squares. Since we also assume that the rank of U_k is $k+1$, we conclude that c is the unique solution of the normal equations

$$U_{k-1}^*\ U_{k-1}c = -U_{k-1}^*\ u_k\ . \tag{4.11}$$

Upon invoking $U_{k-1} = Q_{k-1}\ R_{k-1}$ in (4.11) and using the fact that $Q_{k-1}^*\ Q_{k-1} = I_{k\times k}$ = the $k\times k$ identity matrix, and the fact that R_{k-1} is a nonsingular matrix, we obtain

$$R_{k-1}c = -Q_{k-1}^*\ u_k\ . \tag{4.12}$$

It is easy to see that

$$Q_{k-1}^* u_k = (r_{0k}, r_{1k}, \ldots, r_{k-1,k})^T \equiv \rho_k\ , \tag{4.13}$$

so that (4.12) becomes

$$R_{k-1}c = -\rho_k \,. \tag{4.14}$$

This is a linear system of k equations in the k unknowns $c_0, c_1, \ldots, c_{k-1}$, and its matrix R_{k-1} is upper triangular. Hence its solution can be achieved easily by back substitution.

Once $c_0, c_1, \ldots, c_{k-1}$ are determined, we set $c_k = 1$, and compute the γ_j from (2.6), provided $\sum_{i=0}^{k} c_i \neq 0$.

4.2.2 Determination of the γ_j for RRE

Again as we mentioned in Section 2, the γ_j for $s_{0,k}$ in RRE can be determined by solving the overdetermined system

$$U_k \gamma = 0 \tag{4.15}$$

by least squares subject to the constraint

$$\sum_{j=0}^{k} \gamma_j = 1 \,. \tag{4.16}$$

This amounts to minimizing the positive definite quadratic form $\gamma^* U_k^* U_k \gamma$ subject to (4.16). Consequently, the lemma in Appendix A applies, and the γ_j can be obtained by solving the linear system of $k+2$ equations

$$\begin{aligned} U_k^* U_k \gamma &= \lambda \tilde{e} \\ \sum_{j=0}^{k} \gamma_j &= 1 \end{aligned} \tag{4.17}$$

for $\gamma_0, \gamma_1, \ldots, \gamma_k$, and λ. Here

$$\tilde{e} = (1, 1, \ldots, 1)^T \,. \tag{4.18}$$

As is stated in the same lemma, λ turns out to be strictly positive, and is given by

$$\lambda = \gamma^* U_k^* U_k \gamma \qquad \text{at the solution.} \tag{4.19}$$

The γ_j can be obtained by first solving the linear system

$$U_k^* U_k d = \tilde{e} \tag{4.20}$$

for $d = (d_0, d_1, \ldots, d_k)^T$, and letting

$$\lambda = \left[\sum_{j=0}^{k} d_j \right]^{-1} , \tag{4.21}$$

and finally setting

$$\gamma = \lambda d . \tag{4.22}$$

As far as the solution of the system in (4.20) is concerned, we accomplish this again by using the QR factorization of U_k. Again by $Q_k^* Q_k = I_{(k+1)\times(k+1)}$ = the $(k+1)\times(k+1)$ identity matrix, we can rewrite (4.20) in the form

$$R_k^* R_k d = \tilde{e} . \tag{4.23}$$

This system can be solved by forward and back substitution as the matrix R_k is upper triangular.

4.3 Treatment of the Case rank $(U_k) = k$

Up to this point we discussed the case in which the vectors $u_0, u_1, \ldots, u_k$ are linearly independent. Since these vectors are being introduced one by one, we can view this case as adding the vector u_k to the linearly independent set $\{u_0, u_1, \ldots, u_{k-1}\}$ and obtaining the linearly independent set $\{u_0, u_1, \ldots, u_k\}$. We now consider the case in which $\{u_0, u_1, \ldots, u_{k-1}\}$ is a linearly independent set, but $\{u_0, u_1, \ldots, u_k\}$ is not, i.e., rank $(U_k) = k$. This exhibits itself through $r_{kk} = 0$ in the QR factorization step.

If we are applying MPE, then we can compute the γ_j by solving the (nonsingular) system in (4.14) and employing (2.6), provided $\sum_{i=0}^{k} c_i \neq 0$ there. We then compute $s_{0,k}$.

If we are applying RRE, we can compute $s_{0,k}$ as follows: First, by the linear dependence of $u_0, u_1, \ldots, u_k$, there exist constants $\alpha_0, \alpha_1, \ldots, \alpha_k$, not all zero such that $\sum_{i=0}^{k} \alpha_i u_i = 0$. This implies that the linear system in (4.15) is consistent. Also we can write $U_k = Q_k R_k$, where Q_k and R_k are as in (4.1)-(4.3), $q_0, q_1, \ldots, q_{k-1}$ are uniquely determined and q_k is arbitrary in (4.1), and $r_{kk} = 0$ in (4.3). Multiplying both sides of (4.15) by Q_k^*, and using the fact that $Q_k^* Q_k = I_{(k+1)\times(k+1)}$, we obtain the system of $k+2$ equations

$$R_k \gamma = 0 \quad \text{and} \quad \sum_{j=0}^{k} \gamma_j = 1 . \tag{4.24}$$

Now, by $r_{kk} = 0$, this system actually consists of the $k+1$ inhomogeneous equations

$$[R_{k-1} \mid \rho_k]\gamma = 0 \quad \text{and} \quad \sum_{j=0}^{k} \gamma_j = 1 \tag{4.25}$$

in the $k+1$ unknowns $\gamma_0, \gamma_1, \ldots, \gamma_k$. Since a least squares solution for the linear system in (2.8) always exists, a solution for the γ_j always exists too. Consequently, the equations in (4.25) always have a solution for RRE. Once we determine a set of γ_j's, we compute $s_{0,k}$.

Comparing (4.25) with (4.14) and (2.6), we see that if $s_{0,k}$ exists for MPE when rank $(U_k) = k$, then it is equal to $s_{0,k}$ for RRE.

If the vector sequence $x_0, x_1, x_2, \ldots$, is generated as in (3.2), then, as explained in Section 3, rank $(U_k) = k+1$ for $k < k_0$, where k_0 is the degree of the minimal polynomial of A with respect to $x_0 - s$. The smallest value of k for which rank $(U_k) = k$ is k_0, and at $k = k_0$ we already reach the solution, i.e., $s_{0,k_0} = s$. That is the first time $r_{kk} = 0$ occurs, we have $s_{0,k} = s$, and stop.

If the vector sequence $x_0, x_1, x_2, \ldots$, is not generated linearly, and rank $(U_{k-1}) = k$, but rank$(U_k) = k < k+1$, then we can compute $s_{0,k}$ first, and then take $s_{0,k}$ or a nearby vector as x_0, and restart the computation. Other strategies for continuing the computation can likewise be devised, but we shall not pursue this matter further.

It should be mentioned, however, that, due to roundoff, the chances of encountering the case rank $(U_k) < k+1$ in practice are extremely small. We have thus not included the treatment of this case in the computer program given in the appendix.

4.4 Summary of Implementations

We now summarize the major steps of the implementations, as they have been described above. We assume that all the matrices U_k have full rank.

Suppose that, starting with x_0, we have constructed the matrices Q_{k-1} and R_{k-1}.

We now read x_{k+1} and compute $u_k = x_{k+1} - x_k$. Following this, using MGS, we compute the scalars $r_{0k}, r_{1k}, \ldots, r_{kk}$ and the orthonormal vector q_k, which we use to augment the matrices Q_{k-1} and R_{k-1} to give Q_k and R_k, respectively.

We next proceed to the computation of the γ_j. For MPE, we first solve the upper triangular $k \times k$ system in (4.14) for $c_0, c_1, \ldots, c_{k-1}$ by back substitution, and then use (2.6) to obtain the γ_j. For RRE, we solve the $(k+1) \times (k+1)$ system in (4.23) for d, and then determine the γ_j by (4.21) and (4.22). The solution of the system in (4.23) can be achieved very simply by forward and back substitution as R_k is

upper triangular.

Once the γ_j have been determined, we compute the ξ_j by (4.6) and the η_j by (4.9), and finally, $s_{0,k}$ by (4.8).

Next we read x_{k+2}, and proceed similarly, until a suitable stopping criterion is met.

It should be noted that, strictly speaking, neither r_{kk} nor q_k is needed for determining $s_{0,k}$, and their computation can be completed after x_{k+2} has been introduced. In the computer program that we give in the appendix, though, we chose to compute r_{kk} and q_k before the computation of $s_{0,k}$.

Finally, it is not difficult to see that these implementations are very appropriate for vector computers as their handling of the x_i, u_i, and q_i can be entirely vectorized. The computer program given in the appendix to this work has been written to take full account of this.

5. ESTIMATION OF RESIDUAL NORMS

5.1 General Considerations for Linear and Nonlinear Systems

Let s be the solution of the linear or nonlinear system of equations

$$x = F(x), \tag{5.1}$$

and let us define the residual for an arbitrary vector x by

$$r(x) = F(x) - x . \tag{5.2}$$

Let x_0 be a given initial approximation, and generate the sequence of vectors $x_1, x_2, \ldots$, according to the fixed point iterative method

$$x_{j+1} = F(x_j), \quad j = 0, 1, \ldots . \tag{5.3}$$

Consequently, the residual for x_j is given by

$$r(x_j) = F(x_j) - x_j = x_{j+1} - x_j = u_j , \tag{5.4}$$

thus is readily available.

Let us assume that MPE or RRE is applied to the sequence $x_0, x_1, x_2, \ldots$, and that we are computing the sequence $s_{n,1}, s_{n,2}, \ldots$. Let us assume also that we would like to stop the computation as soon as some norm of $r(s_{n,k})$ becomes $\leq \varepsilon$ for some k, $\varepsilon > 0$ being a preassigned level of accuracy. The most direct way of doing this would be by actually computing the vectors $s_{n,1}, r(s_{n,1}), s_{n,2}, r(s_{n,2}), \ldots$, which is very costly.

