## NEURAL NETWORKS: WHAT NON-LINEARITY TO CHOOSE Vladik Kreinovich, Chris Quintan*

Abstract. Neural networks are now one of the most successful learning formalisms. Neurons transform inputs $x_{1}, \ldots, x_{n}$ into an output $f\left(w_{1} x_{1}+\ldots+w_{n} x_{n}\right)$, where $f$ is a non-linear function and $w_{i}$ are adjustable weights. What $f$ to choose? Usually the logistic function is chosen, but sometimes the use of different functions improves the practical efficiency of the network.

We formulate the problem of choosing $f$ as a mathematical optimization problem and solve it under different optimality criteria. As a result, we get a list of functions $f$ that are optimal under these criteria. This list includes both the functions that were empirically proved to be the best for some problems, and some new functions that may be worth trying.

## 1. FORMULATION OF THE PROBLEM.

Neural networks are now one of the most successful learning formalisms (see, e.g., the recent survey in (Hecht-Nielsen 1991]). After the initial success of linear neural models, in which the output $y$ is equal to the linear combination of the input signals $x_{i}$, ie. $y=w_{1} x_{1}+w_{2} x_{2}+\ldots$, it was shown in [Minsky Papert 1968] that if we only have linear neurons, then we end up with only linear functions and this severely limits the number of problems that we can solve using the network. The next step, then, is to consider non-linear neurons, in which the output signal is equal to $f\left(w_{1} x_{1}+w_{2} x_{2}+\ldots\right)$, where $f(y)$ is a given nonlinear function. A natural question arises: what function $f(y)$ do we choose?

Why is this problem important? It is a very important problem because although neural networks help us to design good learning procedures, these procedures are far from being reliable. Sometimes these procedures do not work; sometimes they work but demand too much time, and too big a sample, to learn. Naturally, we might think that this is because the function $f$ that we used was not the best one. Sometimes the use of different functions can improve the
$m^{i}$ cases of earning, such as average earning rime or average number oi errors, is optimal in these we do not know how to compute any of these possible characteristics. How can we And for 5 which $j(f)$ is optimal if we cannot compute $f(f)$ even or a single $f$ ? There does an seem to be a iike:y answer.

However, we will show that this problem is solvable (and give the solution) using advanced math, namely, group theory. (For a general idea of this approach, see [Kreinovich 1990].)

## 2. WHAT IS KNOWN.

The first non-linear neuron was proposed in [Gowan 1967]. Cowan chose the logistic function, $s_{0}(y)=1 /(1+\exp (-y))$, because it leads to a good approximation of the behavior of real (i.e. biological) neurons. The properties of neural networks with different $f$ were studied by properties useful for learning and is therefore an adequate choice. His analysis restricted the class of possible functions, but there are still many other functions win Hecht-Nielsen, who in [1987] added a demand that any function must be approximable by some neural network. This

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approach leads to non-trivial mathematics, but the final result (see, e.g., [Kreinovich 1991]) is that any smooth function $f$ will work, so this additional demand does not help us to choose $f$.


## 3. WHAT DO WE PROPOSE

### 3.1. Motivations of the proposed mathematical definitions.

We must choose a family of functions, not a single function. We speak about choosing $f$, but the expression for $f(y)$ will change if we change the units in which we measure all the signals (input, output and intermediate), so in mathematical terms, it is better to speak about choosing a family of functions $f$. It is reasonable to suggest that if an $f$ belongs to this family, then this family must contain $k f$ for positive real numbers $k$. This corresponds to changing units. Also, it must contain $f+c$, where $c$ is a constant. This is equivalent to adding a constant bias and therefore does not change the abilities of the resulting network. Since we are talking about non-linear phenomena, we can also assume that some non-linear "rescaling" transformations $x \rightarrow g(x)$ are also applicable, i.e., the family must include the composition $g(f(y))$ for each of functions $f$. This family must not be too big, therefore, it must be determined by finitely many parameters and should ideally be obtained from one function $f(y)$ by applying all these transformations. Without loss of generality, we can assume that this set of transformations is closed under composition and under inverse, i.e., if $z \rightarrow g_{1}(z)$ and $z \rightarrow g_{2}(z)$ are possible transformations, then $z \rightarrow g_{1}\left(g_{2}(z)\right)$ and $z \rightarrow g_{1}^{-1}(z)$ are possible transformations, where by $g_{1}^{-1}$ we denoted an inverse function $g_{1}^{-1}(z)=w$ if and only if $g_{1}(w)=z$. In mathematical terms this means that these transformations form a group, and therefore a family is obtained by applying to some function $f(y)$ all transformations from some finite-dimensional transformation group $G$ that includes all linear transformations (and maybe some non-linear ones).

All these transformations correspond to appropriate "rescalings". Rescaling is something that is smoothly changing the initial scale. This means that if we have two different transformations, there must be a smooth transition between them. In mathematical terms, the existence of this continuous transition is expressed by saying that the group is connected, and the fact that both the transformations and the transitions are smooth is expressed by saying that this is a Lie group (see, e.g., Chevalley, 1946).

