

NASA Contractor Report 198158

ICASE Report No. 95-37

1N-34
49926
p. 30



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(NASA-CR-198158) SPECTRAL METHODS
ON ARBITRARY GRIDS Final Report
(ICASE) 30 p

N95-27247

Unclass

G3/34 0049926

Contract No. NAS1-19480
May 1995

Institute for Computer Applications in Science and Engineering
NASA Langley Research Center
Hampton, VA 23681-0001



Operated by Universities Space Research Association

SPECTRAL METHODS ON ARBITRARY GRIDS*

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Abstract

Stable and spectrally accurate numerical methods are constructed on arbitrary grids for partial differential equations. These new methods are equivalent to conventional spectral methods but do not rely on specific grid distributions. Specifically, we show how to implement Legendre Galerkin, Legendre collocation, and Laguerre Galerkin methodology on arbitrary grids.

*Research supported by AFOSR grant no. 93-0090, ARPA grant no. N00014-91-J-4016, and NSF grant no. DMS-9211820. The research was also partially supported by the National Aeronautics and Space Administration under NASA Contract No. NAS1-19480 while the second author was in residence at the Institute for Computer Applications in Science and Engineering (ICASE), NASA Langley Research Center, Hampton, VA 23681-0001.

1 Introduction

Conventional spectral methods impose rigid requirements on the computational grids used. The grid points are the nodes of Gauss-like quadrature formulas (Gauss, Gauss Radau, or Gauss Lobatto (GL) formulas). These nodes are denser at the boundaries than in the middle of the domain. Although this property is suitable for boundary-layer problems, it may create difficulties for other types of problems, particularly those with disparate length scales that occur in multiple regions of the domain (e.g., diffusive burning or detonation and reacting mixing layers). The principle reason for the degradation in performance on these disparate problems is that the predetermined node points do not, in general, coincide with the features that are being resolved. Extensive mappings can concentrate the node points into regions more ideally suited for accurate resolution but present a serious limitation for complicated problems. For this reason, spectral multidomain techniques have an obvious advantage for complicated problems [1]–[3].

Another complication that conventional spectral methods have, is their implementation in complex geometries. Meshes that are predetermined present a significant constraint. Flexible mesh distributions are easily extended to geometries that are not tensor products of straight lines (to be shown in a later work).

Spectral methods that are not constrained to specific nodal points would clearly be more flexible than conventional spectral methods. Specifically, a distribution of points that more closely approximates the disparate features in the domain could be adopted from the outset. Subsequent adaptation to solution features in the domain need not rely on smooth mappings. In addition, these “arbitrary-grid spectral techniques” could be used in conjunction with multidomain ideas. We focus on formalizing these ideas within the context of spectral techniques.

In this paper, we present some ideas for constructing spectral methods with arbitrary grids. We demonstrate these ideas for a case of spectral solutions of hyperbolic equations; however, these ideas can be applied to any partial differential equation. To illustrate the basic idea, consider the following hyperbolic system of equations in conservation form:

$$\frac{\partial U}{\partial t} = \frac{\partial F(U)}{\partial x} \quad -1 \leq x \leq 1 \quad (1)$$

with arbitrary initial and boundary conditions. For spectral methods, a polynomial (in the spatial variable x) of degree N , $U_N(x, t)$, and a projection operator I_N are sought such that

$$I_N \left[\frac{\partial U_N}{\partial t} - \frac{\partial I_N F(U_N)}{\partial x} \right] = 0 \quad (2)$$

Of the spectral techniques, the most popular method is the Chebyshev collocation method, in which $I_N f(x)$ collocates $f(x)$ at the Chebyshev GL points $\xi_j = \cos(\frac{\pi j}{N})$. Note that we

have here two projections; one involves the differentiation of $F(U_N)$, and the other involves the way that the equation is satisfied. Thus, the first application of the operator I_N occurs when we approximate $\frac{\partial F(U)}{\partial x}$ by the derivative of the interpolation polynomial that interpolates $F(U)$ at the Chebyshev GL quadrature nodes. The second application of I_N occurs when we satisfy the approximate equation

$$\left[\frac{\partial U_N}{\partial t} - \frac{\partial I_N F(U_N)}{\partial x} \right] = 0$$

at the Chebyshev GL points.

The basic premise for unstructured spectral methods is that equation (2) does not have to be satisfied in the same manner in which the operation $\frac{\partial I_N F(U_N)}{\partial x}$ is carried out. In particular, the derivative operation can be carried out by interpolation at any selected points; the equation is satisfied by either a Galerkin formulation or by a collocation method at a different set of points. Most importantly, the equation must be satisfied correctly.

Mathematically speaking, we can replace equation (2) with

$$I_N \left[\frac{\partial U_N}{\partial t} - \frac{\partial J_N F(U_N)}{\partial x} \right] = 0 \quad (3)$$

where $I_N \neq J_N$.

In reference [4], a particular case with this approach has been discussed. The operator J_N was defined by the Chebyshev collocation operator, and I_N was the Legendre collocation operator. In the constant-coefficient case ($F(U) = U$), this method reduces to the Legendre collocation method with an efficient way of calculating the derivative by using Chebyshev collocation points. We now generalize this notion to an arbitrary set of points, which enables us to apply spectral methods in circumstances for which the grid points are *not* nodes of some Gauss quadrature formula.

The method discussed in this paper is different from using a transformation to redistribute the grid points. The use of a transformation to redistribute the grid points involves approximation of the solution by a polynomial in the transformed variable; as a result, the approximation is not a polynomial in the original variable. Our method utilizes a polynomial in the original variable. Moreover, the new method can be applied to totally unstructured grids.

Finally, it should be noted that the new method has many similarities with spectral elements, although the method of derivation is different. For instance, Patera [5] or Korcza et. al [1] used global polynomials (Lagrangian interpolants), passed through the Chebyshev collocation points, to obtain spectral elements. However, their work was not generalized to arbitrary grids.

2 The Differentiation Matrix for Unstructured Grids

Consider the set of points $(x_0 = 1, x_1, x_2, \dots, x_{N-1}, x_N = -1)$, where the points x_1, x_2, \dots, x_{N-1} are arbitrary. Let $f(x)$ be a function defined everywhere in $[-1, 1]$. The interpolation polynomial $f_N(x)$ that collocates $f(x)$ at the points x_j is given by

$$f_N(x) = J_N f = \sum_{j=0}^N f(x_j) L_j(x) \quad (4)$$

where the Lagrange polynomials $L_j(x)$ are defined by

$$L(x) = (x - x_1)(x - x_2) \dots (x - x_{N-2})(x - x_{N-1}) \quad (5)$$

$$L_j(x) = \frac{(1 - x^2)L(x)}{(1 - x_j^2)(x - x_j)L'(x_j)} \quad 1 \leq j \leq N - 1 \quad (6)$$

$$L_0(x) = \frac{(1 + x)L(x)}{2L(1)} \quad (7)$$

$$L_N(x) = \frac{(1 - x)L(x)}{2L(-1)} \quad (8)$$

The Lagrange polynomial evaluated at the discrete points x_k for $k \neq j$, is equal to 0; $L_j(x_k) = \delta_{j,k}$.