Indeed, the computation of $s_{n,k}$ involves about k vector additions and k scalar-vector multiplications, that of $r(s_{n,k})$, by (5.2), amounts to one additional fixed point iteration and one vector addition, and the computation of the norm of $r(s_{n,k})$ requires an additional inner product. In addition, the number of the vector operations increases with increasing k. In view of this, the most desirable situation is one that enables us to estimate some norm of $r(s_{n,k})$ without having to compute either $s_{n,k}$ or $r(s_{n,k})$.

5.2 Residual Computation for Linear Systems

We now devise a strategy by which the l_2-norms of the residuals $r(s_{n,k})$ can be obtained exactly without the need to compute either $s_{n,k}$ or $r(s_{n,k})$, when the sequence $x_0, x_1, x_2, \ldots$, is being generated linearly by the iterative method in (3.2), i.e., when $F(x) = Ax+b$ in (5.3). The case in which $F(x)$ is nonlinear will be considered at the end of this section.

When $F(x) = Ax+b$, the residual for an arbitrary vector x, by (5.2), becomes

$$r(x) = Ax + b - x \,. \tag{5.5}$$

Consequently, by (2.10), (3.2), and (2.1), we have

$$r(s_{0,k}) = \sum_{j=0}^{k} \gamma_j u_j = U_k \gamma \,, \tag{5.6}$$

and the l_2-norm of $r(s_{0,k})$ is thus

$$\| r(s_{0,k}) \| = (r(s_{0,k}), r(s_{0,k}))^{1/2} = (\gamma^* U_k^* U_k \gamma)^{1/2} \,. \tag{5.7}$$

By invoking $U_k = Q_k R_k$ in (5.7), we obtain

$$\| r(s_{0,k}) \| = (\gamma^* R_k^* R_k \gamma)^{1/2} \,. \tag{5.8}$$

We now analyze $\gamma^* R_k^* R_k \gamma$ for MPE and RRE separately.

5.2.1 l_2 -Norm of Residual with MPE

Let us compute $R_k \gamma$ first. By (4.12)-(4.14) we have

$$[R_{k-1} \mid \rho_k] \begin{bmatrix} c \\ 1 \end{bmatrix} = 0 \,. \tag{5.9}$$

By dividing both sides of (5.9) by $\sum_{i=0}^{k} c_i$ with $c_k = 1$, and invoking (2.6), we obtain

$$[R_{k-1} | \rho_k]\gamma = 0 . \tag{5.10}$$

Substituting (5.10) in $R_k\gamma$, we finally have

$$R_k\gamma = (0,0,...,0,r_{kk}\gamma_k)^T , \tag{5.11}$$

from which we obtain

$$(\gamma^* R_k^* R_k \gamma)^{1/2} = r_{kk} | \gamma_k | . \tag{5.12}$$

Consequently, for linearly generated sequences

$$\| r(s_{0,k}) \| = r_{kk} | \gamma_k | \tag{5.13}$$

exactly, with $r(x)$ as defined in (5.5).

5.2.2 l_2 -Norm of Residual with RRE

By (4.19) we have immediately

$$(\gamma^* U_k^* U_k \gamma)^{1/2} = \sqrt{\lambda} , \tag{5.14}$$

with λ as determined from (4.23) and (4.21). Consequently, for linearly generated sequences

$$\| r(s_{0,k}) \| = \sqrt{\lambda} \tag{5.15}$$

exactly, with $r(x)$ as defined in (5.5).

The results given in (5.13) and (5.15) assume exact arithmetic. Due to roundoff errors, however, the actually computed residual norms may be getting farther from (5.13) and (5.15), especially when k is increasing. In this case it may be appropriate to compute $s_{0,k}$ and the norm of its residual every once in a while to make sure that roundoff has not started to dominate the computations. Although such a test is not included in the computer program given in the appendix, it is quite easy to incorporate it there.

5.3 Practical Residual Estimation in Extrapolation for Nonlinear Systems

We now consider the problem of error estimation for the case in which $F(x)$ in (5.1) is nonlinear. Let us assume that the sequence $x_0, x_1, x_2, \ldots$, is convergent, its limit, of course, being s, the solution of (5.1). Therefore, for n sufficiently large, $x_n, x_{n+1}, \ldots$, are all very close to s, and we have

$$x_{n+1} - s = F'(s)(x_n - s) + \varepsilon_n , \tag{5.16}$$

where $F'(x)$ is the Jacobian matrix of the vector valued function $F(x)$, and ε_n is a vector whose norm

is $O(\|x_n - s\|^2)$ as $n \to \infty$. This implies that the sequence $x_0, x_1, x_2, \ldots$, behaves linearly at infinity, in the sense that

$$x_{j+1} \simeq F'(s)x_j + (s - F'(s)s) \tag{5.17}$$

for all sufficiently large j. Thus, for n sufficiently large, we can take

$$r(s_{n,k}) \simeq U_k^{(n)}\gamma \tag{5.18}$$

c.f., (5.6), and

$$\|r(s_{n,k})\| \simeq (\gamma^* R_k^{(n)*} R_k^{(n)} \gamma)^{1/2}, \tag{5.19}$$

c.f., (5.8), where we have retained the index n in $U_k^{(n)}$ and $U_k^{(n)} = Q_k^{(n)} R_k^{(n)}$. The norm in (5.19) is the l_2-norm as before. Consequently, we can take (5.19) as an estimate for the l_2-norm of the residual $r(s_{n,k})$ without having to compute either $s_{n,k}$ or $r(s_{n,k})$, since it is given by (5.12) for MPE and by (5.14) for RRE.

In case n is not large enough, (5.19) may not be very realistic. In this case we may choose to compute $s_{n,k}$ and $r(s_{n,k})$ not for all k, but for $k = p, 2p, 3p, \ldots$, say, for some integer $p > 1$. This obviously reduces the cost.

When we are using MPE or RRE in the cycling mode, which is one of the best modes of usage, things become simpler. To see this let us recall how cycling can be performed.

Step 1. Fix the integer k. Pick $s_k^0 \equiv x_0$ and set $q = 0$.

Step 2. Generate x_1, by (5.3). If $\|r(s_k^{(q)})\| = \|x_1 - x_0\| = \|u_0\| \leq \varepsilon$, then stop. Otherwise, generate $x_2, \ldots, x_{k+1}$ by (5.3).

Step 3. Compute $s_k^{(q+1)} \equiv s_{0,k}$ by MPE or RRE.

Step 4. Replace x_0 by $s_k^{(q+1)}$, and q by $q+1$, and go to Step 2.

(That $r(s_k^{(q)}) = u_0$ in Step 2 follows from (5.4).)

Consequently, no extra computation for residuals is necessary, as u_0 *is* the true residual in each cycle.

6. OPERATION COUNT AND STORAGE REQUIREMENTS

In most applications, N, the dimension of the vectors, is extremely large, while k takes on very small values. Consequently, the major part of the computational effort is spent in handling the large vectors, the rest being negligible.

As we can easily see, most of the vector computations take place in the QR factorization. At the kth stage that leads to $s_{0,k}$, the vector x_{k+1} is provided first. Starting with this, we need one vector addition to form $u_k = x_{k+1} - x_k$, and, following that, k vector additions, $k+1$ scalar-vector multiplications, and $k+1$ inner products to form the orthonormal vector q_k and the scalars $r_{0k}, r_{1k}, \ldots, r_{kk}$ by MGS. The computation of $s_{0,k}$, if desired, requires k vector additions and k scalar-vector multiplications by (4.8). The computation of the γ_i, ξ_i, and η_i is negligible, as it involves work with $k \times k$ or $(k+1) \times (k+1)$ triangular matrices for very small values of k.

As for the storage requirements, it is clear that x_0 needs to be saved. At the kth stage q_k needs to be saved, in addition to the previously saved $q_0, q_1, \ldots, q_{k-1}$. We also need two or three more auxiliary vectors of dimension N. Similarly the elements of the matrix R_k all need to be saved, but their storage requirements are negligible.

In view of the above, if only $s_{0,K}$ is needed for some preassigned K, then, recalling that the vector q_K need not be computed, the total operation count is $\frac{1}{2}(K^2+5K+2)$ vector additions, $\frac{1}{2}(K^2+5K)$ scalar-vector multiplications, and $\frac{1}{2}(K^2+3K+2)$ inner products, which amounts to $\sim 2K^2N$ floating point operations (scalar additions and multiplications). As for the storage requirements, we need $(K+1)N$ storage locations for $x_0, q_0, q_1, \ldots, q_{K-1}$, and $2N$ storage locations for two additional auxiliary vectors. No additional storage locations are required for $s_{0,k}$ as $s_{0,k}$ can overwrite x_0 at the end of the computation.

In many cases it turns out that the accuracy that can be achieved with m cycles of MPE or RRE, each cycle being of width K, is comparable to that obtained for $s_{0,mK}$. If we compare the computational costs of each of these strategies, we see that, roughly speaking, the former is m times less expensive computationally than the latter, and requires m times less storage. Thus, as a computational strategy, cycling possesses important advantages.

It is very instructive to compare the implementations for MPE and RRE, as they are given in this work, with the vector epsilon algorithm (VEA) of Wynn [21]. VEA is defined recursively by

$$\varepsilon_{-1}^{(n)} = 0 \quad \text{and} \quad \varepsilon_0^{(n)} = x_n, \quad n = 0,1,\ldots ,$$
$$\varepsilon_{k+1}^{(n)} = \varepsilon_{k-1}^{(n+1)} + \frac{\overline{\Delta\varepsilon_k^{(n)}}}{(\Delta\varepsilon_k^{(n)}, \Delta\varepsilon_k^{(n)})}, \quad k \geq 0, \ n \geq 0, \tag{6.1}$$

where $\Delta\varepsilon_k^{(n)} = \varepsilon_k^{(n+1)} - \varepsilon_k^{(n)}$, and $\bar{z} = (\bar{z}_1, \ldots, \bar{z}_N)^T$ if $z = (z_1, \ldots, z_N)^T$. Thus, the computation of $\varepsilon_k^{(n)}$ for $k \geq 2$ requires two vector additions, one scalar-vector multiplication, and one inner product. For $\varepsilon_1^{(n)}$ only one vector addition is required. Now as is suggested by experience and as can be justified heuristically, for given K, $\varepsilon_{2K}^{(0)}$ for VEA and $s_{0,K}$ for MPE or RRE would have comparable performance. The total operation count for determining $\varepsilon_{2K}^{(0)}$ is $4K^2$ vector additions, $2K^2+K$ scalar-vector multiplications, and $2K^2+K$ inner products, which amounts to $\sim 10K^2N$ floating points operations (scalar additions and multiplications). As for the storage requirements, we need $(2K+1)N$ storage locations to save $\varepsilon_0^{(2K)}, \varepsilon_1^{(2K-1)}, \ldots, \varepsilon_{2K}^{(0)}$, and $2N$ storage locations for two auxiliary vectors. Consequently, VEA is about five times more expensive than either MPE or RRE as far as operation counts are concerned. As far as storage requirements are concerned, VEA is about twice as expensive as either MPE or RRE. In addition, since $x_0, x_1, \ldots, x_{2K}$ are needed for $\varepsilon_{2K}^{(0)}$, whereas, only $x_0, x_1, \ldots, x_{K+1}$ are needed for either MPE or RRE, VEA is about twice as expensive as MPE or RRE with respect to the number of vectors they utilize.

