# "OPTIMAL PAYLOAD RATE LIMIT ALGORITHM FOR ZERO-G MANIPULATORS" 

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#### Abstract

An algorithm for continuously computing safe maximum relative velocities for two bodies joined by a manipulator is discussed. The maximum velocities are such that if the brakes are applied at that instant, the ensuing travel between the bodies will be less than or equal to a predetermined amount. This paper deals with an improvement in the way this limit is computed for space manipulators. The new method is explained, test cases are posed, and the results of these tests are displayed and discussed.


## I. Introduction

## A. What is a payload rate limit algorithm?

The rate limit for a payload is, in effect, a "speed limit" for the rates the payload is allowed to achieve relative to the manipulating body. Observance of this rate limit ensures that the payload's relative motion can be arrested at any time during a maneuver (e.g., in an emergency such as a joint runaway detection) without exceeding a specified amount of overtravel after application of the brakes. Any method employed to compute these rates must consider items like brake torque capability of the manipulator and masses of the payload and manipulator base.

The Remote Manipulator System (RMS) on NASA's Space Shuttle currently utilizes a rate limit algorithm to compute a single (constant) rate limit for each payload to be manipulated during a mission, and these rates are loaded into the on-board computer system prior to the mission. These rates are designed to ensure a stopping distance of the end-effector (not the payload c.g.) in two feet or less when the arm is in an outstretched position (worst-case).

## B. What is wrong with the current rate limit algorithm?

The algorithms developed to determine the rate limit for a payload to be carried by the RMS were developed by SPAR Aerospace, Inc. in August 1979 with revisions in February 1983 [Ref. 1,2]. These algorithms were designed for the original payload mass range of up to $65,000 \mathrm{lbm}$. (vs. Shuttle mass of $220,000 \mathrm{lbm}$ ), and also assumed a worstcase RMS configuration (fully extended). While these algorithms have performed well for the range of payloads seen so far, they produce rate limits approaching zero for payloads whose masses exceed $100,000 \mathrm{lbm}$. They also produce only one limit, based on the worstcase configuration, when in fact the RMS's ability to arrest relative motion is dependent on arm configuration and the direction of the motion to be arrested. This means the rate
limits calculated are lower than necessary (conservative), so the payload is always manipulated at rates lower than actually required for safety.

## C. Why should we care?

The RMS manipulator is basically a dexterous crane that has to perform largeexcursion travel in many of its tasks. If used, the low rate limits that the current algorithm produces for the massive payloads will cause increases in the time required to perform the job because the arm will move much slower than is necessary for safety. The low rate limits also push all the commands down into the low end of the command scale. This lessens the number of bit-states which can represent the command, thereby decreasing resolution and accuracy of the command actually sent to the joint servos.

The current algorithm must be changed to consider payloads more massive than 65000 lbm ., and must consider the fact that the Shuttle also moves during the braking process. Since the algorithm will require at least an update to handle the larger payloads, this is a good time to look into deriving a better method. This method should be dependent on the arm configuration and commanded velocity.

Some analysis on improving the algorithm has been done before, but the algorithm arising from that analysis was still a single limit for a payload [Ref. 3]. That method did, however, take into account the fact that the Shuttle moves during manipulation, and was the basis of this current effort.

## II. Proposed New Payload Rate Limit Algorithm

## A. What should it do?

Any new method should include better models of system dynamics, and yet be computationally simple. Simplicity will enable calculation of the limit in real-time as a function of the current arm configuration and the requested direction of motion. The new method should also consider relative rotational motion, and allow a maximum rotation angle as an additional criterion.

The end goal of this algorithm should be to produce a rate limit which is higher than the currently-used conservative value yet is still safe.

## B. How can it be done?

## 1. System Dynamics

By considering the system dynamics of the payload and the Shuttle over the process of:

Phase 1. beginning with relative rates (possibly zero)
between Shuttle and payload,
Phase 2. accelerating the payload and Shuttle to some new commanded relative rate, and finally

Phase 3. applying the brakes until all relative rates have been arrested,
one can see that, since we assume no external forces, system momenta is conserved throughout this process. The dynamics from the time of brake actuation to arrest of relative motion (Phase 3) can be viewed as an inelastic collision, in which both bodies have initial relative motion, collide and stick together, and then proceed on as one body.