We use $\frac{dJ_N f(x)}{dx}$ as the approximation to $\frac{df(x)}{dx}$. Note that $\frac{df_N}{dx}$ has two alternative representations; the first is obtained by differentiating (4) as

$$\frac{df_N(x)}{dx} = \sum_{j=0}^N f(x_j) L'_j(x) \quad (9)$$

The second representation stems from the fact that $\frac{dJ_N f(x)}{dx}$ is a polynomial of degree $N - 1$; therefore,

$$\frac{df_N(x)}{dx} = \sum_{j=0}^N f'_N(x_j) L_j(x) \quad (10)$$

Equations (9) and (10) are used to relate the grid-point values of the derivative $f'_N(x_j)$ to those of the function. The most obvious way is to equate the expressions in (9) and (10) at the grid points x_k ($0 \leq k \leq N$) to obtain:

$$f'_N(x_k) = \sum_{j=0}^N L'_j(x_k) f(x_j) \quad (11)$$

To rewrite expression (11) in matrix form, we first denote

$$\vec{f}' = [f'_N(x_0), \dots, f'_N(x_N)]^T, \quad \vec{f} = [f(x_0), \dots, f(x_N)]^T$$

which yields

$$\vec{f}' = D\vec{f} \quad (12)$$

where the differentiation matrix D is given by

$$D = (d_{j,k}) = [L'_k(x_j)] \quad (13)$$

Another method for expressing the equivalency between (9) and (10) is to state that the difference between these expressions (which is identically 0) is orthogonal to all polynomials of degree $\leq N$:

$$\int_{-1}^1 \sum_{j=0}^N [f(x_j)L'_j(x) - f'_N(x_j)L_j(x)] L_k(x) dx = 0 \quad 0 \leq k \leq N \quad (14)$$

The system of equations that follows from (14) can be rewritten as

$$\sum_{j=0}^N m_{k,j} f'_N(x_j) = \sum_{j=0}^N s_{k,j} f(x_j) \quad 0 \leq k \leq N \quad (15)$$

where

$$m_{k,j} = \int_{-1}^1 L_j(x) L_k(x) dx \quad (16)$$

and

$$s_{k,j} = \int_{-1}^1 L'_j(x) L_k(x) dx \quad (17)$$

In the matrix form, equation (15) becomes

$$\mathbf{M}\vec{f}' = \mathbf{S}\vec{f} \quad (18)$$

where

$$\mathbf{M} = (m_{k,j}) \quad 0 \leq j, k \leq N \quad (19)$$

and

$$\mathbf{S} = (s_{k,j}) \quad 0 \leq j, k \leq N \quad (20)$$

Equations (14) and (11) are different manifestations of the same fact: (9) and (10) are equivalent. Therefore, the differentiation matrices derived from (14) must be the same as the matrix derived from (11) (with the assumption that \mathbf{M} is invertible):

$$D = \mathbf{M}^{-1}\mathbf{S} \quad (21)$$

To prove this directly, we show that

$$\mathbf{M}\mathbf{D} = \mathbf{S} \quad (22)$$

By writing the (i, k) term on the left-hand side of (22), we obtain

$$(\mathbf{M}\mathbf{D})_{i,k} = \sum_{j=0}^N m_{i,j} d_{j,k}$$

If we substitute (13) and (16) into (2), then we get

$$(\mathbf{M}\mathbf{D})_{i,k} = \int_{-1}^1 L_i(x) \left[\sum_{j=0}^N L_j(x) L'_k(x_j) \right]$$

We now use the fact that every polynomial of degree N is *identical with its N -degree interpolation polynomial*. Thus, because $L'_k(x)$ is a polynomial of degree $N - 1$ and

$$\sum_{j=0}^N L'_k(x_j) L_j(x)$$

is its interpolant at the points x_j ($0 \leq j \leq N$) then

$$\sum_{j=0}^N L_j(x) L'_k(x_j) = L'_k(x)$$

which yields

$$(\mathbf{M}\mathbf{D})_{i,k} = \int_{-1}^1 L_i(x) L'_k(x) = s_{i,k}$$

(which is apparent from (17)). This establishes expression (22).

□

Thus, we have defined a new method, based on the arbitrary distribution of points, to approximate the derivative of a function. The attractive features of the representation (21) of the differentiation matrix are summarized in lemma 3.1 and lemma 3.2:

Lemma 3.1:

The matrix \mathbf{M} defined in (16) is a symmetric positive-definite matrix.

Proof:

The fact that \mathbf{M} is symmetric follows immediately from the definition (16). In fact,

$$m_{k,j} = \int_{-1}^1 L_j(x) L_k(x) dx = m_{j,k}$$

We must show that \mathbf{M} is positive definite. Let \vec{V} be an $N + 1$ component vector:

$$\vec{V} = (v_0, \dots, v_N)$$

Then,

$$\vec{V}^T \mathbf{M} \vec{V} = \sum_{i=0}^N \sum_{j=0}^N m_{i,j} v_i v_j$$

Recall the definition of $m_{i,j}$ from (16). We get

$$\vec{V}^T \mathbf{M} \vec{V} = \int_{-1}^1 \sum_{i=0}^N v_i L_i(x) \sum_{j=0}^N v_j L_j(x) dx \geq 0 \quad (23)$$

Clearly, the equality sign holds only if \vec{V} is the null vector.

□

Equation (23) can be interpreted in a different way. Let $v(x)$ be the polynomial of degree N defined by

$$v(x_j) = v_j \quad 0 \leq j \leq N$$

so that

$$v(x) = \sum_{j=0}^N v_j L_j(x)$$

Then, (23) can be rewritten as

$$\vec{V}^T \mathbf{M} \vec{V} = \int_{-1}^1 v(x)^2 dx \quad (24)$$

Thus, every vector \vec{V} can be identified with a polynomial $v(x)$ that takes the values of its components at the grid points x_j . The vector space norm

$$\vec{V}^T \mathbf{M} \vec{V}$$

is equivalent to the function space norm

$$\int_{-1}^1 v(x)^2 dx$$

Next, we will consider the properties of the matrix \mathbf{S} .

Lemma 3.2:

Let \mathbf{S} be defined in (17), and let \vec{V} be defined as before. Then,

$$\vec{V}^T \mathbf{S} \vec{V} = \frac{1}{2}(v_0^2 - v_N^2) \quad (25)$$

Proof:

We start by showing that \mathbf{S} is almost antisymmetric. From the definition (17)

$$s_{k,j} = \int_{-1}^1 L'_j(x) L_k(x) dx$$

and integration by parts, we get

$$s_{k,j} = L_j(1)L_k(1) - L_j(-1)L_k(-1) - s_{j,k}$$

We now use the definition of the Lagrange polynomials (6)–(8) and note that the boundary terms vanish for $0 \neq j, k \neq N$ to yield

$$s_{k,j} + s_{j,k} = \delta_{k,0}\delta_{j,0} - \delta_{k,N}\delta_{j,N}$$

Thus,

$$\begin{aligned} \vec{V}^T \mathbf{S} \vec{V} &= \sum_{k=0}^N \sum_{j=0}^N v_j v_k s_{k,j} \\ &= \frac{1}{2} \sum_{k=0}^N \sum_{j=0}^N v_j v_k (s_{k,j} + s_{j,k}) \\ &= \frac{1}{2} \sum_{k=0}^N \sum_{j=0}^N v_j v_k (\delta_{k,0}\delta_{j,0} - \delta_{k,N}\delta_{j,N}) \\ &= \frac{1}{2}(v_0^2 - v_N^2) \end{aligned}$$

which completes the proof of (25).