We note that, in the epsilon family of vector extrapolation methods, VEA seems to be the most advantageous as far as the operation count, storage requirements, and numerical stability are concerned. For more details, see [19].

7. SOME PRACTICAL CONSIDERATIONS FOR ENHANCING CONVERGENCE AND STABILITY

In this section we would like to make a few remarks, which we believe are of practical importance with regard to enhancing the convergence and stability of vector extrapolation methods as they are applied to iterative procedures. Most of these remarks are based on the known theoretical results concerning vector extrapolation methods, some of which have been discussed in Section 3.

7.1 Effect of Iteration Before Extrapolation

In most problems of interest the vector sequence $x_0, x_1, \ldots$, converges extremely slowly so that there is not much difference between $\|x_n - s\|$ and $\|x_0 - s\|$ even for appreciably large values of n. The result in (3.6), however, suggests that there may be a large difference between $\|s_{n,k} - s\|$ and

$\|x_n - s\|$ (hence $\|x_0 - s\|$) if n is sufficiently large. If the vectors x_j are produced by an iterative procedure such as (3.2), then this implies that it may be very useful to start the extrapolation procedure after a number of iterations with (3.2). One heuristic argument in favor of this strategy runs as follows: The initial error $x_0 - s$, in general, has components in the direction of all eigenvectors and principal vectors of A. After a few iterations the components in the direction of those eigenvectors and principal vectors corresponding to zero eigenvalues of A are totally eliminated, while those corresponding to the eigenvalues that are close to zero are diminished. Consequently, the error vector $x_n - s$ has mostly contributions from the eigenvectors and principal vectors corresponding to the large eigenvalues. Precisely these contributions are now diminished by the extrapolation procedure.

7.2 A Simple "Averaging" of the Iteration Process and Its Effect on Convergence and Stability

Assume that (3.2) or (5.3) result from the discrete solutions of continuum problems. Then, for a convergent scheme, the largest eigenvalues of A or of $F'(s)$, the Jacobian matrix of $F(x)$ at $x = s$, may be very close to 1 in the complex plane in some cases. This may cause the extrapolation process not to be very effective. The process may even suffer from a large amount of numerical instability.

One way of dealing with this problem is by applying extrapolation methods not to the sequence $x_0, x_1, x_2, \ldots$, but to $y_0, y_1, y_2, \ldots$, where $y_j = x_{jp}$, for some positive integer p. This strategy has been successfully implemented in [17].

Another way would be by changing (5.3), in general, to read

$$x_{j+1} = x_j + \omega(F(x_j) - x_j) = (1-\omega)x_j + \omega F(x_j), \quad j = 0, 1, \ldots, \tag{7.1}$$

where ω is a scalar different than 1. (The sequence generated by taking $\omega = 1$ is the one generated by (5.3).) Thus x_{j+1} is now a weighted "average" of x_j and $F(x_j)$, in which the weights $1-\omega$ and ω need not be both positive.

By picking ω appropriately we can cause the spectrum of the Jacobian matrix of $(1-\omega)x + \omega F(x)$ at $x = s$, namely, $(1-\omega)I + \omega F'(s)$, to be increasingly favorable to $s_{n,k}$ for large values of n.

Let us take a look at the following example: Suppose the eigenvalues of $F'(s)$ are all positive and lie in the interval $[\varepsilon, 1-\eta]$ for some $\varepsilon > 0$ and $\eta > 0$ close to zero. Consequently, the sequence $x_0, x_1, \ldots$, obtained from (5.3) converges, provided x_0 is sufficiently close to s in case $F(x)$ is nonlinear, and unconditionally in case $F(x)$ is linear. If we pick $\omega = 2$, then the eigenvalues of

$(1-\omega)I + \omega F'(s)$ lie in the interval $[-1+2\varepsilon, 1-2\eta]$ so that the sequence obtained from (7.1) also converges. (If $\varepsilon = \eta$, then this sequence converges more quickly than the one obtained from (5.3).) The new spectrum has two important properties relevant to vector extrapolation methods: 1) The largest positive eigenvalue of $F'(s)$, namely, $1-\eta$, has moved away from 1. 2) Negative eigenvalues close to -1 have been created. Both of these properties enhance the stability of vector extrapolation processes both mathematically and numerically. (This follows from [12, Theorem 4.1], [16, Theorem 3.2], and [18, Theorems 4.1 and 5.2].) It should be noted that 2 is also that value of ω for which the spectral radius of $(1-\omega)I + \omega F'(s)$ is minimal when $\varepsilon = \eta$.

7.2.1 Special Considerations for Linear Systems

When $F(x) = Ax + b$, and the vector sequence is generated by the iterative procedure in (7.1), the approximations $s_{0,k}$ are independent of ω, as has been shown by Israeli and Sidi [8]. That is to say, the convergence properties of the $s_{0,k}$ are not changed by varying ω. Nevertheless, varying ω may influence the stability properties of the numerical implementations.

First, if the sequence obtained from (3.2) is divergent, then all the computations leading to $s_{0,k}$ will suffer a large loss of accuracy, especially for increasing k. By changing ω in (7.1) appropriately, we can cause the sequence to converge (or diverge very slowly), thus avoiding the numerical problem caused by the unboundedness of the original sequence.

Next, if the sequence obtained from (3.2) is slowly converging on account of the largest eigenvalues of A all being very close to 1 in the complex plane, then the vectors $u_0, u_1, u_2, \ldots$, are near being linearly dependent. Consequently, the l_2 condition number of the matrices U_k may be very large. This may have a negative influence on the QR factorization of U_k by MGS that we have chosen for our implementation. This influence exhibits itself in the computed matrices Q_k being far from unitary and the computed $s_{0,k}$ not being very accurate. If, by picking ω appropriately in (7.1), we can change the spectrum in such a way that it now contains both positive and negative large eigenvalues, then the vectors $u_0, u_1, u_2, \ldots$, will be far from being linearly dependent numerically. This will result in better conditioned matrices U_k, which, in turn, will result in the computed matrices Q_k being closer to unitary and the computed $s_{0,k}$ being quite accurate.

The numerical aspects of MGS and its use in the solution of least squares problems and the comparison of these with the Householder QR factorization and least squares solutions are discussed at

length in [7, Sections 5.2.8, 5.2.9, and 5.3.6].

7.2.2 Application to Jacobi Iteration for Consistently Ordered Matrices

The observations above can be used very effectively in the solution of linear systems whose matrices are consistently ordered. Such matrices arise frequently, for example, in the finite difference solutions of elliptic equations.

Suppose iterative methods of the form (3.2) are being used in the solution of such a system. If the method used is the Jacobi iteration method, then it is known that the nonzero eigenvalues of A come in pairs of the form $\pm\mu$, see, e.g., Varga [20, Chapter 4]. Consequently, if the eigenvalues of A are real, then they are in the interval $[-1+\delta, 1-\delta]$ for some δ, $0 < \delta < 1$, provided $\rho(A) < 1$. As a result, the nonzero eigenvalues of A^2 are in the interval $[\varepsilon, 1-\eta]$, for some $\varepsilon > 0$, where $1-\eta = (1-\delta)^2 \simeq 1-2\delta$ if $\delta \ll 1$. Furthermore, if $2M$ is the number of the distinct nonzero eigenvalues of A, then the number of the distinct nonzero eigenvalues of A^2 is M whether the eigenvalues of A are real or not.

This implies that the approximation $s^1_{2n,2k}$ obtained from the Jacobi iterative method and the approximation $s^2_{n,k}$ obtained from the double Jacobi iterative method

$$\begin{aligned} y &= Ax_j + b \\ x_{j+1} &= Ay + b\,, \qquad j = 0, 1, \ldots, \end{aligned} \tag{7.2}$$

have the same asymptotic behavior as $n \to \infty$. In addition, since the largest eigenvalues of A^2 are twice as far from 1 as those of A, $s^2_{n,k}$ is more stable than $s^1_{2n,2k}$ as $n \to \infty$ both mathematically and numerically.

We can now couple the double Jacobi iteration method with the simple averaging procedure that was discussed above. Tha new iteration procedure then is

$$\begin{aligned} y &= Ax_j + b \\ z &= Ay + b \\ x_{j+1} &= (1-\omega)x_j + \omega z\,, \qquad j = 0, 1, \ldots, \end{aligned} \tag{7.3}$$

for some $\omega \neq 0$. As explained before, by varying ω we can cause the spectrum of the iteration matrix of (7.3), namely, $(1-\omega)I + \omega A^2$, to become favorable to $s_{n,k}$. In particular, by picking $\omega = 2$ we can cause this spectrum to lie in the interval $[-1+2\varepsilon, 1-2\eta] = [-1+2\varepsilon, 1-4\delta+\delta^2]$. This enlarges the distance of the largest positive eigenvalue of the Jacobi iteration matrix A from 1 even further, and introduces

negative eigenvalues close to -1. This causes $s_{n,k}^2$ to become more stable. Furthermore, if $\epsilon \geq \delta$, the convergence rate of x_n from (7.3) with $\omega = 2$ is as good as that of x_n from (7.2).