What family is the best? Among all such families, we want to choose the best one. In formalizing what "the best" means we follow the general idea outlined in KKreinovich 1990). The criteria so choose may je computational simpiicity, sficiency oi training, or something eise. in mathemarical optimization problems. zumeric initeria are most trequently used, when to every family we assign some value expressing its performance, and choose a family tor which this value is maximal. However. it is not necessary to restric: surselves to such numeric criteria only. For example. if ve have several diferent Samilies that have the same raining ability .1. ve can cionse berween them the one that has the minimal computational compiexity $\tilde{C}$. In inis case. the actuai criterion that we use to compare two families is not numeric, but more complicated: a family $F_{1}$ is better than the family $F_{2}$ if and only if either $A\left(F_{1}\right)>A\left(F_{2}\right)$ or $A\left(F_{1}\right)=A\left(F_{2}\right)$ and $C\left(F_{1}\right)<C\left(F_{2}\right)$. A criterion can be even more complicated. What a criterion must do is to allow us for every pair of families to tell whether the first family is better with respect to this criterion (we'll denote it by $F_{1}>F_{2}$ ), or the second is better ( $F_{1}<F_{2}$ ) or these families have the same quality in the sense of this criterion (we'll denote it by $F_{1} \sim F_{2}$ ). Of course, it is necessary to demand that these choices be consistent, e.g., if $F_{1}>F_{2}$ and $F_{2}>F_{3}$ then $F_{1}>F_{3}$.

Another natural demand is that this criterion must choose a unique optimal family (i.e., a family that is better with respect to this criterion than any other family). The reason for this demand is very simple. If a criterion does not choose any family at all, then it is of no use. If several different families are "the best" according to this criterion, then we still have a problem to choose among those "best". Therefore, we need some additional criterion for that choice. For example, if several families turn out to have the same training ability, we can choose among them a family with minimal computational complexity. So what we actually do in this case is abandon that criterion for which there were several "best" families, and consider a new
"composite" criterion instead: $F_{1}$ is better than $F_{2}$ according to this new criterion if either it was better according to the old criterion or according to the old criterion they had the same quality and $F_{1}$ is better than $F_{2}$ according to the additional criterion. In other words, if a criterion does not allow us to choose a unique best family it means that this criterion is not ultimate; we have to modify it until we come to a final criterion that will have that property.

The next natural condition that the criterion must satisfy is connected with the following. Suppose that instead of a neuron with the transformation function $f(y)$ we consider a neuron with a function $\bar{f}(y)=f(y+a)$, where $a$ is a constant. This new neuron can be easily simulated by the old ones: namely, the output of this new neuron is $\bar{f}\left(w_{1} x_{1}+w_{2} x_{2}+\ldots\right)=f\left(w_{1} x_{1}+w_{2} x_{2}+\ldots+a\right)$, so it is equivalent to an old neuron with an additional constant input $a$. Likewise, the old neuron is equivalent to the new neuron with an additional constant input -a. Therefore, the networks that are formed by these new neurons have precisely the same abilities as those that are built from the old ones. We cannot claim that the new neurons have the same quality as the old ones, because adding $a$ can increase computational complexity and thus slightly worsen the overall quality. But it is natural to demand that adding a does not change the relative quality of the neurons, i.e., if a family $\{f(y)\}$ is better that a family of $\{g(y)\}$, then for every $a$ the family $\{f(y+a)\}$ must be still better than the family $\{g(y+a)\}$.
3.2. Definitions and the main result. By a transformation we mean a smooth (differentiable) function from real numbers into real numbers. By an appropriate transformation group $G$ we mean a finite-dimensional connected Lie group of transformations. By a family of functions we mean the set of functions that is obtained from a smooth (everywhere defined) non-constant function $f(y)$ by applying all the transformations from some appropriate transformation group $G$. Let us denote the set of all the families by $F$.

A pair of relations ( $<, \sim$ ) is called consistent [Kreinovich 1990, Kreinovich Kumar 1990] if it satisfies the following conditions: (1) if $a<b$ and $b<c$ then $a<c$; (2) $a \sim a$; (3) if $a \sim b$ then $b \sim a$; (4) if $a \sim b$ and $b \sim c$ then $a \sim c$; (5) if $a<b$ and $b \sim c$ then $a<c$; (6) if $a \sim b$ and $b<c$ then $a<c$; (7) if $a<b$ then $b<a$ or $a \sim b$ are impossible.

Assume a set $A$ is given. Its elements will be called alternatives. By an optimality criterion we mean a consistent pair $(<, \sim)$ of relations on the set $A$ of all alternatives. If $a>b$, we say that $a$ is better than $b$; if $a \sim b$, we say that the alternatives $a$ and $b$ are equivalent with respect to this criterion. We say that an alternative $a$ is optimal (or best) with respect to a criterion $(<, \sim)$ if for every other alternative $b$ either $a>b$ or $a \sim b$.

We say that a criterion is final if there exists an optimal alternative, and this optimal alternative is unique.
Comment. In the present section we consider optimality criteria on the set $F$ of all families.
By the result of adding $a$ to a function $f(y)$ we mean a function $\bar{f}(y)=f(y+a)$. By the result of adding $a$ to a family $F$ we mean the set of the functions that are obtained from $f \in F$ by adding $a$. This result will be denoted by $F+a$. We say that an optimality criterion on $F$ is shift-invariant if for every two families $F$ and $G$ and for every number $a$, the following two conditions are true:
i) if $F$ is better than $G$ in the sense of this criterion (i.e., $F>G$ ), then $F+a>G+a$;
ii) if $F$ is equivalent to $G$ in the sense of this criterion (i.e., $F \sim G$ ), then $F+a \sim G \div a$. Comment. As we have already remarked, the demands that the optimality criterion is final and shift-invariant are quite reasonable. The only problem with them is that at first glance they may seem rather weak. However, they are not, as the following Theorem shows:

By a logistic function we mean $s_{0}(y)=1 /(1+\exp (-y))$.