□

As before, equation (25) has a natural interpretation in the polynomial space. Let $v(x)$ be the polynomial of degree N such that $v(x_j) = v_j$. Then,

$$\begin{aligned}
\vec{V}^T \mathbf{S} \vec{V} &= \sum_{k=0}^N \sum_{j=0}^N v_j v_k s_{k,j} \\
&= \sum_{k=0}^N \sum_{j=0}^N v_j v_k \int_{-1}^1 L_k(x) L_j'(x) dx \\
&= \int_{-1}^1 v(x) v'(x) dx \\
&= \frac{1}{2} [v^2(1) - v^2(-1)]
\end{aligned}$$

Note that

$$v(1) = v_0, \quad v(-1) = v_N$$

Thus, (25) is an integration-by-parts formula.

The last issue that we will discuss in this section is the relationship between differentiation matrices, based on different grid-point distributions. Consider two grids x_j and y_j ($j = 0, \dots, N$). Let the Lagrange polynomial $L_j^x(x)$ be defined as in (6)–(8), and let $L_j^y(x)$ be defined in a similar way, based on the set of points y_j . This defines two differentiation matrices (see (11)):

$$D_x = (d_{j,k}^x) = [(L_k^x)'(x_j)] \quad (26)$$

and

$$D_y = (d_{j,k}^y) = [(L_k^y)'(y_j)] \quad (27)$$

We now show that the two matrices are similar.

Theorem 3.1:

Define the matrix T by

$$T = (t_{ij}) = [L_j^x(y_i)] \quad (28)$$

Then define

$$(T^{-1})_{ij} = [L_j^y(x_i)] \quad (29)$$

and

$$D_y = T D_x T^{-1} \quad (30)$$

Proof:

1. Because L_k^y is a polynomial of degree N ,

$$\sum_{j=0}^N L_j^x(x) L_k^y(x_j) = L_k^y(x)$$

If we substitute $x = y_m$, then we get

$$\sum_{j=0}^N L_j^x(y_m) L_k^y(x_j) = L_k^y(y_m) = \delta_{k,m}$$

which proves 1.

2. Again, the Lagrange polynomials, based on the grid points y_j , are polynomials of degree N ; therefore, their derivative can be represented as

$$(L_i^y)'(x) = \sum_{j=0}^N L_i^y(x_j) (L_j^x)'(x) \quad (31)$$

By the same token,

$$(L_j^x)'(x) = \sum_{l=0}^N (L_j^x)'(x_l) L_l^x(x) \quad (32)$$

Now, we substitute $x = y_m$ in (31) and (32) to get

$$(L_i^y)'(y_m) = \sum_{j=0}^N \sum_{l=0}^N L_i^y(x_j) (L_j^x)'(x_l) L_l^x(y_m) \quad (33)$$

The left-hand side is the (m, i) element of D_y , whereas the right-hand side is the (m, i) element of $T^{-1} D_x T$; thus, (30) has been proved.

□

3 The Legendre Galerkin Method Based on Arbitrary Grids

Consider now the linear form of (1):

$$U_t(x, t) = U_x(x, t) \quad -1 \leq x \leq 1 \quad (34)$$

$$U(x, 0) = f(x) \quad (35)$$

$$U(1, t) = g(t) \quad (36)$$

We introduce a new method for the discretization of (34), based on the differentiation matrix introduced in the last section. Note that the differentiation matrix uses the arbitrary grid x_j . With the new method, we seek a vector

$$\vec{u} = [u_0(t), \dots, u_N(t)]^T$$

that satisfies

$$\mathbf{M} \frac{d\vec{u}}{dt} = \mathbf{S}\vec{u} - \tau \vec{e}_0 [u_0 - g(t)] \quad (37)$$

where

$$\vec{e}_0 = (1, 0, 0, \dots, 0)^T$$

The discussion on imposing the initial condition is deferred until later in the paper because of subtle issues that involve convergence. Here, we generally will not use

$$u_j(0) = f(x_j) \quad 0 \leq j \leq N$$

unless the grid points x_j have special properties.

The structure of the matrices \mathbf{M} and \mathbf{S} , indicated in (34) and (37), leads immediately to the following stability result:

Theorem 4.1:

The method described in (37) is stable for $\tau \geq \frac{1}{2}$.

Proof:

We multiply (37) by \vec{u}^T to get

$$\vec{u}^T \mathbf{M} \frac{d\vec{u}}{dt} = \vec{u}^T \mathbf{S} \vec{u} - \tau \vec{u}^T \vec{e}_0 [u_0 - g(t)] \quad (38)$$

We use the symmetry property for \mathbf{M} and the almost skew symmetric property (25) for \mathbf{S} to obtain

$$\frac{1}{2} \frac{d}{dt} \vec{u}^T \mathbf{M} \vec{u} = \frac{1}{2} (u_0^2 - u_N^2) - \tau u_0 [u_0 - g(t)] \quad (39)$$

For stability, we consider the case $g(t) = 0$; from this case we can clearly determine that if $\tau \geq \frac{1}{2}$, then

$$\frac{1}{2} \frac{d}{dt} \vec{u}^T \mathbf{M} \vec{u} \leq 0 \quad (40)$$

and stability exists in the norm induced by the positive-definite matrix \mathbf{M} .

□

The stability result (39) can be represented in a different way in view of the equivalency between vectors and polynomials established in (24). Specifically, let $u_N(x, t)$ be an N th-degree polynomial such that

$$u_N(x_j, t) = u_j(t) \quad 0 \leq j \leq N$$

Then, from (24) we see that

$$\begin{aligned} \frac{1}{2} \frac{d}{dt} \int_{-1}^1 u_N(x, t)^2 dx &= \frac{1}{2} \frac{d}{dt} \tilde{u}^T \mathbf{M} \tilde{u} \\ &= \frac{1}{2} (u_0^2 - u_N^2) - \tau u_0^2 \\ &= \frac{1}{2} [u_N(1, t)^2 - u_N(-1, t)^2] - \tau u_N(1, t)^2 \end{aligned}$$

Thus, for the polynomial $u_N(x, t)$ we have stability in the usual L_2 norm.

Now, we examine equation (37) from yet another point of view. By, multiplying (37) by \mathbf{M}^{-1} , we get

$$\frac{d\tilde{u}}{dt} = \mathbf{M}^{-1} \mathbf{S} \tilde{u} - \tau \mathbf{M}^{-1} \tilde{e}_0 [u_0 - g(t)] \quad (41)$$

or in view of (21), we obtain

$$\frac{d\tilde{u}}{dt} = D \tilde{u} - \tau \mathbf{M}^{-1} \tilde{e}_0 [u_0 - g(t)]. \quad (42)$$

The expression $\mathbf{M}^{-1} \tilde{e}_0$ can be evaluated explicitly.