We note, incidentally, that the iterative method of (7.3) with $\omega = 2$ is known as Abramov's method, see Faddeev and Faddeeva [5, p. 514]. It is quite easy to see that, in this case,

$$x_{j+1} - s = (2A^2 - I)(x_j - s) = T_2(A)(x_j - s),$$

where $T_2(\lambda) = 2\lambda^2 - 1$ is the Chebyshev polynomial of degree two. It should be emphasized that this is not Chebyshev acceleration, however.

8. NUMERICAL EXAMPLES

We have applied MPE and RRE through their new implementations described in the previous sections to several examples. This has been done by employing the computer program that is provided in Appendix B of this work. Some of the results obtained this way will be reported in this section.

We have picked real linear systems of equations whose matrices are symmetric or nonsymmetric. Numerical results for two of these systems, one symmetric and the other nonsymmetric, are included in this work.

Example 1. Consider the vector sequence obtained from (3.2), where A is a 1000×1000 septadiagonal matrix symmetric with respect to both of its main diagonals, and is given by

$$A = 0.06 \times \begin{bmatrix} 5 & 2 & 1 & 1 & & & & & \\ 2 & 6 & 3 & 1 & 1 & & & & \\ 1 & 3 & 6 & 3 & 1 & 1 & & & \\ 1 & 1 & 3 & 6 & 3 & 1 & 1 & & \\ & 1 & 1 & 3 & 6 & 3 & 1 & 1 & \\ & & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ & & & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{bmatrix}.$$

The vector b is such that the exact solution s of (3.1) is $(1,1,\ldots,1)^T$.

All eigenvalues of A are in $(0,1)$, the smallest and the largest being $4.7279\cdots \times 10^{-6}$ and $0.95999\cdots$, respectively. Consequently, the matrix $C = I - A$ is symmetric positive definite. Also, there is a large amount of clustering of eigenvalues near the smallest and the largest ones.

Taking $x_0 = 0$, we generated the vectors $x_1, x_2, \ldots$, by (7.1) once by taking $\omega = 1$ and once by taking $\omega = 2$, and then applied MPE to these two sequences. We also applied the method of conjugate gradients (CG) to the linear system $Cx = b$ starting again with $x_0 = 0$. The results of these computations are shown in Table 1a.

Recall that the Arnoldi method becomes equivalent to CG when C is a symmetric matrix, and MPE, when applied to a linearly generated sequence, becomes equivalent to the Arnoldi method. Also $s_{0,k}$, when applied to a sequence generated linearly as in (7.1), is independent of ω. Consequently, $s_{0,k}$, both for $\omega = 1$ and $\omega = 2$, obtained from MPE, and z_k, obtained from CG, are all the same mathematically. This is verified in Table 1a at least for $k \leq 10$. The differences between the $\omega = 1$ and $\omega = 2$ MPE computations for $k > 10$ can be explained exactly as described at the end of Section 7.2.1. Again, as can be seen from Table 1a, the $\omega = 2$ MPE computation differs from the CG computation starting with $k = 40$ approximately. Since CG involves orthogonalization with respect to only one vector, its absolute accuracy is guaranteed. On the other hand, MPE involves orthogonalization with respect to an ever increasing number of vectors at each stage, thus it cannot be absolutely accurate. In spite of this, the present implementation of MPE seems to be very stable in the sense that $\|s_{0,k} - s\|$ seems to be constantly decreasing with increasing k. Indeed, we have verified this by going up to $k = 100$ in both the $\omega = 1$ and $\omega = 2$ MPE computations.

Our purpose in presenting Table 1a was to demonstrate the good stability properties of the new MPE implementation for large values of k. Otherwise, CG is the method we would normally use for this example, since its operation count and storage requirements are extremely small.

In Table 1b we present the results obtained for the same example with $\omega = 2$ first performing 20 iterations and then using MPE in the cycling mode with $k = 10$, as explained at the end of Section 5. The remarkable effectiveness of this strategy is obvious.

	MPE						CG
	$\omega = 1$			$\omega = 2$			
k	r_{kk}/γ_k	$\|\|r(s_{0,k})\|\|$	$\|\|s_{0,k}-s\|\|$	r_{kk}/γ_k	$\|\|r(s_{0,k})\|\|$	$\|\|s_{0,k}-s\|\|$	$\|\|z_k-s\|\|$
0	1.46D+00	1.46D+00	3.16D+01	2.92D+00	2.92D+00	3.16D+01	3.16D+01
5	1.92D-01	1.92D-01	1.17D+00	3.83D-01	3.83D-01	1.17D+00	1.17D+00
10	1.98D-02	1.98D-02	1.53D-01	3.96D-02	3.96D-02	1.53D-01	1.53D-01
15	2.51D-03	2.51D-03	2.03D-02	5.01D-03	5.01D-03	2.02D-02	2.02D-02
20	4.51D-04	4.53D-04	3.70D-03	6.63D-04	6.63D-04	2.68D-03	2.68D-03
25	1.58D-04	2.26D-04	4.43D-03	8.78D-05	8.78D-05	3.52D-04	3.52D-04
30	6.27D-05	1.11D-04	2.44D-03	1.15D-05	1.15D-05	4.63D-05	4.63D-05
35	2.49D-05	3.19D-05	6.08D-04	1.53D-06	1.53D-06	6.53D-06	6.11D-06
40	6.37D-06	1.42D-05	3.34D-04	5.16D-07	5.30D-07	1.64D-06	8.03D-07
45	7.93D-06	5.42D-06	3.66D-05	7.31D-08	1.29D-07	1.27D-06	1.06D-07
50	9.19D-07	5.29D-06	1.30D-04	3.17D-08	4.29D-08	1.85D-07	1.39D-08

Table 1a - Numerical results for Example 1, starting with $x_0 = 0$.

i	$\|\|r(s_k^{(i)})\|\|$	$\|\|s_k^{(i)}-s\|\|$
0	4.75 D-01	5.91 D+00
1	2.00 D-04	6.94 D-04
2	2.90 D-06	8.78 D-06
3	4.17 D-08	1.74 D-07
4	9.27 D-10	3.70 D-09
5	2.18 D-11	9.11 D-11
6	5.49 D-13	2.83 D-12
7	4.26 D-14	1.77 D-13
8	6.16 D-15	9.46 D-14

Table 1b - MPE applied to Example 1 in the cycling mode. Starting with the zero vector, first 20 iterations are performed. Following that MPE is applied in the cycling mode with $k = 10$. The l_2-norm of the error in the initial (zero) vector is 3.16D+01. The vectors are obtained by "averaging" the iterative process (3.2) with $\omega = 2$.

Example 2. Consider the linear nonsymmetric system of equations $\tilde{C}x = \tilde{b}$, where $\tilde{C}$ is the block - tri-diagonal matrix

$$\tilde{C} = \begin{bmatrix} B & -I & & & \\ -I & B & -I & & \\ & -I & B & -I & \\ & & \ddots & \ddots & \ddots \\ & & & -I & B \end{bmatrix} \quad \text{with } B = \begin{bmatrix} 4 & a & & \\ b & \ddots & \ddots & \\ & \ddots & \ddots & a \\ & & b & 4 \end{bmatrix},$$

and $a=-1+\delta$, $b=-1-\delta$, $\delta \neq 0$. (See [10, p. 122].) Again, the vector $\tilde{b}$ is such that the exact solution s is $(1,1,...,1)^T$. The iterative method that we pick for this system is Jacobi's method, so that $A = I - \frac{1}{4}\tilde{C}$.

Now the matrix $\tilde{C}$ is consistently ordered. Thus the suggestions put forth in Section 7.2.2 can be successfully employed in this case.

In our numerical experiments we took $\delta = 0.2$. The matrices B and I in $\tilde{C}$ were all 10×10 and $\tilde{C}$ was 200×200, exactly as in [10]. The extrapolation method for which we give numerical results is RRE. We first applied RRE in the cycling mode in conjunction with Jacobi iteration. The vector obtained at the end of each cycle is denoted $\bar{s}_k^{(i)}$. Next we applied RRE in the cycling mode in conjunction with double Jacobi iteration. The vector obtained at the end of each cycle is denoted $\hat{s}_k^{(i)}$ now. Finally, we applied RRE in the cycling mode extended as follows: The vector sequence is generated by the iterative procedure of (7.3) with $\omega = 2$, i.e., by the "averaged" double Jacobi iteration with $\omega = 2$. In each cycle n_i+k_i+1 such iterations are performed, and extrapolation is applied to the last k_i+2 of the vectors, i.e., in each cycle s_{n_i,k_i} is computed. The vector obtained at the end of each cycle now is denoted $\tilde{s}_{n_i,k_i}^{(i)}$. The index i denotes the cycle number in each case.

In Table 2 we give the l_2 - norms of the errors $\bar{s}_k^{(i)} - s$ (k = 20), $\hat{s}_k^{(i)} - s$ (k = 10), and $\tilde{s}_{n_i,k_i}^{(i)}$ ($n_i = 5$, $k_i = 5$ all i). Thus the number of basic Jacobi iterations performed to obtain the approximations $\bar{s}_k^{(i)}$, $\hat{s}_k^{(i)}$, and $\tilde{s}_{n_i,k_i}^{(i)}$ in each cycle is 21, 22, and 22 respectively. We see that $\bar{s}_{20}^{(i)}$ and $\hat{s}_{10}^{(i)}$ have comparable accuracy, as expected. The number of vector operations for $\bar{s}_{20}^{(i)}$, however, is over three

times that for $\hat{s}_{10}^{(i)}$. Also the storage requirement for $\bar{s}_{20}^{(i)}$ is about twice that for $\hat{s}_{10}^{(i)}$. The performance of $\tilde{s}_{5,5}^{(i)}$ is only slightly inferior. The number of vector operations for $\tilde{s}_{5,5}^{(i)}$ is about one tenth that for $\bar{s}_{20}^{(i)}$, while its storage needs are about one third those of $\bar{s}_{20}^{(i)}$.

i	$\|\|\bar{s}_{20}^{(i)}-s\|\|$	$\|\|\hat{s}_{10}^{(i)}-s\|\|$	$\|\|\tilde{s}_{5,5}^{(i)}-s\|\|$
1	6.66D-02	7.47D-02	1.34D-01
2	2.02D-04	2.36D-04	5.86D-04
3	2.53D-07	4.26D-07	1.14D-05
4	2.90D-10	2.05D-09	3.04D-08
5	2.03D-12	5.96D-12	2.15D-10
6	1.35D-13	6.48D-14	1.07D-12
7	3.61D-14	3.13D-14	1.75D-14

Table 2 - RRE applied to Example 2 in the cycling mode. The initial vector is zero, and the l_2-norm of the error associated with it is 1.41D+01.