THEOREM 1. If a family $F$ is optimal in the sense of some optimality criterion that is final and shift-invariant, then every function $f$ from $F$ is equal either to $a+b s_{0}(K y+l)$ for some $a, b, K$ and $l$, or a linear function $a+b x$, or to $a+b \exp (K x)$ for some $a, b, K$.
Comments. 1. Logistic, hyperbolic-tangent, linear and exponential functions are really among the most popular [Kosko 1992].
2. We assumed that $f$ must be smooth. If we consider $f$ that can be not smooth in some points, then it is natural to assume that on the intervals on which $f$ is smooth, it must coincide with one of these functions. Such piecewise smooth functions have also been successfully used, the most popular are threshold functions that are obtained from the smooth ones by restricting their values to $[0, \infty)$ or $[0,1]$ [Kosko 1992].
(The proofs are given in Section 5).

## 4. OPTIMIZATION OF NEURAL NETWORKS: RELATED RESULTS

4.1, Scale-invariance instead of shift invariance. In the above text we assumed that the optimality criterion is shift-invariant. The same arguments can be used to motivate the demand that the optimality criterion is invariant with respect to scaling transformations $f(y) \rightarrow f(a y)$ for some $a>0$. Let us analyze the consequences of this demand.

Definition. By the result of rescaling a function $f(y)$ by a real number $a>0$ we mean a function $\bar{f}(y)=f(a y)$. By the result of rescaling a family $F$ by $a$ we mean the set of the functions, that are obtained from rescaling $f \in F$ by $a$. This result will be denoted by $a F$. We say that an optimality criterion on $F$ is scale-invariant if for every two families $F$ and $G$ and for every number $a$ the following two conditions are true:
i)' if $F$ is better than $G$ in the sense of this criterion (i.e., $F>G$ ), then $a F>a G$;
$\left.i^{\prime}\right)^{\prime}$ if $F$ is equivalent to $G$ in the sense of this criterion (i.e., $F \sim G$ ), then $a F \sim a G$.
THEOREM 2. If a family $F$ is optimal in the sense of some optimality criterion that is final and scale-invariant, then every function $f$ from $F$ is equal to $f(y)=\left(A+B y^{-\alpha}\right) /\left(C+D y^{-\alpha}\right)$ for some $A, B, C, D$ and $\alpha>0$.
Comments. 1. In particular, for $A=0, B=C=D=1$ and $\alpha=2$ we get the Cauchy function $f(y)=1 /\left(1-y^{2}\right)$, that is used in neural networks (see, e.g., [Hecht-Nielsen 1991]. If $B=0$ and $\alpha$ is an integer, $\alpha>1$, we get ratio-polynomial signal functions that have also been successfully used [Kosko 1992]. For $\alpha=1, B=0$ and $A=C=D=1$ this function equals to the expression $x /(1+x)$, which was analyzed in Munro 1986]. So this Theorem gives a list of possible optimal non- linear neurons that generalizes Cauchy and ratio-polynomial functions.
2. By comparing the results of Theorems 1 and 2 one can conclude that a scale-invariant criterion cannot be shift-invariant: indeed, in this case we could apply Theorem 2, so $f$ must be described by the above expression. But these functions are different from the functions from Theorem 1, and so due to Theorem 1 this criterion is not shift-invariant.

But what if we still want our criterion to be both shift- and scale-invariant? For standard neurons with non-linearity of the type $y=f\left(w_{1} x_{1}+\ldots+w_{n} x_{n}\right)$ it is impossible; in section 4.3 we'll show that it is possible for a more general type of neurons.
4.2. More general families of neurons. A natural way to define a finite-dimensional family of functions is to fix finitely many functions $f_{i}(y)$ and consider their arbitrary linear combinations $\sum_{i} C_{i} f_{i}(y)$.

Definition. Let's fix an integer $m$. By a basis we mean a set of $m$ smooth functions $f_{i}(y), i=1,2, \ldots, m$. By a $m$-dimensional family of functions we mean all functions of the type $f(y)=\sum_{i} C_{i} f_{i}(y)$ for some basis $f_{i}(y)$, where $C_{i}$ are arbitrary constants. The set of all $m$-dimensional families will be denoted by $F_{m}$.