Theorem 4.2:

Let \mathbf{M} be the mass matrix defined in (16). Define the residual vector \vec{r} by

$$\mathbf{M}^{-1} \tilde{e}_0 = \vec{r} = (r_0, \dots, r_N)^T$$

Then,

$$r_j = \frac{P'_{N+1}(x_j) + P'_N(x_j)}{2} \quad (43)$$

where $P_N(x)$ is the Legendre polynomial of order N .

Proof:

We must verify that if \vec{r} satisfies (43), then the expression

$$\mathbf{M}\vec{r} = \vec{e}_0$$

is also satisfied. Substituting (16) into expression (3) yields

$$\begin{aligned} (\mathbf{M}\vec{r})_i &= \sum_{j=0}^N m_{i,j} r_j \\ &= \int_{-1}^1 L_i(x) \sum_{j=0}^N L_j(x) r_j dx \end{aligned} \quad (44)$$

Substituting expression (43) into (44) yields

$$(\mathbf{M}\vec{r})_i = \int_{-1}^1 L_i(x) \sum_{j=0}^N L_j(x) \left[\frac{P'_{N+1}(x_j) + P'_N(x_j)}{2} \right] dx \quad (45)$$

Because P'_{N+1} and P'_N are polynomials of degree $\leq N$, they coincide with their N th-degree interpolation polynomials; therefore,

$$\sum_{j=0}^N L_j(x) \left[\frac{P'_{N+1}(x_j) + P'_N(x_j)}{2} \right] = \left[\frac{P'_{N+1}(x) + P'_N(x)}{2} \right]$$

so that

$$\begin{aligned} (\mathbf{M}\vec{r})_i &= \int_{-1}^1 L_i(x) \left[\frac{P'_{N+1}(x) + P'_N(x)}{2} \right] dx \\ &= L_i(1) \left[\frac{P_{N+1}(1) + P_N(1)}{2} \right] - L_i(-1) \left[\frac{P_{N+1}(-1) + P_N(-1)}{2} \right] \\ &\quad - \int_{-1}^1 L'_i(x) \left[\frac{P_{N+1}(x) + P_N(x)}{2} \right] dx \end{aligned}$$

Recall that

$$P_N(1) = 1, P_N(-1) = (-1)^N$$

and that P_N and P_{N+1} are orthogonal to all polynomials of degree $< N$; the last two terms in the right-hand side of (46) vanish, and we are left with

$$(\mathbf{M}\vec{r})_i = L_i(1) = \delta_{i,0}$$

which proves theorem 4.2.

□

Theorem 4.2 sheds a new light on the connection between the method defined in (37) that uses the arbitrary set of grid points x_j and the Legendre Galerkin method. *They are the same method.*

Theorem 4.3:

The method defined in (37) is equivalent to the Legendre Galerkin method.

Proof:

Define

$$u_N(x, t) = \sum_{j=0}^N u_j(t) L_j(x)$$

where $u_j(t)$ are the elements of \vec{u} defined in (37). Then, $u_N(x, t)$ satisfies the error equation

$$\frac{\partial u_N(x, t)}{\partial t} = \frac{\partial u_N(x, t)}{\partial x} - \tau \left[\frac{P'_{N+1}(x) + P'_N(x)}{2} \right] [U_N(1, t) - g(t)] \quad (46)$$

The error equation is satisfied because both sides of expression (46) are polynomials of degree N . The two sides agree at $N + 1$ points x_j ($j = 0, \dots, N$) by virtue of (37), which indicates that they are equivalent. Because the right-hand side is orthogonal to all polynomials of degree N that vanish on the boundary $x = 1$, this error equation is the same equation that is satisfied by the Legendre Galerkin method [6].

□

As equation (46) demonstrates, the precise method for imposing the boundary conditions affects the overall behavior of the method. Section 3 shows that two differentiation operators defined on different grids are similar and, thus, have the same eigenvalues. We now show that the modified differentiation matrix also has this property. Equation (42) produces a modified differentiation matrix (i.e., a differentiation matrix that takes into account the boundary conditions):

$$D - \tau R$$

where the boundary matrix R is defined as

$$R_{i,j} = r_i \delta_{j,0} \quad (47)$$

Suppose now that we have two grids x_j, y_j ($j = 0, \dots, N$). We have shown in theorem 2.3 that D_x and D_y are similar:

$$D_y = T D_x T^{-1}$$

where the matrices T and T^{-1} are defined in (28) and (29). We show now that the same similarity transformation exists for the modified differentiation matrices. That is,

$$D_y - \tau R_y = T (D_x - \tau R_x) T^{-1} \quad (48)$$

or (with theorem 3.1)

$$R_y = T R_x T^{-1}$$

Consider element (i, j) of the right-hand side:

$$\begin{aligned} (T R_x T^{-1})_{i,j} &= \sum_{l=0}^N \sum_{m=0}^N T_{i,l} (R_x)_{l,m} T_{m,j}^{-1} \\ &= \sum_{l=0}^N \sum_{m=0}^N L_l^x(y_i) r_l \delta_{0,m} L_j^y(x_m) \end{aligned}$$

We recall that

$$r_l = \left[\frac{P'_{N+1}(x_l) + P'_N(x_l)}{2} \right] = R_N(x_l)$$

where $R_N(x)$ is a polynomial of degree N and is, therefore, equal to its interpolant. Thus,

$$(T R_x T^{-1})_{i,j} = R_N(y_i) L_j^y(x_0) = R_N(y_i) \delta_{j,0}$$

which proves that the similarity transformation is valid even for the modified derivative matrix.

The Legendre Galerkin method defined by equation (37) is stable; therefore, the initial error is not amplified. However, the effects of initial conditions must be carefully taken into account. We know that polynomials based on arbitrary grid distributions may generally be nonconvergent (the Runge phenomenon).

The initial error can be decreased with the number of mesh points N by constructing the Chebyshev interpolation as an initial condition. Thus, let

$$\xi_j = \cos\left(\frac{\pi j}{N}\right) \quad 0 \leq j \leq N \quad (49)$$

$$L^\xi(x) = (1 - x^2) T'_N(x) \quad (50)$$

$$L_j^\xi(x) = \frac{L^\xi(x)}{(x - \xi_j)(L^\xi)'(\xi_j)} \quad (51)$$

The Chebyshev approximation for the initial condition is, then,

$$C_N f(x) = \sum_{j=0}^N f(\xi_j) L_j^\xi(x) \quad (52)$$

so that the recommended initial approximation will be

$$f(x_j) \sim \sum_{j=0}^N f(\xi_j) L_j^\xi(x_j)$$

This approximation will provide a convergent approximation for the initial condition. Of course, the Chebyshev approximation is not the only possibility; any other spectral or pseudospectral approximation would do as well.