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APPENDIX A

Lemma: Let T be an $m \times m$ hermitian positive definite matrix, and let $z_1, z_2, \ldots, z_m$ be complex variables. Denote $z = (z_1, z_2, \ldots, z_m)^T$. Then the solution to the problem

$$\begin{aligned} &\text{minimize } z^* T z \\ &\text{subject to } \sum_{i=1}^{m} z_i = 1 \end{aligned} \tag{A.1}$$

can be obtained by solving the linear system of $m+1$ equations

$$\begin{aligned} T z &= \lambda \tilde{e} \\ \sum_{i=1}^{m} z_i &= 1, \end{aligned} \tag{A.2}$$

where $z_1, \ldots, z_m$, and λ are unknowns, and

$$\tilde{e} = (1, 1, \ldots, 1)^T . \tag{A.3}$$

The unknown λ turns out to be real and positive, and is given by

$$\lambda = z^* T z \quad \text{at the solution.} \tag{A.4}$$

The solution of (A.2) can be achieved by first solving the system

$$T h = \tilde{e} \tag{A.5}$$

for $h = (h_1, \ldots, h_m)^T$, and letting

$$\lambda = \left(\sum_{i=1}^{m} h_i \right)^{-1} , \tag{A.6}$$

and finally setting

$$z = \lambda h . \tag{A.7}$$

Proof: We start by expressing the problem in terms of real variables. Let us write T in the form

$$T = M + iN, \quad M \text{ and } N \text{ real } m \times m \text{ matrices} . \tag{A.8}$$

Then, by the assumption that T is hermitian, it follows that

$$M^T = M \quad \text{and} \quad N^T = -N . \tag{A.9}$$

Writing

$$z = x + iy, \quad x \text{ and } y \text{ real } m\text{-dimensional vectors} , \tag{A.10}$$

and invoking (A.8) and (A.9) in $z^* T\, z$, we have

$$z^* T\, z = x^T M\, x + y^T M\, y + 2y^T N\, x\,. \tag{A.11}$$

Now the constraint $\sum_{i=1}^{m} z_i = 1$ in (A.1) is equivalent to two real constraints, namely,

$$\sum_{i=1}^{m} x_i = 1 \quad \text{and} \quad \sum_{i=1}^{m} y_i = 0\,. \tag{A.12}$$

We now use the method of Lagrange multipliers to minimize (A.11) subject to (A.12). Introducing the Lagrange multipliers -2μ and -2ν for the constraints $\sum_{i=1}^{m} x_i = 1$ and $\sum_{i=1}^{m} y_i = 0$, respectively, and taking derivatives with respect to the x_i and y_i, we obtain the linear system of equations

$$\begin{aligned} Mx - Ny - \mu\, \tilde{e} &= 0 \\ My + Nx - \nu\, \tilde{e} &= 0, \end{aligned} \tag{A.13}$$

which, upon letting $\lambda = \mu + i\nu$, becomes equivalent to $Tz = \lambda\, \tilde{e}$. We have thus shown the truth of (A.2). Multiplying $Tz = \lambda\, \tilde{e}$ on the left by z^*, and using $\sum_{i=1}^{m} z_i = 1$, we obtain (A.4). Obviously, λ has to be strictly positive. For if λ were zero, then $z=0$ would have to be the solution as T is hermitian positive definite, but this would contradict the constraint $\sum_{i=1}^{m} z_i = 1$. The rest of the proof follows easily from (A.2), and we shall omit it. □

APPENDIX B

In this appendix we give a computer code written in standard FORTRAN 77 that implement MPE and RRE as described in the present work.

The implementation of MPE and RRE is done in SUBROUTINE MPERRE that forms the heart of this code.

Use of MPE and RRE in the cycling mode is made possible by SUBROUTINE CYCLE.

The vector sequence for extrapolation is generated by calling SUBROUTINE VECTOR, which, in the present code provides the iteration sequence of Example 1 with $\omega = 2$ weighting.

The driving program in the present code is the one that generates some of results shown in Table 1b.