Realizing that the final system velocity $\mathrm{V}_{\mathrm{f}}$ is constant, we can define a reference frame F , translating at velocity $\mathrm{V}_{\mathrm{f}}$. The final rotational rates of the system are expected to be very small (less than 0.2 degrees per second) This means we can allow the frame $F$ to rotate with the system at its final rotational velocity $\omega_{f}$ and still consider it inertial (pseudo-inertial). In this frame, an observer would see that each body would have an initial velocity, but would come to rest at impact. The relative velocities between these two bodies would be the commanded velocity. The momentum equation can be written for each body in this pseudo-inertial frame:

$$
m_{1} V_{1}+\int_{0}^{t} \underset{\rightarrow}{F}(t) d t=0 \text { and } m_{2} V_{2}-\int_{0}^{t} \underset{\rightarrow}{F}(t) d t=0
$$

so

$$
\underline{V}_{1}=-\frac{m_{2}}{m_{1}} \underline{V}_{2}
$$

which says that the initial velocities of the two bodies in this frame will always be opposed and parallel. Because the common direction of these velocities is also the direction of the relative velocity, and we assume the force between the bodies to be opposed to the relative velocity so everything falls on this line, leaving us a scalar equation.

Writing the equation of linear momentum for the system in frame $F$, we get the following equations, which we can treat as scalar since all vectors are parallel.

For the entire system (no external impulse) :

$$
m_{1} V_{1}+m_{2} V_{2}=0
$$

For bodies 1 and 2 individually (external impulse from arm) :

$$
m_{1} V_{1}+\int_{0}^{t} F(t) d t=0 \quad m_{2} V_{2}-\int_{0}^{t} F(t) d t=0
$$

where $\left(V_{2}-V_{1}\right)$ is the commanded velocity.

## 2. Estimation of Impulse

If we had some idea about what kind of impulse we could expect to see, we could extrapolate what the travel motion of the system would be. The impulse does not have to be exactly known, but the estimate must be less than the actual impulse, and still
produce a rate faster than the old value. A way to estimate the impulse is to estimate the initial force or torque, and then estimate the manner in which it approaches zero.

Any force applied by the RMS to the bodies it connects will result from brake slippage and storage of strain energy in the booms and gearboxes. Since only brake slippage is non-conservative, this mechanism alone will be considered. The potential brake torque at each joint is assumed to be known. In this case, the initial resistive force encountered by the bodies can be estimated. By assuming all joints required to contribute to the commanded motion will slip if the brakes are applied, a vector of brake torques can be created which is an estimate of the torques which would be seen at each joint:

$$
(\underline{B})=-\left(\begin{array}{c}
\tau_{1} * \operatorname{sign}\left(\dot{\gamma}_{1}\right) \\
\vdots \\
\tau_{6} * \operatorname{sign}\left(\dot{\gamma}_{6}\right)
\end{array}\right)
$$

where $\operatorname{sign}()$ is -1 if the commanded joint rate is negative, zero if zero, and +1 if positive.
Here we assume that the inverse of the Jacobian matrix relating end-effector states to joint states is known or easily obtained. Multiplication of the transposed inverse of the Jacobian by the brake torque vector just constructed results in the static moments and forces at the end-effector to counteract those brake torques.

$$
\binom{\mathrm{T}}{\overrightarrow{\underline{E}}}=[J]^{-T}(\underline{B})
$$

These loads at the end-effector will not generally be parallel to their counterpart velocities. The components of these loads in the direction of those velocities are an estimate of initial loads the payload and shuttle will encounter, while the components normal to those velocities are assumed to arise from the errors in the assumed joint torque vector.

$$
\begin{aligned}
& T_{\max }=\underline{T} \cdot \widehat{\omega}_{c} \\
& \mathrm{~F}_{\max }=\overrightarrow{\mathrm{F}} \cdot \hat{V}_{\mathrm{c}}
\end{aligned}
$$

where

$$
\widehat{V}_{c}, \hat{\omega}_{c} \text { are the unit vectors of the commanded velocities, and } T_{\max }
$$ $F_{\text {max }}$ are the available torque and force in those directions (scalar). These loads are the estimates of the actual initial loads seen upon applying the manipulator brakes.