We now briefly discuss the issue of implementation. Two methods are available for implementing the arbitrary-grid spectral methods. The first method is to form the matrices \mathbf{M} and \mathbf{S} by carrying out explicitly the integrations in (16) and (17). (This technique is utilized in the two examples presented later in the text.) This procedure is done *once and for all* for every given set of grid points. Then, the equations are solved as described in (37). A more convenient method that does not involve evaluating integrals is to use the differentiation matrix D defined in (13) and solve the system (42) with the identity

$$\mathbf{M}^{-1} \vec{e}_0 = \frac{P'_{N+1}(x_j) + P'_N(x_j)}{2}$$

proven in theorem 4.2. For a large N , the method that will be the most successful is the one with the least sensitivity to round-off errors. This point has not been fully investigated at this time.

Finally, an observation in regard to the maximum allowable time step for the arbitrary-grid spectral schemes. All spatial operators have the same eigenvalues, regardless of the spatial distribution of points (48). Therefore, the maximum allowable time step is the same for all schemes. Stability is a matrix property, and depends on all the points in the distribution. This observation is somewhat counter to the conventional finite-difference notion, in which the maximum time step is governed by the smallest grid spacing.

4 The Legendre Collocation for Unstructured Grids

The Legendre collocation for unstructured grids involves the approximation of the integrals in (16) and (17) by the Gauss-Lobatto-Legendre (GLL) quadrature formula. Let $(\eta_0 = 1, \eta_1, \dots, \eta_{N-1}, \eta_N = -1)$ be the nodes of the GLL quadrature formula and ω_l , $0 \leq l \leq N$ be the weights. We define a new mass matrix \mathbf{M}_c by

$$\mathbf{M}_c(i, j) = \sum_{l=0}^N L_j(\eta_l) L_k(\eta_l) \omega_l \quad (53)$$

where the $L_j(x)$ are the Lagrange polynomials at the points $(x_0 = 1, x_1, x_2, \dots, x_{N-1}, x_N = -1)$. Note that this is an arbitrary set of grid points.

The matrix \mathbf{M}_c may be different from \mathbf{M} because the GLL formula is exact to order $2N-1$ and $L_j(x)L_k(x)$ is a polynomial of order $2N$. The matrix \mathbf{M}_c is, however, a symmetric and positive-definite matrix.

By introducing quadrature to equation (17), we define a new stiffness matrix \mathbf{S}_c as

$$\mathbf{S}_c(i, j) = \sum_{l=0}^N L'_j(\eta_l) L_k(\eta_l) \omega_l \quad (54)$$

Note that because of the exactness of the GLL formula, the sum on the right-hand side of (54) is the same as the integral in the right-hand side of (17); therefore,

$$\mathbf{S}_c = \mathbf{S}$$

For this reason, the property (25) is true for the stiffness matrix \mathbf{S}_c also.

The uniqueness of the differentiation matrix D also yields

$$\mathbf{M}_c^{-1} \mathbf{S}_c = \mathbf{M}^{-1} \mathbf{S}$$

which does not contradict the fact that

$$\mathbf{M}_c \neq \mathbf{M}$$

because the matrices \mathbf{S}_c and \mathbf{S} are singular.

In the Legendre collocation method of (34) with arbitrary grids, we seek a vector

$$\vec{u} = [u_0(t), \dots, u_N(t)]^T$$

that satisfies

$$\mathbf{M}_c \frac{d\vec{u}}{dt} = \mathbf{S}_c \vec{u} - \tau \vec{e}_0 [u_0 - g(t)] \quad (55)$$

where

$$\vec{e}_0 = (1, 0, 0, \dots, 0)^T$$

Alternatively,

$$\frac{d\vec{u}}{dt} = D\vec{u} - \tau \mathbf{M}_c^{-1} \vec{e}_0 [u_0 - g(t)] \quad (56)$$

The stability of (55) follows immediately from the fact that \mathbf{M}_c is symmetric positive definite and \mathbf{S}_c satisfies (25). Our aim is to show that (55) is equivalent to the usual Legendre collocation method.

Theorem 4.1:

Let \mathbf{M}_c be the mass matrix defined in (53). We define the residual vector \vec{r} by

$$\mathbf{M}_c^{-1} \vec{e}_0 = \vec{r} = (r_0, \dots, r_N)^T$$

Then,

$$r_j = P'_N(x_j)(1 + x_j) \frac{1}{2N^2} \quad (57)$$

where $P_N(x)$ is the Legendre polynomial of order N .

Proof:

We start by noting that the nodes η_l of the GLL formula are the zeroes of the polynomial

$$P'_N(x)(1 - x^2)$$

Also, because $P'_N(1 + x)$ is a polynomial of degree N ,

$$\sum_{j=0}^N L_j(x) P'_N(x_j)(1 + x_j) = P'_N(x)(1 + x)$$

Therefore,

$$\begin{aligned} 2N^2 (\mathbf{M}_c \vec{r})_i &= \sum_{j=0}^N \mathbf{M}_c(i, j) P'_N(x_j)(1 + x_j) \\ &= \sum_{j=0}^N \sum_{l=0}^N L_j(\eta_l) L_i(\eta_l) P'_N(x_j)(1 + x_j) \omega_l \\ &= \sum_{l=0}^N L_i(\eta_l) P'_N(\eta_l)(1 + \eta_l) \omega_l \\ &= 2N^2 \delta_{0,i} \end{aligned}$$

which proves the theorem.

□

Equation (56) also can be viewed as shown in the following example. We seek a polynomial (in x) $u_N(x, t)$ of the form

$$u_N(x, t) = \sum_{j=0}^N u_N(x_j, t) L_j(x) \quad (58)$$

such that

$$\frac{du_N(x_k, t)}{dt} = \sum_{j=0}^N u_N(x_j, t) L'_j(x_k) - \tau R_N(x_k) [u_N(1, t) - g(t)] \quad (59)$$

where

$$R_N(x) = (1 + x) P'_N(x)$$

This approach is equivalent to the Legendre Collocation Method [6].

The extension of the arbitrary-grid Legendre collocation method from the linear case (34) to the solution of the nonlinear case (1) is immediate. The issue of implementation could be significant. To avoid computing the points η_l , the best choice is to use the formulation (56) rather than (55). In this case, \mathbf{M}_c and \mathbf{S}_c do not need to be computed.

At this stage, note that for the case

$$R_N(x) = P_N(x)$$

we have the Legendre Tau method, with the additional property of an improved time step. However, we do not have the representation of the Legendre Tau method in the form of (55).