Our definitions of the optimality criterion, final criterion and shift-invariant criterion can be applied to these families.
THEOREM 3. If an m-dimensional family $F$ is optimal in the sense of some optimality criterion that is final and shift-invariant, then every function from $F$ is equal to a linear combination of the functions of the type $y^{p} \exp (\alpha y) \sin (\beta y+\phi)$, where $p$ is a non-negative integer, $\alpha, \beta$ and $\phi$ are real numbers.
Comment. In particular, for $p=1, \alpha=\beta=0$ we get linear functions; for $p=\beta=0$ we get exponential functions; for $p=\alpha=\phi=0$ we get a sine function, that has been successfully used [Braham 1989, Braham Hamblen 1990].
4.3. What if we demand scale-invariance and shift-invariance of the optimality criterion? Neurons with different non-linearity. In all the above cases we considered neurons that transform the input signals $x_{1}, \ldots, x_{n}$ into $f\left(w_{1} x_{1}+\ldots+w_{n} x_{n}\right)$. This means that we added only one type of non-linearity: $y \rightarrow f(y)$ for some non-linear $f$. Let us consider neurons with the most general type of non-linearity $y_{1}, \ldots, y_{p} \rightarrow f\left(y_{1}, \ldots, y_{p}\right)$, where $f$ is an arbitrary non-linear function of $p$ variables. As in the above case, linear transformations are easy to implement, therefore we can consider neurons of the type

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x_{1}, \ldots, x_{n} \rightarrow f\left(w_{11} x_{1}+\ldots+w_{1 n} x_{n}, w_{21} x_{1}+\ldots+w_{2 n} x_{n}, \ldots, w_{p 1} x_{1}+\ldots+w_{p n} x_{n}\right)
$$

where $w_{i j}$ are weights.
Definition. Let's fix integers $m$ and $p$. By a basis we mean a set of $m$ smooth functions $f_{i}\left(y_{1}, \ldots, y_{p}\right), i=1,2, \ldots, m$. By a $m$-dimensional family of functions of $p$ variables, or $m, p$ family for short, we mean a family that is formed by all functions of the type $f\left(y_{1}, \ldots, y_{p}\right)=$ $\sum_{i} C_{i} f_{i}\left(y_{1}, \ldots, y_{p}\right)$ for some basis $f_{i}\left(y_{1}, \ldots, y_{p}\right)$, where $C_{i}$ are arbitrary constants. The set of all $m, p$-dimensional families will be denoted by $F_{m, p}$.
Comment. Since it is easy to implement arbitrary linear transformations, it is reasonable to demand (like we did above) that the relative quality of the family does not change if we apply a shift or a rescaling to all its functions. So we arrive at the following defnitions:

Definitions. Suppose that a vector $\vec{a}=\left(a_{1}, \ldots, a_{p}\right)$ is given. By the result of adding $\vec{a}$ to a function $f\left(y_{1}, \ldots, y_{p}\right)$ we mean a :unction $\vec{F}\left(y_{1}, \ldots, y_{p}\right)=f\left(y_{1}-a_{1}, \ldots y_{p} \div a_{p}\right)$. $3 y$ ine resuit of adding $\vec{a}$ to a tamily $F$ we mean the ser of the functions. that are obtained fom $f \equiv E$ by adding $\vec{a}$. This result will be denoted by $F+\bar{a}$.

3y the resuit of rescaling a function $f^{\prime}\left(y_{1}, \ldots, y_{p}\right)$ by a rector $\bar{a}=\left(a_{1}, \ldots\right)$ with $I_{i}>0$ we mean a iunction $f\left(y_{1}, \ldots, y_{p}\right)=;\left(a_{1} y_{1}, \ldots, a_{p} y_{p}\right)$. By the resuit of rescaiing a iamily $\vec{F}$ by $\vec{a}$ we mean the set of the functions, that are obtained from rescaling $f \in F$ by $\vec{a}$. This result will be denoted by $\vec{a} F$.

Now we can apply the definitions of the optimality criterion, final criterion, shift- and scaleinvariant criterion to $m$-dimensional families.
THEOREM 4. If an $m$, $p$-dimensional family $F$ is optimal in the sense of some optimality criterion that is final, shift- and scale-invariant, then every function $f$ from $F$ is equal to a polynomial of $y_{1}, \ldots, y_{p}$.
Comment. So if we consider neurons with a more complicated non-linearity, we get the GMBH network (see [Hecht-Nielsen 1991, Section 5.6.1]), which historically is the first commercially successful non-linear neuron.
4.4. A related question: how to modify the weights during the training? This question is related to the following problem: usually during training weights are changed linearly (like $w \rightarrow w+c$ for some constant $a$ ). However, sometimes some weights become so big that the
output of the corresponding neurons is close to 1 . These neurons àre called saturated (see, e.g., [Wasserman 1989, pp. 90-91]. Saturation extends the training time, bringing the whole training process into a paralysis. To overcome paralysis, non-linear weight transformations are used: $w \rightarrow g(w)$ for some non-linear $g$. The main purpose of this transformation is to get rid of big values, i.e., to transform the whole set of real numbers $(-\infty, \infty)$ into, some limited interval of values $[-\Delta, \Delta]$. But there are many functions with this property. So the natural question arises: what $g$ is the best to choose?

The same arguments as in Sections 2 and 3 can be used to conclude that what we really need to choose is a family of functions, not just a function $g$, and that it is reasonable to assume that this family is obtained from some function $g(w)$ by applying all the transformations from some appropriate transformation group.

If a function $g$ is better than a function $\bar{g}$, then it is reasonable to assume that it will still be better if we first make one more standard training step $w \rightarrow w+a$ and only then apply the non-linear transformation. These two consequent steps are equivalent to one transformation $w \rightarrow g(w+a)$. So the above demand means that if $g(w)$ is better than $\bar{g}(w)$, then $g_{1}(w)=g(w+a)$ must be better than $\bar{g}_{1}(w)=\bar{g}(w+a)$. In other words, the optimality criterion must be shiftinvariant. Let's turn to formal definitions.