So far this method predicts the initial loads seen at the beginning of the braking, but its behavior of these loads over time is unknown. Since we are only having to compute a conservative limit, we only need assume a conservative load profile, i.e. the assumed profile's integral must always be less than or equal to the actual impulse. Nominally, loads due to brake slippage could be thought of as constant over time, but simulation has shown that these loads tend to drop off over the braking maneuver. To be conservative, we assume a profile in which the loads begin at the predicted value ( $\mathrm{F}_{\mathrm{max}}$, $\mathrm{T}_{\max }$ ), but then linearly ramp down to zero over the time of the maneuver.

## 3. Calculation of Maximum Safe Speed

## a. Linear Velocity

Once we know or assume a force-time profile (impulse), we can determine the distance travelled over that time and make sure all distances travelled are less than 2 feet (or some other criterion).

To do this, we must look at the equations of motion of these bodies along the line of action in frame F :

$$
\mathrm{F}(\mathrm{t})=\mathrm{A}-\mathrm{Bt} \quad\{0<\mathrm{t}<\text { tend }\}
$$

where $F(t)$ is the force exerted on each body,
tend is the time at which motion stops
A is the maximum force $F_{\text {max }}$ computed above
$B$ is the slope at which $F(t)$ ramps down,

$$
\text { i.e., } F_{\max } / \text { tend }
$$

so the acceleration of each mass is

$$
\mathrm{a}(\mathrm{t})=(\mathrm{A}-\mathrm{Bt}) / \mathrm{m}
$$

and the velocity of each mass is

$$
\begin{aligned}
& v(t)= \\
& \qquad \int_{0}^{\text {terd }} a(t) d t+V(t=0) \\
& =\left(A t-B t^{2} / 2\right) / m \quad(0<t<t \text { end })
\end{aligned}
$$

and the distance travelled by each mass is

$$
\begin{aligned}
& \mathbf{x}(\mathrm{t})= \\
& \quad \int_{0}^{t_{\operatorname{tax}}} V(\mathrm{t}) \mathrm{dt}+X(\mathrm{t}=0) \\
& =\left(A t^{2} / 2-B t^{3} / 6\right) / \mathrm{m}\left(0<t<t_{\text {end }}\right)
\end{aligned}
$$

We can also see that since

$$
x_{1}\left(t=t_{\text {end }}\right)+x_{2}\left(t=t_{\text {end }}\right)<2 \text { feet }
$$

we can back out a relationship between the $\mathrm{F}_{\max }$ and the amount of time required to arrest the motion, tend.

$$
2 \text { feet }=\left(\mathrm{At}_{\text {end }} 2 / 2-\mathrm{Btend}^{3 / 6}\right)\left(1 / \mathrm{m}_{1}+1 / \mathrm{m}_{2}\right)
$$

or, substituting for $A=F_{\max }$, and $B=F_{\max } /$ end, we get

$$
t_{\text {end }}=\quad \sqrt{\left[3 * d^{*} \mathrm{~m}_{1} * \mathrm{~m}_{2} /\left(\left(\mathrm{m}_{1}+\mathrm{m}_{2}\right) * \mathrm{~F}_{\max }\right)\right]}
$$

where d is the stopping distance allowed, which is 2 feet for the RMS.
Once we know the amount of time required to arrest the motion, we can use the momentum equations for each body:

$$
\mathrm{m}_{1} \mathrm{~V}_{1}+\text { impulse }=0, \quad \mathrm{~m}_{2} \mathrm{~V}_{2}+\text { impulse }=0
$$

and since the impulse is the area under the assumed force-time curve,

$$
\text { impulse }=1 / 2 \mathrm{~F}_{\max } \mathrm{T}_{\mathrm{end}}
$$

and the final allowable command velocity is $\mathrm{V}_{2}+\mathrm{V}_{1}$,

$$
\begin{gathered}
\mathrm{V}_{\max }=\left(\mathrm{F}_{\max } * \operatorname{tend}^{\operatorname{man}}\right) *\left[1 / \mathrm{m}_{1}+1 / \mathrm{m} 2\right] \\
\text { b. Rotational Velocity }
\end{gathered}
$$

The computation of the maximum safe rotational velocity closely parallels that of the translational. An assumption of the torque-time profile (linearly approaching zero) is made, and the equation of angular momentum is written for the system about the endeffector tip. This requires that all body inertias be computed about that point.