We give no further explanations about the code and its use, as the different parts of the code are documented in detail.

```
C*************************************************************************CYC00010
C  IMPLEMENTATION OF MPE AND RRE WITH QR FACTORIZATION FOR LEAST          CYC00020
C  SQUARES. (QR PERFORMED BY MODIFIED GRAM-SCHMIDT PROCESS)               CYC00030
C  MPE AND RRE ARE APPLIED IN THE CYCLING MODE.                           CYC00040
C*************************************************************************CYC00050
C  THE COMPONENTS OF THE INITIAL VECTOR X, NAMELY, X(I),I=1,...,NDIM,     CYC00060
C  CAN BE PICKED RANDOMLY. WE ACHIEVE THIS, E.G., BY INVOKING THE         CYC00070
C  IMSL VERSION 10 SUBROUTINE DRNUN THAT GENERATES PSEUDORANDOM           CYC00080
C  NUMBERS FROM A UNIFORM (0,1) DISTRIBUTION.                             CYC00090
C  OTHER CHOICES FOR X(1),...,X(NDIM) ARE POSSIBLE, SUCH AS X(I)=0,       CYC00100
C  I=1,...,NDIM. IN THIS CASE REPLACE THE STATEMENT                       CYC00110
C     CALL DRNUN(NDIM,X)                                                  CYC00120
C  BY THE DO LOOP                                                         CYC00130
C     DO 10 I=1,NDIM                                                      CYC00140
C     X(I)=0                                                              CYC00150
C 10  CONTINUE                                                            CYC00160
C*************************************************************************CYC00170
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)                                 CYC00180
      PARAMETER (METHOD=1,N0=20,N=0,KMAX=10,NCYCLE=15,NDIM=1000)          CYC00190
      PARAMETER (EPSC=1D-10,IPRES=-1,IPRES1=1)                            CYC00200
      DIMENSION X(NDIM),S(NDIM),Y(NDIM),Z(NDIM)                           CYC00210
      DIMENSION Q(NDIM,0:KMAX-1),R(0:KMAX,0:KMAX)                         CYC00220
      DIMENSION C(0:KMAX),GAMMA(0:KMAX),XI(0:KMAX-1)                      CYC00230
      EXTERNAL VECTOR                                                     CYC00240
C                                                                         CYC00250
C  INITIAL VECTOR DETERMINATION.                                          CYC00260
C                                                                         CYC00270
C     CALL DRNUN(NDIM,X)                                                  CYC00280
      DO 10 I=1,NDIM                                                      CYC00290
      X(I)=0                                                              CYC00300
  10  CONTINUE                                                            CYC00310
C                                                                         CYC00320
C  END OF INITIAL VECTOR DETERMINATION.                                   CYC00330
C                                                                         CYC00340
      CALL CYCLE(METHOD,X,S,N0,N,KMAX,NCYCLE,NDIM,Y,Z,VECTOR,Q,R,         CYC00350
     *C,GAMMA,XI,RESC,EPSC,IPRES,IPRES1)                                  CYC00360
      STOP                                                                CYC00370
      END                                                                 CYC00380
                                                                          CYC00390
      SUBROUTINE CYCLE(METHOD,X,S,N0,N,KMAX,NCYCLE,NDIM,Y,Z,VECTOR,Q,R,   CYC00400
     *C,GAMMA,XI,RESC,EPSC,IPRES,IPRES1)                                  CYC00410
C*************************************************************************CYC00420
C  THIS SUBROUTINE APPLIES MPE AND RRE IN THE CYCLING MODE.               CYC00430
C  MPE AND RRE ARE INVOKED BY CALLING SUBROUTINE MPERRE.                  CYC00440
C*************************************************************************CYC00450
C  THE ARGUMENTS METHOD,NDIM,Y,Z,VECTOR,Q,R,C,GAMMA,XI,IPRES,IPRES1       CYC00460
C  ARE AS IN SUBROUTINE MPERRE.                                           CYC00470
C                                                                         CYC00480
C  X      : INITIAL VECTOR. INPUT ARRAY OF DIMENSION NDIM. (DOUBLE        CYC00490
C           PRECISION)                                                    CYC00500
C  S      : THE FINAL APPROXIMATION PRODUCED BY THE SUBROUTINE. OUTPUT    CYC00510
C           ARRAY OF DIMENSION NDIM. (DOUBLE PRECISION)                   CYC00520
C  N0     : NUMBER OF ITERATIONS PERFORMED BEFORE CYCLING IS STARTED,     CYC00530
C           I.E., BEFORE MPE OR RRE IS APPLIED FOR THE FIRST TIME.        CYC00540
C           INPUT. (INTEGER)                                              CYC00550
C  N      : NUMBER OF ITERATIONS PERFORMED BEFORE MPE OR RRE IS APPLIED    CYC00560
C           IN EACH CYCLE AFTER THE FIRST CYCLE. INPUT. (INTEGER)         CYC00570
C  KMAX   : WIDTH OF EXTRAPOLATION. ON EXIT FROM SUBROUTINE MPERRE IN     CYC00580
C           EACH CYCLE, THE ARRAY S IS, IN FACT, THE APPROXIMATION        CYC00590
C           S(N0,KMAX) IN THE FIRST CYCLE,AND S(N,KMAX) IN THE FOLLOWING CYC00600
C           CYCLES. INPUT. (INTEGER)                                      CYC00610
C  NCYCLE: MAXIMUM NUMBER OF CYCLES ALLOWED. INPUT. (INTEGER)             CYC00620
C  RESC   : L2-NORM OF THE RESIDUAL FOR S AT THE END OF EACH CYCLE.       CYC00630
C           RETRIEVED AT THE END OF THE NEXT CYCLE. OUTPUT. (DOUBLE       CYC00640
C           PRECISION)                                                    CYC00650
C  EPSC   : AN UPPER BOUND ON RESC/RESP, SOME RELATIVE RESIDUAL FOR S,    CYC00660
```

```
C           USED IN THE STOPPING CRITERION. HERE RESP IS THE L2-NORM     CYC00670
C           OF THE RESIDUAL FOR S(N0,KMAX) AT THE END OF THE FIRST       CYC00680
C           CYCLE, I.E., ON EXIT FROM SUBROUTINE MPERRE THE FIRST TIME.  CYC00690
C           IF RESC.LE.EPSC*RESP AT THE END OF SOME CYCLE, THEN ONE      CYC00700
C           ADDITIONAL CYCLE IS PERFORMED, AND THE CORRESPONDING         CYC00710
C           S(N,KMAX) IS ACCEPTED AS THE FINAL APPROXIMATION, AND THE    CYC00720
C           SUBROUTINE IS EXITED. INPUT. (DOUBLE PRECISION)              CYC00730
C*********************************************************************CYC00740
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)                                CYC00750
      PARAMETER (EPS=0)                                                  CYC00760
      DIMENSION X(NDIM),S(NDIM),Y(NDIM),Z(NDIM)                          CYC00770
      DIMENSION Q(NDIM,0:KMAX-1),R(0:KMAX,0:KMAX)                        CYC00780
      DIMENSION C(0:KMAX),GAMMA(0:KMAX),XI(0:KMAX-1)                     CYC00790
      EXTERNAL VECTOR                                                    CYC00800
      DO 40 IC=1,NCYCLE                                                  CYC00810
      IF (IPRES.EQ.1.OR.IPRES1.EQ.1) THEN                                CYC00820
      WRITE(6,101) IC                                                    CYC00830
 101  FORMAT(/,'  CYCLE NO. ',I3)                                        CYC00840
      END IF                                                             CYC00850
      NN=N                                                               CYC00860
      IF (IC.EQ.1) NN=N0                                                 CYC00870
      IF (IPRES.EQ.1.OR.IPRES1.EQ.1) THEN                                CYC00880
      WRITE(6,102) NN                                                    CYC00890
 102  FORMAT(/,'  NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS',I3)       CYC00900
      WRITE(6,103) KMAX                                                  CYC00910
 103  FORMAT(/,'  WIDTH OF EXTRAPOLATION IS ',I3)                        CYC00920
      END IF                                                             CYC00930
      DO 20 J=0,NN-1                                                     CYC00940
      CALL VECTOR(X,Y,NDIM)                                              CYC00950
      DO 10 I=1,NDIM                                                     CYC00960
      X(I)=Y(I)                                                          CYC00970
  10  CONTINUE                                                           CYC00980
  20  CONTINUE                                                           CYC00990
      CALL MPERRE(METHOD,X,S,KMAX,KOUT,NDIM,Y,Z,VECTOR,Q,R,C,            CYC01000
     *GAMMA,XI,RES,RES1,EPS,IPRES,IPRES1)                                CYC01010
      IF (IC.EQ.1) RESP=R(0,0)                                           CYC01020
      RESC=R(0,0)                                                        CYC01030
      IF (RESC.LE.EPSC*RESP) RETURN                                      CYC01040
      DO 30 I=1,NDIM                                                     CYC01050
      X(I)=S(I)                                                          CYC01060
  30  CONTINUE                                                           CYC01070
  40  CONTINUE                                                           CYC01080
      RETURN                                                             CYC01090
      END                                                                CYC01100
                                                                         CYC01110
      SUBROUTINE MPERRE(METHOD,X,S,KMAX,KOUT,NDIM,Y,Z,VECTOR,Q,R,C,      CYC01120
     *GAMMA,XI,RES,RES1,EPS,IPRES,IPRES1)                                CYC01130
C*********************************************************************CYC01140
C  THIS SUBROUTINE APPLIES THE MINIMAL POLYNOMIAL EXTRAPOLATION (MPE)   CYC01150
C  OR THE REDUCED RANK EXTRAPOLATION  (RRE) METHODS TO A VECTOR         CYC01160
C  SEQUENCE X0,X1,X2,..., THAT IS OFTEN GENERATED BY A FIXED POINT      CYC01170
C  ITERATIVE TECHNIQUE.                                                 CYC01180
C  BOTH MPE AND RRE ARE ACCELERATION OF CONVERGENCE (OR EXTRAPOLATION)  CYC01190
C  METHODS FOR VECTOR SEQUENCES. EACH METHOD PRODUCES A TWO-DIMENSIONAL CYC01200
C  ARRAY S(N,K) OF APPROXIMATIONS TO THE LIMIT OR ANTILIMIT OF THE      CYC01210
C  SEQUENCE IN QUESTION.                                                CYC01220
C  THE IMPLEMENTATIONS EMPLOYED IN THE PRESENT SUBROUTINE GENERATE      CYC01230
C  THE SEQUENCES S(0,0)=X0,S(0,1),S(0,2),... .                          CYC01240
C*********************************************************************CYC01250
C  AUTHOR        :  AVRAM SIDI                                          CYC01260
C                   COMPUTER SCIENCE DEPARTMENT                         CYC01270
C                   TECHNION-ISRAEL INSTITUTE OF TECHNOLOGY             CYC01280
C                   HAIFA 32000,ISRAEL                                  CYC01290
C  E-MAIL ADDRESS:  CSSSIDI@TECHNION.BITNET                             CYC01300
C*********************************************************************CYC01310
C  METHOD: IF METHOD.EQ.1, THEN MPE IS EMPLOYED. IF METHOD.EQ.2, THEN   CYC01320
```

```
C           RRE IS EMPLOYED. INPUT.. (INTEGER)                                  CYC01330
C  X      : THE VECTOR X0. INPUT ARRAY OF DIMENSION NDIM. (DOUBLE              CYC01340
C           PRECISION)                                                          CYC01350
C  S      : THE APPROXIMATION S(0,K) PRODUCED BY THE SUBROUTINE FOR            CYC01360
C           EACH K. ON EXIT, S IS S(0,KOUT). OUTPUT ARRAY OF DIMENSION          CYC01370
C           NDIM. (DOUBLE PRECISION)                                            CYC01380
C  KMAX   : A NONNEGATIVE INTEGER.THE MAXIMUM WIDTH OF EXTRAPOLATION           CYC01390
C           ALLOWED. THUS THE NUMBER OF THE VECTORS X0,X1,X2,...,               CYC01400
C           EMPLOYED IN THE PROCESS IS KMAX+2 AT MOST. INPUT. (INTEGER)         CYC01410
C  KOUT   : A NONNEGATIVE INTEGER. KOUT IS DETERMINED BY A SUITABLE            CYC01420
C           STOPPING CRITERION, AND DOES NOT EXCEED KMAX.THE VECTORS            CYC01430
C           ACTUALLY EMPLOYED BY THE EXTRAPOLATION PROCESS ARE                  CYC01440
C           X0,X1,X2,...,XP, WHERE P=KOUT+1. OUTPUT. (INTEGER)                  CYC01450
C  NDIM   : DIMENSION OF THE VECTORS. INPUT. (INTEGER)                         CYC01460
C  Y      : WORK ARRAY OF DIMENSION NDIM. (DOUBLE PRECISION)                   CYC01470
C  Z      : WORK ARRAY OF DIMENSION NDIM. (DOUBLE PRECISION)                   CYC01480
C  VECTOR: A USER-SUPPLIED  SUBROUTINE WHOSE CALLING SEQUENCE IS               CYC01490
C              CALL VECTOR(Y,Z,NDIM); Y,NDIM INPUT,Z OUTPUT.                   CYC01500
C           Y,Z,NDIM ARE EXACTLY AS DESCRIBED ABOVE.FOR A FIXED POINT           CYC01510
C           ITERATIVE TECHNIQUE FOR SOLVING THE LINEAR OR NONLINEAR             CYC01520
C           SYSTEM T=F(T), DIM(T)=NDIM, Y AND Z ARE RELATED BY Z=F(Y).          CYC01530
C           THUS X1=F(X0), X2=F(X1), ETC.                                       CYC01540
C           VECTOR SHOULD BE DECLARED IN AN EXTERNAL STATEMENT IN THE           CYC01550
C           CALLING PROGRAM.                                                    CYC01560
C  Q      : WORK ARRAY OF DIMENSION (NDIM,0:KMAX-1). FOR EACH K, ITS           CYC01570
C           ELEMENTS ARE THOSE OF THE ORTHOGONAL MATRIX OBTAINED  FROM          CYC01580
C           QR FACTORIZATION OF THE MATRIX U                                    CYC01590
C                     U = ( U0 | U1 | ... | UK ), K=0,1,2,...,                  CYC01600
C           WHERE U0=X1-X0, U1=X2-X1, U2=X3-X2, ETC. OUTPUT. (DOUBLE            CYC01610
C           PRECISION)                                                          CYC01620
C  R      : WORK ARRAY OF DIMENSION (0:KMAX,0:KMAX). FOR EACH K, ITS           CYC01630
C           ELEMENTS ARE THOSE OF THE UPPER TRIANGULAR MATRIX OBTAINED          CYC01640
C           FROM QR FACTORIZATION OF THE MATRIX U DESCRIBED ABOVE.              CYC01650
C           OUTPUT. (DOUBLE PRECISION)                                          CYC01660
C  C      : WORK ARRAY OF DIMENSION (0:KMAX). FOR EACH K, C FOR MPE IS         CYC01670
C           THE LEAST SQUARES SOLUTION OF THE SYSTEM U*C=0 SUBJECT TO           CYC01680
C           THE CONSTRAINT C(K)=1. (DOUBLE PRECISION)                           CYC01690
C  GAMMA  : WORK ARRAY OF DIMENSION (0:KMAX). FOR EACH K, THE GAMMA'S          CYC01700
C           ARE SUCH THAT                                                       CYC01710
C           S(0,K)=GAMMA(0)*X0+GAMMA(1)*X1+...+GAMMA(K)*XK.                     CYC01720
C           FOR EACH K, GAMMA FOR RRE IS THE LEAST SQUARES SOLUTION OF          CYC01730
C           THE SYSTEM U*GAMMA=0 SUBJECT TO THE CONSTRAINT                      CYC01740
C           GAMMA(0)+GAMMA(1)+...+GAMMA(K)=1. (DOUBLE PRECISION)                CYC01750
C  XI     : WORK ARRAY OF DIMENSION (0:KMAX-1). FOR EACH K, THE XI'S           CYC01760
C           ARE SUCH THAT                                                       CYC01770
C           S(0,K)=X0+XI(0)*U0+XI(1)*U1+...+XI(J)*UJ, J=K-1.                    CYC01780
C           (DOUBLE PRECISION)                                                  CYC01790
C  RES    : L2-NORM OF THE RESIDUAL FOR S(0,K) FOR A LINEAR SYSTEM             CYC01800
C           T=A*T+B (OR AN ESTIMATE FOR IT FOR A NONLINEAR SYSTEM               CYC01810
C           T=F(T)) FOR EACH K. ON EXIT, THIS K IS KOUT. OUTPUT.                CYC01820
C           (DOUBLE PRECISION)                                                  CYC01830
C  RES1   : L2-NORM OF THE RESIDUAL ACTUALLY COMPUTED FROM S(0,K) FOR          CYC01840
C           EACH K. (THE RESIDUAL VECTOR FOR ANY VECTOR VEC IS TAKEN            CYC01850
C           AS (F(VEC)-VEC).) ON EXIT, THIS K IS KOUT. OUTPUT.                  CYC01860
C           (DOUBLE PRECISION)                                                  CYC01870
C  EPS    : AN UPPER BOUND ON RES/R(0,0), THE RELATIVE RESIDUAL FOR S,         CYC01880
C           USED IN THE STOPPING CRITERION. NOTE THAT R(0,0)=L2-NORM            CYC01890
C           OF THE RESIDUAL FOR X0, THE INITIAL VECTOR. IF, FOR SOME K,         CYC01900
C           RES.LE.EPS*R(0,0), THEN THE CORRESPONDING S(0,K) IS ACCEPTED        CYC01910
C           AS THE FINAL APPROXIMATION, AND THE SUBROUTINE IS EXITED            CYC01920
C           WITH KOUT=K. IF S(0,KMAX) IS NEEDED, THEN EPS SHOULD BE             CYC01930
C           SET EQUAL TO ZERO. INPUT. (DOUBLE PRECISION)                        CYC01940
C  IPRES  : IF IPRES.EQ.1, THEN RES IS PRINTED FOR ALL K, K=0,1,... .           CYC01950
C           OTHEWISE, IT IS NOT. INPUT. (INTEGER)                               CYC01960
C  IPRES1: IF IPRES1.EQ.1, THEN RES1 IS COMPUTED AND PRINTED FOR ALL            CYC01970
C           K, K=0,1,... . OTHERWISE, IT IS NOT. INPUT. (INTEGER)               CYC01980
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C*************************************************************************CYC01990
C  THE ABOVE MENTIONED QR FACTORIZATION IS PERFORMED BY EMPLOYING        CYC02000
C  THE MODIFIED GRAM-SCHMIDT PROCESS.                                    CYC02010
C*************************************************************************CYC02020
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)                                CYC02030
      PARAMETER (EPS1=1D-32,EPS2=1D-16)                                  CYC02040
      DIMENSION X(NDIM),S(NDIM),Y(NDIM),Z(NDIM)                          CYC02050
      DIMENSION Q(NDIM,0:KMAX-1),R(0:KMAX,0:KMAX)                        CYC02060