Definitions. Let us fix some real number $\Delta>0$. We say that a function is bounded if its values always belong to $[-\Delta, \Delta]$. By a family of weight transformations we mean the set of functions that is obtained from a smooth (everywhere defined) non-constant bounded function $f(w)$ by applying all the transformations from some appropriate transformation group $G$. Let us denote the set of all such families by $F_{w}$. Let us consider optimality. criteria on this set $F_{w}$.

THEOREM 5. If a family $F_{w}$ is optimal in the sense of some optimality criterion that is final and shift-invariant, then every bounded function $g$ from $F$ is equal to $a+b s_{0}(K y+l)$ for some $a, b, K$ and $l$.

Comment. Such functions really proved to be the best in overcoming paralysis [Wasserman 1989, pp. 90-91].

## 5. PROOFS.

Proof of Theorem 1. The idea of this proof is as ioilows: irst we prove that the appropriate transformation group consists of tractionally-inear functions (in part i), then we prove that the optimal family is shift-invariant in part 2), and from that in part 3 we conclude that any function : irom $\vec{F}$ sacisîes some iunc:ional equations, whose solutions are inown.

1. By an appropriate group we meant a connected finite-dimensional Lie group of iransiorpations of the set of real numbers $R$ onto itself that contains all linear transformations. Norbert Wiener asked to classify such groups for an $n$-dimensional space with arbitrary $n$, and $t$ is classification was obtained in [Guillemin Sternberg 1964] and [Singer Sternberg 1965]. In our case (when $n=1$ ) the only possible groups are the group of all linear transformations and the group of all fractionally-linear transformations $x \rightarrow(a x+b) /(c x+d)$. In both cases the group consists only of fractionally linear transformations (the simplifed proof for the 1 -dimensional case is given in (Kreinovich 1987]; for other applications of this result see [Kreinovich Kumar 1990, 1991], [Kreinovich Corbin 1991], [Kreinovich Quintana 1991]).
2. Let us now prove that the optimal family $F_{\text {opt }}$ exists and is shift-invariant in the sense that $F_{o p t}=F_{o p t}+a$ for all real numbers $a$. Indeed, we assumed that the optimality criterion is final, therefore there exists a unique optimal family $F_{\text {opt }}$. Let's now prove that this optimal family is shift-invariant (this proof is practically the same as in [Kreinovich 1990]). The fact that $F_{o p t}$ is optimal means that for every other $F$, either $F_{o p t}>F$ or $F_{o p t} \sim F$. If $F_{o p t} \sim F$ for some $F \neq F_{\text {opt }}$, then from the definition of the optimality criterion we can easily deduce that $F$ is also


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optimal, which contradicts the fact that there is only one optimal family. So for every $F$ either $\bar{F}_{\text {opt }}>F$ or $F_{\text {opt }}=F$.

Take an arbitrary $a$ and let $F=F_{o p t}+a$. If $F_{o p t}>F=F_{o p t}+a$, then from the invariance of the optimality criterion (condition ii) we conclude that $F_{\text {opt }}-a>F_{\text {opt }}$, and that conclusion contradicts the choice of $F_{\text {opt }}$ as the optimal family. So $F_{o p t}>F=F_{o p t}+a$ is impossible, and therefore $F_{o p t}=F=F_{o p t}+a$, i.e., the optimal family is really shift- invariant.
3. Let us now deduce the actual form of the functions $f$ from the optimal family. If $f(y)$ is such a function, then the result $f(y+a)$ of adding $a$ to this function $f$ belongs to $F+a$, and so, due to 2 ., it belongs to $F$. But all the functions from $f$ can be obtained from each other by fractionally linear transformations, so $f(y+a)=(A+B f(y)) /(C+D f(y))$ for some $A, B, C$ and $D$. So we arrive at a functional equation for $f$. Let us reduce this equation to a one with a known solution. For that purpose, let us use the fact that fractionally linear transformations are projective transformations of a line, and for such transformations the cross ratio is preserved [Aczel 1966, Section 2.3], i.e., if $g(y)=(A+B f(y)) /(C+D f(y))$, then

$$
\frac{g\left(y_{1}\right)-g\left(y_{3}\right)}{g\left(y_{2}\right)-g\left(y_{3}\right)} \frac{g\left(y_{2}\right)-g\left(y_{4}\right)}{g\left(y_{1}\right)-g\left(y_{4}\right)}=\frac{f\left(y_{1}\right)-f\left(y_{3}\right)}{f\left(y_{2}\right)-f\left(y_{3}\right)} \frac{f\left(y_{2}\right)-f\left(y_{4}\right)}{f\left(y_{1}\right)-f\left(y_{4}\right)}
$$

for all $y_{i}$. In our case this is true for $g(y)=f(y+a)$, therefore for all $a$ the following equality is true:

$$
\frac{f\left(y_{1}+a\right)-f\left(y_{3}+a\right)}{f\left(y_{2}+a\right)-f\left(y_{3}+a\right)} \frac{f\left(y_{2}+a\right)-f\left(y_{4}+a\right)}{f\left(y_{1}+a\right)-f\left(y_{4}+a\right)}=\frac{f\left(y_{1}\right)-f\left(y_{3}\right)}{f\left(y_{2}\right)-f\left(y_{3}\right)} \frac{f\left(y_{2}\right)-f\left(y_{4}\right)}{f\left(y_{1}\right)-f\left(y_{4}\right)}
$$