The assumed torque-time profile is

$$
\mathrm{T}(\mathrm{t})=\mathrm{A}-\mathrm{Bt}(0<\mathrm{t}<\text { tend })
$$

where $A$ is $T_{\max }$ (initial torque)
$B$ is the slope at which it ramps down ( $\mathrm{T}_{\mathrm{max}} / \mathrm{t}_{\text {end }}$ ).
For each body,

$$
\mathrm{T}(\mathrm{t})=\mathrm{I} \alpha
$$

where
I is the scalar moment of inertia at the end-effector about the rotation axis, and
$\alpha \quad$ is the rotational acceleration.
Solving for $\alpha, \Delta \omega$, and $\Delta \theta$

$$
\alpha=\mathrm{I}^{-1}[\mathrm{~A}-\mathrm{Bt}]
$$

and

$$
\Delta \omega=I^{-1}\left[A t-B t^{2} / 2\right]
$$

and

$$
\Delta \theta=\mathrm{I}^{-1}\left[\mathrm{At}^{2} / 2-\mathrm{Bt}^{3} / 6\right]
$$

As before, we set the travel constraint

$$
10 \text { degrees }=\Delta \theta 1+\Delta \theta 2 \quad(\text { at } t=\text { tend })
$$

or

$$
\theta_{\max }=10 \text { degrees }=\left(\mathrm{I}_{1}^{-1}+\mathrm{I}_{2}^{-1}\right)\left[\mathrm{T}_{\max } \text { tend }{ }^{2 / 3}\right]
$$

which yields the expression for tend

$$
\text { tend }=\quad \sqrt{\left(\left(\theta_{\max } / 57.3 * \mathrm{I}_{1} \mathrm{I}_{2} * 3\right) /\left(\left(\mathrm{I}_{1}+\mathrm{I}_{2}\right) * \mathrm{~T}_{\max }\right)\right), ~}
$$

which yields the expression of commanded velocity

$$
\omega_{\max }=\left(\mathrm{T}_{\max } \operatorname{tend} / 2\right) *\left[I_{1}^{-1}+I_{2}^{-1}\right]
$$

much as the translational equation.
This $V_{\text {max }}$ and $\omega_{\text {max }}$ are the magnitudes of the allowable velocity in the commanded direction. This command is what would be sent on to the robot controller, if the operator desired to go the maximum safe speed.

## 4. Application of Limit to Command

Obviously, an operator would not want to go the maximum safe speed in all situations, so what to do with the knowledge of this instantaneous speed limit raises several possibilities. One way to apply the limit would be to use some constant conservative limits under normal operation, and use the maximum limits whenever a hand controller has been fully deflected in some axis. Another method would be to use the maximum limit as the upper end of the hand-controller's range, i.e. if the hand-controller is deflected $50 \%$ in some direction, then the commanded velocity would be $50 \%$ of the maximum safe velocity. This latter appears to be possibly unwieldy for the operator, since a constant deflection of the controller would produce varying velocity commands as the arm's configuration changes. Either of these methods could be implemented as an operator-requested mode.

## III. Tests, Results and Conclusions

A. How can we evaluate this new algorithm?

In order to evaluate the performance of the algorithm, dynamic simulation of the Shuttle/RMS/Payload system was required. This was done with a dynamic batch
simulation program developed at the Johnson Space Center which uses an extensive model of the Shuttle and the RMS. This program (called MIRRORS, for Model for Integrated Robotics Research and Operational Requirements Synthesis) is a spin-off of the PDRSS (Payload Deployment and Retrieval System Simulator, which is a spin-off of the SVDS (Space Vehicle Dynamic Simulator), which was written during the Apollo Program. The program has been checked against actual flight data from Shuttle/RMS missions as well as other simulators, and is used routinely for RMS maneuver simulation.

The underlying idea behind the evaluation was to get a prediction of safe velocity from the algorithm and then start the simulation with those velocities and the brakes on, watching the ensuing travel. This was done while varying command direction, payload mass and arm configuration. The commands given were all single-axis commands:

$$
\begin{array}{ll}
\mathbf{X} & - \\
\mathbf{Y} & \text { translate in positive Orbiter X-axis (toward nose) } \\
\mathbf{Z} & \text { translate in positive Y-axis (toward starboard wing) } \\
& \text { translate in positive Z-axis (down toward bay) } \\
\mathbf{R} & \\
\mathbf{P} & \text { rotate in positive direction about } \mathrm{X} \text {-axis (roll) } \\
\mathbf{Y}- & \text { rotate in positive direction about Y-axis (pitch) } \\
\text { rotate in positive direction about Z-axis (yaw). }
\end{array}
$$

The algorithm was coded into the flight software module of the MIRRORS program. Tests were run in the following manner for commands in each rotational and translational axis:

1. For given command direction, determine from the algorithm the maximum safe speed.
2. Initialize simulation with payload moving at that speed in the commanded direction, and with the brakes just applied.
3. Let simulation run until motion arrested, compare amount of travel with specified maximum ( 2 feet or 10 degrees in all cases).