      DIMENSION C(0:KMAX),GAMMA(0:KMAX),XI(0:KMAX-1)                     CYC02070
      IF (IPRES.EQ.1.AND.IPRES1.EQ.1) THEN                               CYC02080
      WRITE(6,301)                                                       CYC02090
 301  FORMAT(/,'  K          RES             RES1')                      CYC02100
      ELSE IF (IPRES.EQ.1.AND.IPRES1.NE.1) THEN                          CYC02110
      WRITE(6,302)                                                       CYC02120
 302  FORMAT(/,'  K          RES')                                       CYC02130
      ELSE IF (IPRES.NE.1.AND.IPRES1.EQ.1) THEN                          CYC02140
      WRITE(6,303)                                                       CYC02150
 303  FORMAT(/,'  K          RES1')                                      CYC02160
      END IF                                                             CYC02170
      DO 10 I=1,NDIM                                                     CYC02180
      Y(I)=X(I)                                                          CYC02190
  10  CONTINUE                                                           CYC02200
      DO 250 K=0,KMAX                                                    CYC02210
C                                                                        CYC02220
C  COMPUTATION OF THE VECTOR XJ, J=K+1, FROM XK, AND COMPUTATION OF UK   CYC02230
C                                                                        CYC02240
      CALL VECTOR(Y,Z,NDIM)                                              CYC02250
      DO 20 I=1,NDIM                                                     CYC02260
      Y(I)=Z(I)-Y(I)                                                     CYC02270
  20  CONTINUE                                                           CYC02280
C                                                                        CYC02290
C  DETERMINATION OF THE ORTHONORMAL VECTOR QK FROM UK BY THE MODIFIED    CYC02300
C  GRAM-SCHMIDT PROCESS                                                  CYC02310
C                                                                        CYC02320
      DO 50 J=0,K-1                                                      CYC02330
      SUM=0                                                              CYC02340
      DO 30 I=1,NDIM                                                     CYC02350
      SUM=SUM+Q(I,J)*Y(I)                                                CYC02360
  30  CONTINUE                                                           CYC02370
      R(J,K)=SUM                                                         CYC02380
      DO 40 I=1,NDIM                                                     CYC02390
      Y(I)=Y(I)-R(J,K)*Q(I,J)                                            CYC02400
  40  CONTINUE                                                           CYC02410
  50  CONTINUE                                                           CYC02420
      SUM=0                                                              CYC02430
      DO 60 I=1,NDIM                                                     CYC02440
      SUM=SUM+Y(I)**2                                                    CYC02450
  60  CONTINUE                                                           CYC02460
      R(K,K)=DSQRT(SUM)                                                  CYC02470
      IF (R(K,K).GT.EPS1*R(0,0).AND.K.LT.KMAX) THEN                      CYC02480
      HP=1D0/R(K,K)                                                      CYC02490
      DO 70 I=1,NDIM                                                     CYC02500
      Q(I,K)=HP*Y(I)                                                     CYC02510
  70  CONTINUE                                                           CYC02520
      ELSE IF (R(K,K).LE.EPS1*R(0,0)) THEN                               CYC02530
      EEE=EPS1                                                           CYC02540
      WRITE(6,304) K,K,EEE                                               CYC02550
 304  FORMAT(/,'  R(',I3,',',I3,') .LE.',1P,D8.1,'*R(0,0).',/)          CYC02560
      END IF                                                             CYC02570
C                                                                        CYC02580
C  END OF COMPUTATION OF THE VECTOR QK                                   CYC02590
C                                                                        CYC02600
      IF (METHOD.EQ.1) THEN                                              CYC02610
C                                                                        CYC02620
C  COMPUTATION OF THE GAMMA'S FOR MPE                                    CYC02630
C                                                                        CYC02640
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```
      DO 90 I=K-1,0,-1                                                  CYC02650
      CI=-R(I,K)                                                        CYC02660
      DO 80 J=I+1,K-1                                                   CYC02670
      CI=CI-R(I,J)*C(J)                                                 CYC02680
   80 CONTINUE                                                          CYC02690
      C(I)=CI/R(I,I)                                                    CYC02700
   90 CONTINUE                                                          CYC02710
      C(K)=1D0                                                          CYC02720
      SUM=0                                                             CYC02730
      DO 100 I=0,K                                                      CYC02740
      SUM=SUM+C(I)                                                      CYC02750
  100 CONTINUE                                                          CYC02760
      IF (DABS(SUM).LE.EPS2) THEN                                       CYC02770
      WRITE(6,311) K                                                    CYC02780
  311 FORMAT(/,'  S(  0,',I3,') IS NOT DEFINED.',/)                     CYC02790
      GO TO 250                                                         CYC02800
      END IF                                                            CYC02810
      DO 110 I=0,K                                                      CYC02820
      GAMMA(I)=C(I)/SUM                                                 CYC02830
  110 CONTINUE                                                          CYC02840
      RES=R(K,K)*DABS(GAMMA(K))                                         CYC02850
C                                                                       CYC02860
C  END OF COMPUTATION OF THE GAMMA'S FOR MPE                            CYC02870
C                                                                       CYC02880
      ELSE IF (METHOD.EQ.2) THEN                                        CYC02890
C                                                                       CYC02900
C  COMPUTATION OF THE GAMMA'S FOR RRE                                   CYC02910
C                                                                       CYC02920
      DO 130 I=0,K                                                      CYC02930
      CI=1D0                                                            CYC02940
      DO 120 J=0,I-1                                                    CYC02950
      CI=CI-R(J,I)*C(J)                                                 CYC02960
  120 CONTINUE                                                          CYC02970
      C(I)=CI/R(I,I)                                                    CYC02980
  130 CONTINUE                                                          CYC02990
      DO 150 I=K,0,-1                                                   CYC03000
      CI=C(I)                                                           CYC03010
      DO 140 J=I+1,K                                                    CYC03020
      CI=CI-R(I,J)*GAMMA(J)                                             CYC03030
  140 CONTINUE                                                          CYC03040
      GAMMA(I)=CI/R(I,I)                                                CYC03050
  150 CONTINUE                                                          CYC03060
      SUM=0                                                             CYC03070
      DO 160 I=0,K                                                      CYC03080
      SUM=SUM+GAMMA(I)                                                  CYC03090
  160 CONTINUE                                                          CYC03100
      DO 170 I=0,K                                                      CYC03110
      GAMMA(I)=GAMMA(I)/SUM                                             CYC03120
  170 CONTINUE                                                          CYC03130
      RES=1D0/DSQRT(DABS(SUM))                                          CYC03140
C                                                                       CYC03150
C  END OF COMPUTATION OF THE GAMMA'S FOR RRE                            CYC03160
C                                                                       CYC03170
      END IF                                                            CYC03180
      KOUT=K                                                            CYC03190
      IF (IPRES.EQ.1.AND.IPRES1.NE.1) THEN                              CYC03200
      WRITE(6,321) K,RES                                                CYC03210
  321 FORMAT(I3,2X,1P,D15.2)                                            CYC03220
      END IF                                                            CYC03230
      IF (RES.LE.EPS*R(0,0).OR.R(K,K).LE.EPS1*R(0,0)                    CYC03240
     *    .OR.K.EQ.KMAX.OR.IPRES1.EQ.1) THEN                            CYC03250
C                                                                       CYC03260
C  COMPUTATION OF THE APPROXIMATION S(0,K)                              CYC03270
C                                                                       CYC03280
      XI(0)=1D0-GAMMA(0)                                                CYC03290
      DO 180 J=1,K-1                                                    CYC03300
```

```
      XI(J)=XI(J-1)-GAMMA(J)                                            CYC03310
 180  CONTINUE                                                          CYC03320
      DO 190 I=1,NDIM                                                   CYC03330
      S(I)=X(I)                                                         CYC03340
 190  CONTINUE                                                          CYC03350
      DO 220 J=0,K-1                                                    CYC03360
      HP=0                                                              CYC03370
      DO 200 I=J,K-1                                                    CYC03380
      HP=HP+R(J,I)*XI(I)                                                CYC03390
 200  CONTINUE                                                          CYC03400
      DO 210 I=1,NDIM                                                   CYC03410
      S(I)=S(I)+HP*Q(I,J)                                               CYC03420
 210  CONTINUE                                                          CYC03430
 220  CONTINUE                                                          CYC03440
C                                                                       CYC03450
C  END OF COMPUTATION OF THE APPROXIMATION S(0,K)                       CYC03460
C                                                                       CYC03470
      END IF                                                            CYC03480
      IF (IPRES1.EQ.1) THEN                                             CYC03490
C                                                                       CYC03500
C  EXACT COMPUTATION OF RESIDUAL L2-NORM.                               CYC03510
C                                                                       CYC03520
      CALL VECTOR(S,Y,NDIM)                                             CYC03530
      RES1=0                                                            CYC03540
      DO 230 I=1,NDIM                                                   CYC03550
      RES1=RES1+(Y(I)-S(I))**2                                          CYC03560
 230  CONTINUE                                                          CYC03570
      RES1=DSQRT(RES1)                                                  CYC03580
C                                                                       CYC03590
C  END OF EXACT COMPUTATION OF RESIDUAL L2-NORM.                        CYC03600
C                                                                       CYC03610
      IF (IPRES.EQ.1) THEN                                              CYC03620
      WRITE(6,331) K,RES,RES1                                           CYC03630
 331  FORMAT(I3,2X,1P,2D15.2)                                           CYC03640
      ELSE IF (IPRES.NE.1) THEN                                         CYC03650
      WRITE(6,332) K,RES1                                               CYC03660
 332  FORMAT(I3,2X,1P,D15.2)                                            CYC03670
      END IF                                                            CYC03680
      END IF                                                            CYC03690
      IF (RES.LE.EPS*R(0,0).OR.R(K,K).LE.EPS1*R(0,0)) RETURN            CYC03700
      DO 240 I=1,NDIM                                                   CYC03710
      Y(I)=Z(I)                                                         CYC03720
 240  CONTINUE                                                          CYC03730
 250  CONTINUE                                                          CYC03740
      RETURN                                                            CYC03750
      END                                                               CYC03760
                                                                        CYC03770
      SUBROUTINE VECTOR(X,Y,NDIM)                                       CYC03780
C***********************************************************************CYC03790
C  THIS SUBROUTINE GENERATES THE VECTOR Z FROM THE VECTOR Y BY USING,    CYC03800
C  E.G., A FIXED POINT ITERATION TECHNIQUE.                             CYC03810
C***********************************************************************CYC03820
C  IN THE PRESENT EXAMPLE THE ITERATIVE TECHNIQUE IS OF THE FORM        CYC03830
C  Y=A1*X+B1. HERE A1 IS AN NDIM*NDIM  SEPTADIAGONAL MATRIX SYMMETRIC   CYC03840
C  WITH RESPECT TO BOTH OF ITS DIAGONALS, AND IS DEFINED AS             CYC03850
C  A1=(1-OMEGA)*I+OMEGA*A,WHERE OMEGA IS A SCALAR, I IS THE             CYC03860
C  IDENTITY MATRIX, AND A IS THE MATRIX                                 CYC03870
C                                                                       CYC03880
C             | 5  2  1  1                    |                        CYC03890
C             | 2  6  3  1  1                 |                        CYC03900
C             | 1  3  6  3  1  1              |                        CYC03910
C   A = 0.06*| 1  1  3  6  3  1  1           | .                      CYC03920
C             |    1  1  3  6  3  1  1        |                        CYC03930
C             |       1  1  3  6  3  1  1     |                        CYC03940
C             |          .  .  .  .  .  .  .  |                        CYC03950
C                                                                       CYC03960
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C  B1 IS THE VECTOR DEFINED AS B1=OMEGA*B, THE VECTOR B BEING CHOSEN     CYC03970
C  SUCH THAT THE SOLUTION OF THE SYSTEM T=A*T+B IS THE VECTOR            CYC03980
C  (1,1,...,1).                                                         CYC03990
C  THE ITERATIVE TECHNIQUE USED IS THUS RICHARDSON'S ITERATIVE           CYC04000
C  METHOD APPLIED TO THE SYSTEM (I-A)*T=B.                               CYC04010
C*********************************************************************CYC04020
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)                               CYC04030
      PARAMETER (OMEGA=2D0,TAU=1D0-OMEGA)                               CYC04040
      DIMENSION X(NDIM),Y(NDIM)                                         CYC04050
      N=NDIM                                                            CYC04060
      Y(1)=(5*X(1)+2*X(2)+X(3)+X(4))*6D-2+46D-2                         CYC04070
      Y(2)=(2*X(1)+6*X(2)+3*X(3)+X(4)+X(5))*6D-2+22D-2                  CYC04080
      Y(3)=(X(1)+3*X(2)+6*X(3)+3*X(4)+X(5)+X(6))*6D-2+1D-1              CYC04090
      DO 10 I=4,N-3                                                     CYC04100
      Y(I)=(X(I-3)+X(I-2)+3*X(I-1)+6*X(I)+3*X(I+1)+X(I+2)+X(I+3))*6D-2  CYC04110
     *     +4D-2                                                        CYC04120
   10 CONTINUE                                                          CYC04130
      Y(N-2)=(X(N)+3*X(N-1)+6*X(N-2)+3*X(N-3)+X(N-4)+X(N-5))*6D-2+1D-1  CYC04140
      Y(N-1)=(2*X(N)+6*X(N-1)+3*X(N-2)+X(N-3)+X(N-4))*6D-2+22D-2        CYC04150
      Y(N)=(5*X(N)+2*X(N-1)+X(N-2)+X(N-3))*6D-2+46D-2                   CYC04160
      DO 20 I=1,N                                                       CYC04170
      Y(I)=TAU*X(I)+OMEGA*Y(I)                                          CYC04180
   20 CONTINUE                                                          CYC04190
      RETURN                                                            CYC04200
      END                                                               CYC04210
```