The most general continuous solutions of this functional equation are given by Theorem 2.3.2 from [Aczel 1966]: either $f$ is fractionally linear, or $f(y)=(a+b \tan (k y)) /(c+d \tan (k y))$ for some $a, b, c, d$, or $f(y)=(a+b \tanh (k y)) /(c+d \tanh (k y))$, where $\tanh (z)=\sinh (z) / \cosh (z)$, $\sinh (z)=(\exp (z)-\exp (-z)) / 2)$ and $\cosh (z)=(\exp (z)+\exp (-z)) / 2)$.

If $f(y)$ is fractionally linear $f(y)=(a+b y) /(c+d y)$ and $d \neq 0$, then the denominator is equal to zero for $y=-c / d$. The only way for the function to be defined for this $y$ is that the numerator should also be zero, i.e., $a+b y=a+b(-c / d)=0$. But in this case $a=b(c / d)$, therefore $a+b y=b(c / d+y)=(b / d)(c+d y)$, and the fraction $f(y)$ is always equal to a constant $b / d$. But we assumed that $f$ is not a constant. So $d=0$ and $f$ is linear.

Let us prove that the expressions with tangent are impossible. Indeed, the denominator must be not identically equal to zero, therefore either $c \neq 0$, or $d \neq 0$. If $d \neq 0$, then for $k y=\arctan (-c / d)$ we have $\tan (k y)=-c / d$, and the denominator is equal to zero. As in the linear case we can then conclude that in this case $f$ is constant, and that contradicts to our assumption that it is not. So $d=0$ and $f(y)=(a / d)+(b / d) \tan (k y)$. Hence either $b=0$ and $f=$ const, or $b \neq 0$, and $f$ is not defined, when $\tan (k y)=\infty$, i.e., when $k y=\pi / 2$ and $y=\pi /(2 k)$. So expressions with tangent are really impossible.

Let us now consider the case of hyperbolic tangent. If $k=0$, then $f$ is constant, which is impossible. So $k \neq 0$. If $k<0$, then we can take $\bar{k}=-k$ and use the fact that tanh is an odd function, so $\tanh (k y)=-\tanh (\bar{k} y)$. Therefore, in the following we can assume that $k>0$. Multiplying both the denominator and the numerator by $\cosh (z)$, we conclude that $f(y)=(a \cosh (k y)+b \sinh (k y)) /(c \cosh (k y)+d \sinh (k y))$. We then substitute the expressions for $\sinh$ and cosh in terms of exp, and conclude that $f(y)=(A \exp (k y)+B \exp (-k y)) /(C \exp (k y)+$ $D \exp (-k y)$ ) for some $A, B, C, D$. Multiplying both denominator and numerator by $\exp (-k y)$, we arrive at $f(y)=(A+B \exp (-2 k y)) /(C+D \exp (-2 k y))$. If $D=0$, then we get a linear transformation of the exponential function. If $C=0$, then $f(y)=(B / D)+(A / D) \exp (2 k y)$, which is also a linear transformation of the exponential function. Let us now consider the case, when both $C$ and $D$ are different from 0 .

If $C$ and $D$ have different signs, then for $\exp (2 k y)=-D / C$ the denominator equals to zero, and so, just like in the tangent case, we conclude that $f$ is either identically constant, or not defined in this point $y=\ln (-D / C) /(2 k)$. If $C$ and $D$ have the same signs, then for $l=-\ln (D / C)$ we have $C+D \exp (-2 k y)=C(1+(D / C) \exp (-2 k y))=C(1+\exp (-(2 k y+l))$. If we substitute $\exp (-2 k y)=\exp (-(2 k y+l)) \exp (l)=(C / D) \exp (-(2 k y+l))$ into the numerator, we get $A+(B C / D) \exp (-(2 k y+l))$, and therefore $f(y)=(A+(B C / D) \exp (-(2 k y+l))) /(C(1+$ $\exp (-(2 k y+l))$ ). One can check (by substituting the expression of the logistic function $s_{0}$ in terms of exp) that this expression is equal to $(A / C)+(B / D-A / C) s_{0}(2 k y+l)$. So we get the desired expression for $K=2 k$. Q.E.D.
Proof of Theorem 2. Just like in the proof of Theorem 1, we conclude that $f(a y)=(A+$ $B f(y)) /(C+D f(y))$. This functional equation is almost the same as the one we solved in Theorem 1, with the only exception being that here we have a product instead of a sum. It is well known that if we turn to logarithms, then the logarithm of a product is equal to the sum of logarithms. So in order to reduce this case to the one already analyzed, let us introduce the new variable $Y=\ln y$ (so that $y=\exp Y$ ), and a new function $F(Y)=f(\exp (y))$. For this function, the above functional equation takes the following form: for every $E$ there exist $A, B, C, D$ such that $F(Y+E)=(A+B F(Y)) /(C+D F(Y))(E=\ln a)$. In the proof of Theorem 1, we have already enumerated the solutions of this equation, so $F(Y)$ is either a fractionally-linear function, or a fractionally-linear transformation of $\tan (k y)$ or $\tanh (k y)$. If we know $F$, then using the equation $F(Y)=f(\exp (y))$, we can reconstruct $f(y)=F(\ln y)$ for $y>0$. Similar expressions can be obtained for $y<0$ : in this case we need to use $Y=\ln |y|$. So in order to complete the proof, we must substitute $\ln y$ into the expressions enumerated above, and choose those that are defined everywhere.