5 Unstructured Grids for Unbounded Domains: Laguerre Methods

Consider the equation

$$\frac{\partial U}{\partial t} = -\frac{\partial U}{\partial x} \quad 0 \leq x < \infty \quad (60)$$

$$U(0, t) = g(t) \quad (61)$$

$$U(x, 0) = h(x)$$

Note that the domain is semibounded. Note also that if $g(t) = 0$, then

$$\frac{d}{dt} \int_0^\infty e^{-x} U^2(x, t) dx = - \int_0^\infty e^{-x} U^2(x, t) dx \quad (62)$$

Assume that we have an arbitrary set of grid points

$$(x_0 = 0, x_1, \dots, x_N)$$

In the Galerkin procedure, we approximate the derivative of a function $f(x)$ whose values at x_j are given by the derivative of its interpolant $f_N(x)$. After we define

$$L(x) = (x - x_0) \dots (x - x_N) \quad (63)$$

we define the Lagrange polynomials by

$$L_j(x) = \frac{L(x)}{(x - x_j)L'(x_j)} \quad (64)$$

The derivative of the interpolant $f_N(x)$ has two equivalent expressions:

$$\frac{df_N(x)}{dx} = \sum_{j=0}^N f(x_j) \frac{dL_j(x)}{dx} \quad (65)$$

and

$$\frac{df_N(x)}{dx} = \sum_{j=0}^N f'_N(x_j) L_j(x) \quad (66)$$

In the Galerkin Laguerre method, we express the equivalency between the expressions by

$$\int_0^\infty e^{-x} \sum_{j=0}^N \left[f(x_j) \frac{dL_j(x)}{dx} - f'_N(x_j) L_j(x) \right] L_k(x) dx = 0 \quad 0 \leq k \leq N \quad (67)$$

Equation (67) defines the differentiation matrix D . In fact, if we define

$$m_{k,j} = (L_j, L_k) \quad (68)$$

and

$$s_{k,j} = (L'_j, L_k) \quad (69)$$

where the scalar product (u, v) is defined as

$$(u, v) = \int_0^\infty e^{-x} u(x) v(x) dx$$

then we get

$$D = \mathbf{M}^{-1} \mathbf{S}$$

As before, the differentiation matrix is unique. The manner in which the matrices \mathbf{M} and \mathbf{S} are constructed leads immediately to the following lemma.

Lemma 6.1:

The matrix \mathbf{M} is symmetric positive definite. The matrix \mathbf{S} satisfies

$$\mathbf{S} + \mathbf{S}^T = \mathbf{M} - \text{diagonal}(1, 0, 0, \dots, 0) \quad (70)$$

Proof:

From the definition of the matrix \mathbf{S} , we have

$$\begin{aligned} s_{k,j} &= (L'_j, L_k) \\ &= \int_0^\infty e^{-x} L'_j(x) L_k(x) dx \\ &= -L_j(0) L_k(0) - \int_0^\infty e^{-x} L_j(x) L'_k(x) dx \\ &\quad + \int_0^\infty e^{-x} L_j(x) L_k(x) dx \end{aligned}$$

By using the definition of the matrix \mathbf{M} and the properties of the Lagrange polynomials, we get

$$s_{k,j} = -\delta_{i,0} \delta_{j,0} - s_{j,k} + m_{k,j} \quad (71)$$

which proves (70). □

To discretize (60), we introduce the unknown vector

$$\vec{u} = [u_0(t), \dots, u_N(t)]^T$$

that satisfies

$$\mathbf{M} \frac{d\vec{u}}{dt} = -\mathbf{S} \vec{u} - \tau \vec{e}_0 [u_0 - g(t)] \quad (72)$$

The stability is immediate, as shown in the following lemma.

Lemma 6.2:

Let \vec{u} satisfy (72), with $g(t) = 0$. Then, we have the energy estimate

$$\frac{d}{dt} \vec{u}^T \mathbf{M} \vec{u} = -\vec{u}^T \mathbf{M} \vec{u} - (2\tau - 1) u_0^2 \quad (73)$$

Proof:

Equation (73) follows immediately from multiplying (72) by \vec{u}^T and using (70). □

Lemma 6.1 implies that the method is stable, provided that $\tau \geq \frac{1}{2}$. Note that the energy estimate (73) for the approximation is nearly the same as for the differential equation (64).

We still must show that the method described in (72) is equivalent to the Laguerre Galerkin method. We begin by rewriting (72) as

$$\frac{d\vec{u}}{dt} = -\mathbf{M}^{-1}\mathbf{S}\vec{u} - \tau\mathbf{M}^{-1}\vec{e}_0[u_0 - g(t)] \quad (74)$$

The key issue is to identify the vector

$$\mathbf{M}^{-1}\vec{e}_0$$

which is done in the following theorem.

Theorem 6.1:

Let \mathbf{M} be the mass matrix defined in (68). Define the residual vector \vec{r} by

$$\mathbf{M}^{-1}\vec{e}_0 = \vec{r} = (r_0, \dots, r_N)^T$$

Then

$$r_j = \frac{d}{dx}\mathcal{L}_{N+1}^{(0)}|_{x=x_j}, \quad (75)$$

where $\mathcal{L}_N^{(0)}$ is the Laguerre polynomial of order N .

Proof:

We must verify that \vec{r} satisfies (75) and that

$$\mathbf{M}\vec{r} = \vec{e}_0$$

We begin by expanding $(\mathbf{M}\vec{r})$ as

$$\begin{aligned} (\mathbf{M}\vec{r})_i &= \sum_{j=0}^N m_{i,j} r_j \\ &= \int_0^\infty e^{-x} L_i(x) \sum_{j=0}^N L_j(x) r_j dx \end{aligned} \quad (76)$$

If we substitute (75) into (76), then we get

$$(\mathbf{M}\vec{r})_i = \int_0^\infty e^{-x} L_i(x) \sum_{j=0}^N L_j(x) \left[\frac{d}{dx} \mathcal{L}_{N+1}^{(0)}|_{x=x_j} \right] dx \quad (77)$$

Because $\frac{d}{dx}\mathcal{L}_{N+1}^{(0)}$ is a polynomial of order N , it coincides with its interpolant; therefore,

$$\sum_{j=0}^N L_j(x) \left[\frac{d}{dx} \mathcal{L}_{N+1}^{(0)}|_{x=x_j} \right] = \frac{d}{dx} \mathcal{L}_{N+1}^{(0)}(x)$$

Thus,

$$(\mathbf{M}\vec{r})_i = \int_0^\infty e^{-x} L_i(x) \frac{d}{dx} \mathcal{L}_{N+1}^{(0)}(x) dx$$

If we integrate the right-hand side by parts, we get

$$(\mathbf{M}\vec{r})_i = -L_i(0)\mathcal{L}_{N+1}^{(0)}(0) + \int_0^\infty e^{-x} L_i(x) \mathcal{L}_{N+1}^{(0)}(x) dx - \int_0^\infty e^{-x} L'_i(x) \mathcal{L}_{N+1}^{(0)}(x) dx$$

The last two terms on the right vanish because of the orthogonality of $\mathcal{L}^{(0)}$, and the first term vanishes if $i \neq 0$; thus,

$$(\mathbf{M}\vec{r})_i = -\delta_{i,0}$$

and the theorem is proven. □

Another method for getting the Laguerre method on the grid x_j is to seek a polynomial $u_N(x, t)$ such that

$$\frac{du_N(x_k, t)}{dt} = \sum_{j=0}^N L'_j(x_k) u_N(x_j, t) - \tau \mathcal{L}_N^0(x_k) [u_N(x_0, t) - g(t)] \quad (78)$$

where $(LG)_N(x)$ is the N th-degree Laguerre polynomial. This approach is the Laguerre collocation method.

6 Numerical Results

We now test the previous theoretical results with two numerical examples. The linear equations (34)–(36) are solved with $f(x) = \sin(\pi x)$, $g(t) = \sin[\pi(1+t)]$, and the exact solution $U(x, t) = \sin[\pi(x+t)]$. A variety of grids, from Chebyshev to “randomly generated” grids, are used to test the accuracy and stability of the method. For all calculations, 128-bit arithmetic is used to ensure adequate precision.