CYCLE NO. 1

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 20

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	4.75D-01	4.75D-01
1	5.36D-01	5.36D-01
2	1.52D-02	1.52D-02
3	1.93D-02	1.93D-02
4	4.23D-03	4.23D-03
5	3.79D-03	3.79D-03
6	1.41D-03	1.41D-03
7	1.00D-03	1.00D-03
8	5.16D-04	5.16D-04
9	3.04D-04	3.04D-04
10	2.00D-04	2.00D-04

CYCLE NO. 2

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 0

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	2.00D-04	2.00D-04
1	9.57D-05	9.57D-05
2	9.59D-05	9.59D-05
3	4.58D-05	4.58D-05
4	4.42D-05	4.42D-05
5	1.68D-05	1.68D-05
6	1.91D-05	1.91D-05
7	6.49D-06	6.49D-06
8	7.22D-06	7.22D-06
9	2.56D-06	2.56D-06
10	2.90D-06	2.90D-06

CYCLE NO. 3

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 0

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	2.90D-06	2.90D-06
1	1.18D-06	1.18D-06
2	1.38D-06	1.38D-06
3	6.20D-07	6.20D-07
4	6.64D-07	6.64D-07
5	2.43D-07	2.43D-07
6	2.63D-07	2.63D-07
7	8.58D-08	8.58D-08
8	9.15D-08	9.15D-08
9	3.95D-08	3.95D-08
10	4.17D-08	4.17D-08

CYCLE NO. 4

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 0

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	4.17D-08	4.17D-08
1	2.44D-08	2.44D-08

2	2.40D-08	2.40D-08
3	1.29D-08	1.29D-08
4	1.24D-08	1.24D-08
5	5.49D-09	5.49D-09
6	5.39D-09	5.39D-09
7	1.95D-09	1.95D-09
8	1.96D-09	1.96D-09
9	8.71D-10	8.71D-10
10	9.27D-10	9.27D-10

CYCLE NO. 5

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 0

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	9.27D-10	9.27D-10
1	5.14D-10	5.14D-10
2	5.46D-10	5.46D-10
3	2.74D-10	2.74D-10
4	2.71D-10	2.71D-10
5	1.17D-10	1.17D-10
6	1.09D-10	1.09D-10
7	4.64D-11	4.64D-11
8	4.28D-11	4.28D-11
9	2.07D-11	2.07D-11
10	2.19D-11	2.18D-11

CYCLE NO. 6

NO. OF ITERATIONS PRIOR TO EXTRAPOLATION IS 0

WIDTH OF EXTRAPOLATION IS 10

K	RES	RES1
0	2.18D-11	2.18D-11
1	1.25D-11	1.25D-11
2	1.26D-11	1.26D-11
3	7.23D-12	7.23D-12
4	6.32D-12	6.32D-12
5	3.29D-12	3.28D-12
6	2.85D-12	2.85D-12
7	1.27D-12	1.28D-12
8	1.15D-12	1.15D-12
9	5.79D-13	5.67D-13
10	5.67D-13	5.49D-13

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16. Abstract

The minimal polynomial extrapolation (MPE) and reduced rank extrapolation (RRE) are two very effective techniques that have been used in accelerating the convergence of vector sequences, such as those that are obtained from iterative solution of linear and nonlinear systems of equations. Their definitions involve some linear least squares problems, and this causes difficulties in their numerical implementation. In this work timewise efficient and numerically stable implementations for MPE and RRE are developed. A computer program written in FORTRAN 77 is also appended and applied to some model problems.

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