Let us first consider the case when $F(Y)=(a+b Y) /(c+d Y)$. Substituting $\ln y$ instead of $Y$, we get $f(y)=(a+b \ln y) /(c+d \ln y)$. This function must be smooth. Let us compute the derivative $f^{\prime}: f^{\prime}(y)=((b / y)(c+d \ln y)-(d / y)(a+b \ln y)) /(c+d \ln y)^{2}$. If $b / c-a / d=0$, then, as in Theorem 1, we can conclude that $f$ is identically constant. If $b / c-a / d \neq 0$, then for $y \rightarrow 0$ this derivative tends to $\infty$, so such functions $f$ are not smooth at 0 .

The fact that the expressions with tangent are impossible is proved just like in Theorem 1. So the only remaining case is the case of tanh, in which the function $F(Y)$ can be reduced to $F(Y)=(A+B \exp (-2 k Y)) /(C+D \exp (-2 k Y))$. Substituting $Y=\ln y$, and using the fact that $\exp (-2 k \ln y))=(\exp (\ln y))^{-2 k}=y^{-2 k}$, we conclude that $f(y)=\left(A+B y^{-\alpha}\right) /\left(C+D y^{-\alpha}\right)$, where $\alpha=2 k$. Q.E.D.
Proof of Theorem 3. As in the proof of Theorem 1, we come to a conclusion that the optimal family $F_{\text {opt }}$ exists and is shift-invariant. In particular, for every $i$ the result $f_{1}(y+a)$ of shifting $f_{i}(y)$ must belong to the same family, i.e., $f_{i}(y+a)=C_{i 1}(a) f_{1}(y)+C_{i 2}(a) f_{2}(y)+\ldots+C_{i m}(a) f_{m}(y)$ for some constants $C_{i j}$, depending on $a$. Let us first prove that these functions $C_{i j}(a)$ are differentiable. Indeed, if we take $m$ different values $y_{k}, 1 \leq k \leq m$, we get $m$ linear equations for $C_{i j}(a)$ :

$$
f_{i}\left(y_{k}+a\right)=C_{i 1}(a) f_{1}\left(y_{k}\right)+C_{i 2}(a) f_{2}\left(y_{k}\right)+\ldots+C_{i m}(a) f_{m}\left(y_{k}\right)
$$

from which we can determine $C_{i j}$ using Kramer's rule. Kramer's rule expresses every unknown as a fraction of two determinants, and these determinants polynomially depend on the coefficients. The coefficients either do not depend on $a$ at all $\left(f_{j}\left(y_{k}\right)\right)$ or depend smoothly $\left(f_{i}\left(y_{k}+a\right)\right)$ because $f_{i}$ are smooth functions. Therefore these polynomials are also smooth functions, and so is their fraction $C_{i j}(a)$.

We have an explicit expression for $f_{i}(y+a)$ in terms of $f_{j}(y)$ and $C_{i j}$. So, when $a=0$, the derivative of $f_{i}(y+a)$ with respect to $a$ equals to the derivative of the expression. If we differentiate it, we get the following formula: $f_{k}^{\prime}(y)=c_{i 1} f_{1}(y)+c_{i 2} f_{2}(y)+\ldots+c_{i m} f_{m}(y)$, where $c_{i j}=C_{i j}^{\prime}(0)$. So the set of functions $f_{i}(y)$ satisfies the system of linear differential equations with
coustant coefficients. The general solution of such system is well known [Bellman 1970], so we get the desired expressions. Q.E.D.
Proof of Theorem 4. Just like in Theorem 1, we conclude that an optimal family exists and is both shift- and scale-invariant. This means that the results of adding $\vec{a}$ to $f_{i}$ and rescaling $f_{i}$ by $\vec{a}$ also belong to this optimal family $F$.

For shift-invariance this means that

$$
f_{i}\left(y_{1}+a_{1}, y_{2}+a_{2}, \ldots, y_{p}+a_{p}\right)=C_{i 1}(\vec{a}) f_{1}\left(y_{1}, y_{2}, \ldots, y_{p}\right)+\ldots+C_{i m}(\vec{a}) f_{m}\left(y_{1}, \ldots, y_{p}\right)
$$

In particular, if we take $a_{2}=\ldots=a_{p}=0$, we conclude that

$$
f_{i}\left(y_{1}+a_{1}, y_{2}, \ldots, y_{p}\right)=C_{i 1}\left(a_{1}\right) f_{1}\left(y_{1}, y_{2}, \ldots, y_{p}\right)+C_{i 2}\left(a_{1}\right) f_{2}\left(y_{1}, \ldots, y_{p}\right)+\ldots+C_{i m}\left(a_{1}\right) f_{m}\left(y_{1}, \ldots, y_{p}\right)
$$

If we fix some values $y_{2}, \ldots, y_{p}$ and denote $g_{i}\left(y_{1}\right)=f_{i}\left(y_{1}, y_{2}, \ldots, y_{p}\right)$, we conclude that $g_{i}\left(y_{1}+a_{1}\right)=$ $C_{i 1}\left(a_{1}\right) g_{1}\left(y_{1}\right)+\ldots+C_{i m}\left(a_{1}\right) g_{m}\left(y_{1}\right)$. So the functions $g_{i}$ satisfy the same equations that we have already solved in the proof of Theorem 3, and therefore each of $g_{i}$ is equal to the linear combination of the functions $y_{1}^{p} \exp \left(\alpha y_{1}\right) \sin \left(\beta y_{1}+\phi\right)$ from the formulation of Theorem 3 .