Figure 1 shows the refinement study on five different grids:

1. Uniform grid $x_j = \frac{2j-N}{N}$ ($j = 0, \dots, N$)
2. Chebyshev grid $x_j = \cos(\frac{\pi j}{N})$
3. A linear combination of the uniform grid and Chebyshev grid (i.e., $x_j = 0.5 \frac{2j-N}{N} + 0.5 \cos(\frac{\pi j}{N})$)
4. Chebyshev² (i.e., $x_j = \cos^2(\frac{\pi j}{N})$)
5. (Chebyshev)² for $-1 \leq x \leq 0$ and (Chebyshev) ^{$\frac{1}{2}$} for $0 \leq x \leq 1$, where (Chebyshev) ^{$\frac{1}{2}$} is defined by the grid points $x_j = \cos^{\frac{1}{2}}(\frac{\pi j}{N})$.

The \log_{10} of the L_2 error, plotted against the number of points in the approximating polynomials is shown in Figure 1. The problems are run to the physical time $T = 2$. The convergence is exponential for all cases until machine round off is encountered. These results are consistent with the previous numerical results. (Note that the Chebyshev grid is the least sensitive to round off.)

The Legendre Galerkin method defined by equation (37) is stable; therefore, the initial error is not amplified. However, the effects of initial conditions must be carefully taken into account. We know that polynomials based on arbitrary grid distributions generally may be nonconvergent. This property, called the Runge phenomena, is easily demonstrated by approximating the function $f(x) = \frac{1}{1+(5x)^2}$ ($-1 \leq x \leq 1$) on a uniform grid. The global approximating polynomials oscillate wildly at each end of the domain, which yields a poor approximation in those regions. The Runge phenomena is alleviated by using a grid distribution (like the Chebyshev grid distribution), which clusters points near the boundaries -1 and 1 .

Figure 2 illustrates that a Runge-like phenomena exists within the arbitrary-grid spectral methods if special precautions are not taken in the initialization step. In this problem, the linear equations (34)–(36) are solved with $f(x) = \frac{1}{1+[5(x)]^2}$, $g(t) = \frac{1}{1+[5(1+t)]^2}$, and the exact solution $U(x, t) = \frac{1}{1+[5(x+t)]^2}$. The simulation is run to time $T = 0.001$ (a physical time that occurs well before the influence of the initialization is lost.) (Running to a physical time $T \geq 2$ yields exponential convergence on all grids.) Convergence is achieved only for the Chebyshev grid distribution.

The source of the error in this problem is the failure of the arbitrary grid that approximates the polynomial to converge to the initial condition. For small times (less than 1 convective sweep), erroneous information is left in the domain, and the resulting method is nonconvergent. By changing the problem slightly, however, convergence can be recovered on all grids.

To initialize the problem, we must construct an approximation to the initial condition $f(x)$, based on the grid points x_j ($0 \leq j \leq N$). We want to keep the flexibility and rigid structure of the original grid distribution; however, the interpolation polynomial, based on the grid points x_j , generally is not convergent. Therefore, we use the method outlined in (49) and (52). With this initialization, spectral convergence is recovered.

7 Conclusions

A new technique for implementing spectral methods for hyperbolic equations has been developed that does not require grid points that are nodes of some Gauss quadrature formula. For this reason, this method is referred to as an arbitrary-grid spectral method. Both Galerkin and collocation formulations are presented for the conventional Legendre method, and a Galerkin formulation is presented for the conventional Laguerre method.

The basis for the stability of the unstructured spectral schemes relies on a weighted energy norm in all cases. Stability is proven for the constant coefficient hyperbolic system. All unstructured spectral methods utilize a "weak" imposition of the boundary condition, similar to the technique used in the penalty formulations of the finite element method. With this imposition, the complete differentiation matrix, including boundary conditions, is similar to (i.e., it has the same eigenvalues) the conventional differentiation operator; therefore, this matrix behaves similarly.

The new formulations are demonstrated on two scalar hyperbolic problems. The arbitrary-grid Legendre Galerkin method is used in both cases. Exponential accuracy is shown in both cases on arbitrary grids. Care must be exercised in the initialization procedure to ensure convergence of the new schemes.

8 Acknowledgments

The authors would like to thank Drs. G. Erlebacher, H. O. Kreiss, D. Kopriva, and S. Abarbanel. Their comments and suggestions are greatly appreciated.

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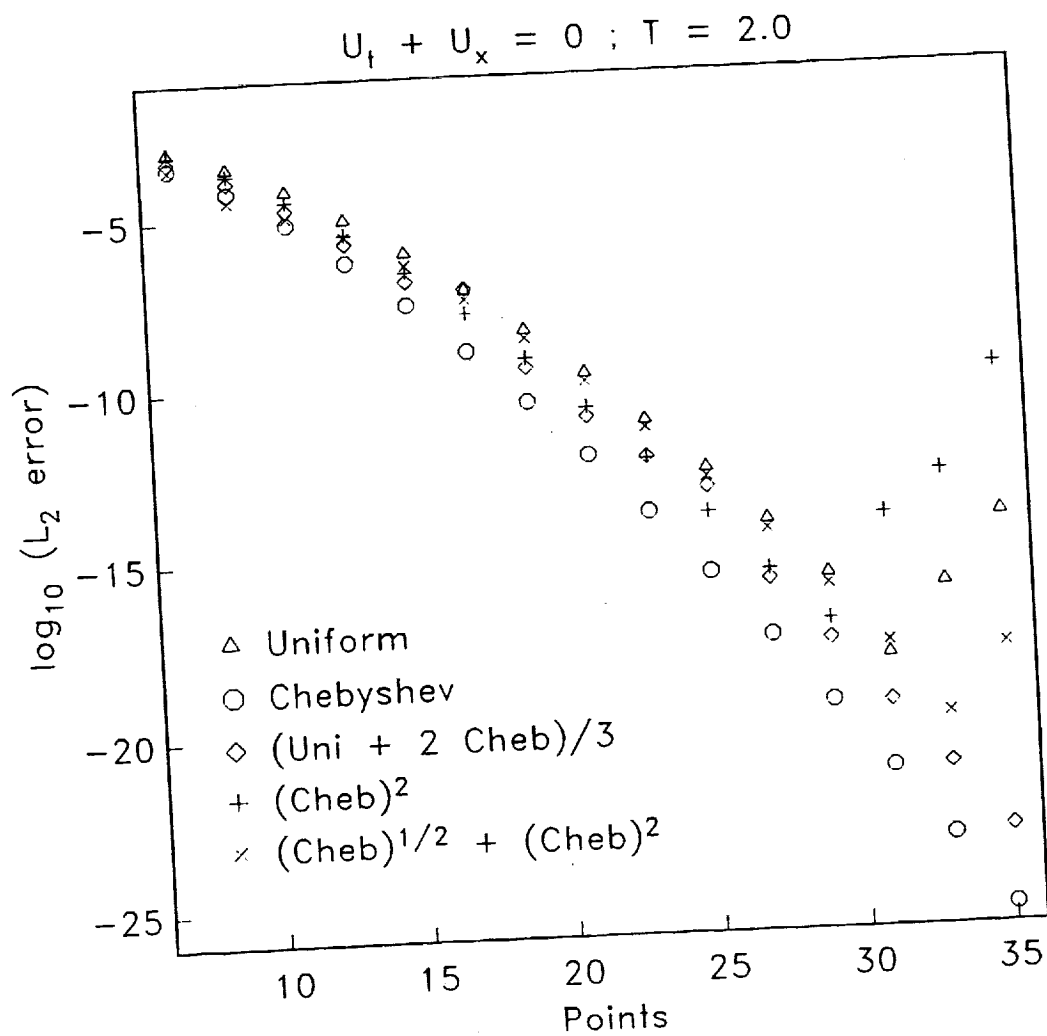


FIGURE 1. Convergence of the arbitrary grid Legendre Galerkin method on various grids.