The test runs were conducted for three different payloads, all 15 foot diameter homogeneous cylinders, grappled on the side at the midpoint, having masses of 32000 , 100000 , and 250000 lbm . The test also used two different initial arm configurations, for a total of $6 * 3 * 2=36$ test runs.

To evaluate the amount of travel, the code was altered to compute the Euclidean distance from the current position to the point where the brakes were applied. For the rotational cases, the angular displacement about the relative rotational eigen-axis (component of the quaternion relating payload attitude relative to the Orbiter) was computed. These are included in the test results below.
B. What were the test results?

The test results for all 36 runs were in general agreement with the desired end goals, in that 33 of 36 runs resulted in payloads travelling 2 feet $/ 10$ degrees or less while at speeds faster than the old method would allow. Three runs, however, did result in up to 2.3 feet of travel, all in the Z direction. The results of all the runs are tabled below, for
each arm configuration (joint angles from shoulder yaw to wrist roll shown in parentheses). The three runs that exceeded the travel limits are also noted.
The format for the data for each run (three numbers) is as follows:
[1] rate limit as predicted by old method (feet/sec or deg/sec)
[2] rate limit as predicted by new method
[3] stopping distance or angle for new rate (feet or degrees)
ARM CONFIGURATION \#1 (-90,90,-71,0,0,0):

| Pay- <br> load | X | $\stackrel{Y}{Y}$ | $\begin{aligned} & \text { ommand } \\ & \text { Z } \\ & \text { feet } \end{aligned}$ | $\underset{\text { degrees/sec }}{\mathbf{R}} \stackrel{\mathrm{P}}{\text { and degrees }}$ |  |  | old rate limit new rate limit stopping distance or angle <br> * Exceeded 2' criterion |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | feet $/ \mathrm{sec}$ and feet |  |  |  |  |  |  |
|  | 0.153 | 0.153 | 0.153 | 0.498 | 0.498 | 0.498 |  |
| 32K | 0.250 | 0.252 | 0.324 | 1.372 | 0.687 | 0.718 |  |
|  | 0.990 | 0.354 | 2.365* | 2.027 | 3.417 | 3.177 |  |
|  | 0.100 | 0.100 | 0.100 | 0.284 | 0.284 | 0.284 |  |
| 100K | 0.160 | 0.161 | 0.207 | 0.779 | 0.400 | 0.411 |  |
| 100K | 0.749 | 0.223 | 2.362* | 1.907 | 3.497 | 3.000 |  |
|  | 0.071 | 0.071 | 0.071 | 0.180 | 0.180 | 0.180 |  |
| 250K | 0.123 | 0.124 | 0.159 | 0.497 | 0.269 | 0.267 |  |
|  | 0.624 | 0.228 | 2.225* | 2.037 | 3.690 | 2.770 |  |

ARM CONFIGURATION \#2 ( $-48,118,-118,-26,-39,3$ ):


The amount of additional computer time required to compute this rate limit was negligible for the simulation, which is already numerically intensive. This is not an indication of its impact on some other manipulator, although the scheme doesn't require much numerical work provided the inverse of the Jacobian matrix is available.

## C. Any conclusions?

Since the tests did produce three runs which exceeded the 2 feet limit, one conclusion is that further work is needed in better estimating the impulse imparted between the bodies. However enough runs ( 33 out of 36 , or $92 \%$ ) not only stopped well within the limit, but at speeds faster than the present method would allow, which indicates the potential worth of this form of electronic safety monitoring. The conclusion of the study conducted to date is that further investigation is warranted to increase accuracy of the impulse estimation. If this can be improved and remain numerically simple, the algorithm will be a useful tool in speeding-up tasks for space robots.

## IV. REFERENCES

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