Likewise scale-invariance means that
$f_{i}\left(a_{1} y_{1}, a_{2} y_{2}, \ldots, a_{p} y_{p}\right)=C_{i 1}(\vec{a}) f_{1}\left(y_{1}, y_{2}, \ldots, y_{p}\right)+C_{i 2}(\vec{a}) f_{2}\left(y_{1}, \ldots, y_{p}\right)+\ldots+C_{i m}(\vec{a}) f_{m}\left(y_{1}, \ldots, y_{p}\right)$.
If we take $a_{2}=\ldots=a_{p}=1$, we conclude that $g_{i}\left(a_{1} y_{1}\right)=C_{i 1}\left(a_{1}\right) g_{1}\left(y_{1}\right)+\ldots+C_{i m}\left(a_{1}\right) g_{m}\left(y_{1}\right)$. This functional equation is almost the same as for shift-invariance, the only difference is that we have a product instead of a sum. We already had such a situation, when we proved Theorem 2, and so we know what trick to apply: we must introduce a new variable $Y=\ln y_{1}$ (so that $y_{1}=\exp (Y)$ ), and new functions $G_{i}(Y)=g_{i}(\exp (Y))$ (so that $g_{i}\left(y_{1}\right)=G_{i}\left(\ln y_{1}\right)$ ). Then for these new functions this functional equation takes the form $G_{i}(A+Y)=\bar{C}_{i 1}(A) G_{1}(Y)+\ldots+\bar{C}_{i m}(A) G_{m}(Y)$. This is precisely the system of functional equations that we already know how to solve. So we can conclude that $G_{i}(Y)$ is a linear combination of these functions $Y^{p} \exp (\alpha Y) \sin (\beta Y+\phi)$ from Theorem 3. When we substitute $Y=\ln y_{1}$, we conclude that $g_{i}\left(y_{1}\right)=G_{i}(Y)=G_{i}\left(\ln y_{1}\right)$ is a linear combination of the functions $\left(\ln y_{1}\right)^{p} \exp \left(\alpha \ln y_{1}\right) \sin \left(\beta \ln y_{1}+\phi\right)$. This expression is rather complicated. The only simplincation that we can apply is :o change exp( $\alpha \ln y_{1}$ ) o (expiln $y_{1}$ ) $)^{\alpha}=y_{1}^{x}$. so we conclude that $g_{i}$ is a linear combination of the functions ( $\left.\ln y_{1}\right)^{p} y_{1}^{\alpha} \sin \left(3 \ln y_{1} \div 0\right)$.

So Sor the same functions $g_{i}$ we have ;wo different expressions odrained from the demancis of shift-invariance and scale-invariance. When can a function $y_{i}$ satisiy both conciusions, !.e.. jeiong to both ciasses? If it sontains ierms with logarithms, it cannot be a inear combination of the functions from Theorem 3, because there are no logarithms among them. The same if it contains sines of logarithms. So the only case when a linear combination of the functions $\left(\ln y_{1}\right)^{p} y_{1}^{\alpha} \sin \left(\beta \ln y_{1}+\phi\right)$ is at the same time the linear combination of the functions $y_{1}^{\beta} \exp \left(\bar{\alpha} y_{1}\right) \sin \left(\bar{\beta} y_{1}+\bar{\phi}\right)$ is when $p=\beta=0$. In this case the above expression turns into $y_{1}^{\alpha}$. and from the equality of these expressions we conclude that $\alpha=\bar{p}$. But $\bar{p}$ is necessarily a non-negative integer, and therefore $\alpha$ is non-negative integer as well. So $g_{i}\left(y_{1}\right)$, which is equai

Similarly we can conclude that $f_{i}$ polynomially depends on all other variabies $y_{2}, \ldots, y_{p}$, and therefore all the functions $f_{i}$ are polynomials of $y_{i}$. Every function from $F$ is a linear combination of these polynomials, and therefore a polynomial itself. Q.E.D.
Proof of Theorem 5. We can repeat the proof of Theorem 1 and come to a conclusion that all these functions are either linear, or exponential, or a logistic function. Neither linear, nor exponential functions are bounded, so only a logistic function is left. Q.E.D.

Acknowledgments. This research was supported by NSF grant No. CDA-9015006, NASA Research grant NAG $9-482$ and a grant from the Institute for Manufacturing and Materials Management. The authors are also thankful to Dr. R. Hecht-Nielsen (San Diego, CA) for encouragement, to Dr. G. E. Hinton (Toronto) for interesting preprints, and to all participants of the NSF II Workshop (Purdue, 1991), especially to Dr. Lokendra Shastri (Philadelphia) for valuable discussions.

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