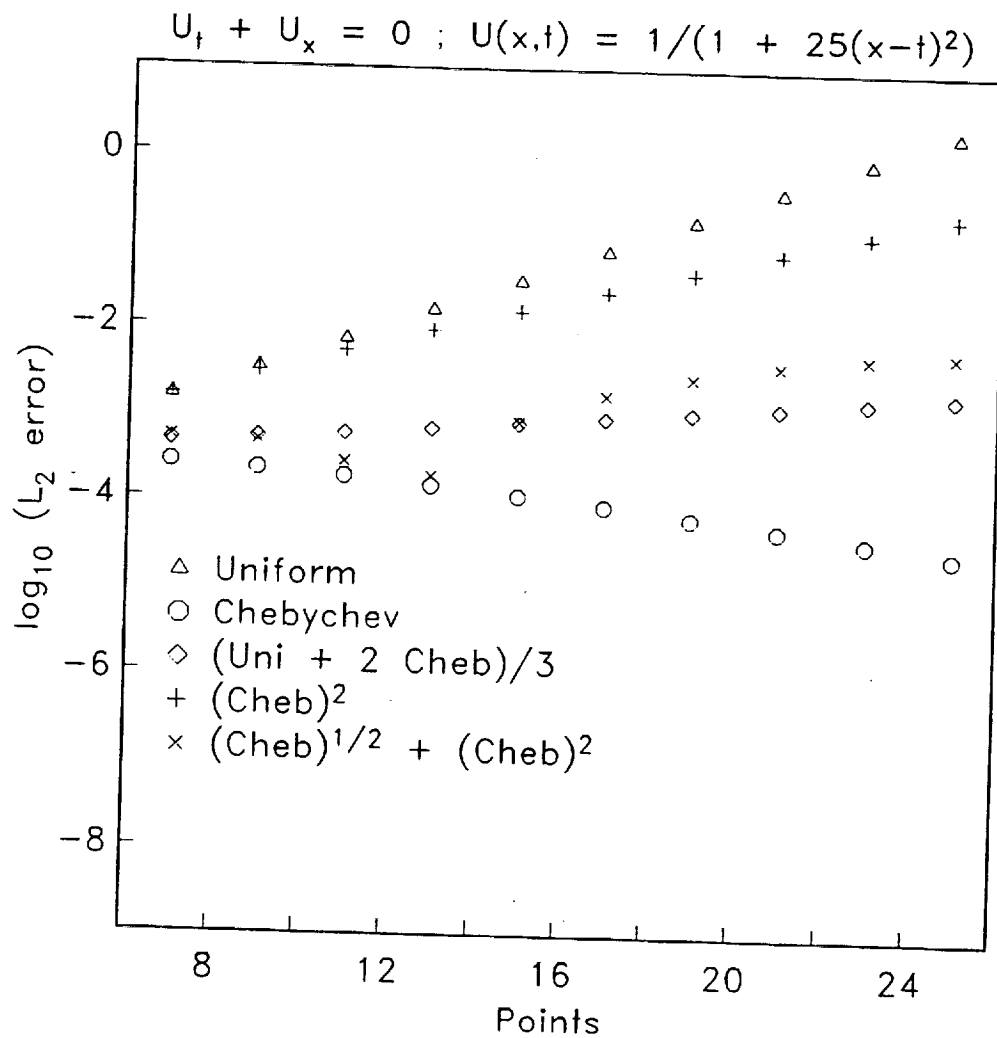


FIGURE 2. Divergence of the arbitrary grid Legendre Galerkin method for improperly imposed initial conditions.

REPORT DOCUMENTATION PAGE			Form Approved OMB No. 0704-0188	
<small>Public reporting burden for this collection of information is estimated to average 1 hour per response, including the time for reviewing instructions, searching existing data sources, gathering and maintaining the data needed, and completing and reviewing the collection of information. Send comments regarding this burden estimate or any other aspect of this collection of information, including suggestions for reducing this burden, to Washington Headquarters Services, Directorate for Information Operations and Reports, 1215 Jefferson Davis Highway, Suite 1204, Arlington, VA 22202-4302, and to the Office of Management and Budget, Paperwork Reduction Project (0704-0188), Washington, DC 20503.</small>				
1. AGENCY USE ONLY(Leave blank)	2. REPORT DATE May 1995	3. REPORT TYPE AND DATES COVERED Contractor Report		
4. TITLE AND SUBTITLE SPECTRAL METHODS ON ARBITRARY GRIDS		5. FUNDING NUMBERS C NAS1-19480 WU 505-90-52-01		
6. AUTHOR(S) Mark H. Carpenter David Gottlieb				
7. PERFORMING ORGANIZATION NAME(S) AND ADDRESS(ES) Institute for Computer Applications in Science and Engineering Mail Stop 132C, NASA Langley Research Center Hampton, VA 23681-0001		8. PERFORMING ORGANIZATION REPORT NUMBER ICASE Report No. 95-37		
9. SPONSORING/MONITORING AGENCY NAME(S) AND ADDRESS(ES) National Aeronautics and Space Administration Langley Research Center Hampton, VA 23681-0001		10. SPONSORING/MONITORING AGENCY REPORT NUMBER NASA CR-198158 ICASE Report No. 95-37		
11. SUPPLEMENTARY NOTES Langley Technical Monitor: Dennis M. Bushnell Final Report To be submitted to the Journal of Computational Physics				
12a. DISTRIBUTION/AVAILABILITY STATEMENT Unclassified-Unlimited Subject Category 34			12b. DISTRIBUTION CODE	
13. ABSTRACT (Maximum 200 words) Stable and spectrally accurate numerical methods are constructed on arbitrary grids for partial differential equations. These new methods are equivalent to conventional spectral methods but do not rely on specific grid distributions. Specifically, we show how to implement Legendre Galerkin, Legendre collocation, and Laguerre Galerkin methodology on arbitrary grids.				
14. SUBJECT TERMS Spectral Methods; Numerical Stability; Unstructured Grids			15. NUMBER OF PAGES 29	
			16. PRICE CODE A03	
17. SECURITY CLASSIFICATION OF REPORT Unclassified	18. SECURITY CLASSIFICATION OF THIS PAGE Unclassified	19. SECURITY CLASSIFICATION OF ABSTRACT	20. LIMITATION OF ABSTRACT	

