# On Time-Triggered Ethernet in NASA's Lunar Gateway

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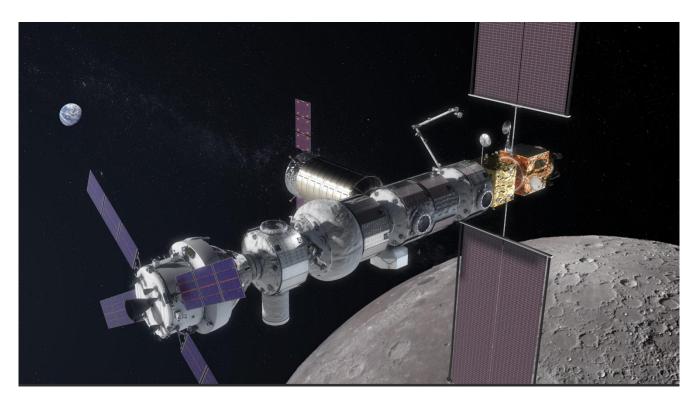
July 30, 2020



- Introduction to Gateway
- Time-Triggered Ethernet (TTE) backbone
- TTE, A Fault-Tolerant Interconnect
- TTE, An Integration Framework
- A Unique Challenge, Classical Ethernet
- Conclusion



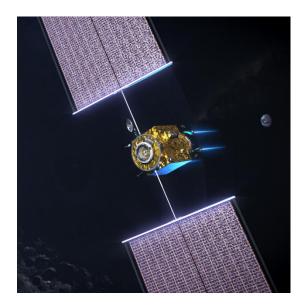
 Lunar outpost under development to support a sustainable human presence on and around the moon





- Made up of multiple modules:
  - Power and Propulsion Element (PPE)
    - Solar electric propulsion spacecraft that will provide power, communications, and attitude control
  - Habitation and Logistics Outpost (HALO)
    - Provides life support, command and control, energy storage and power distribution



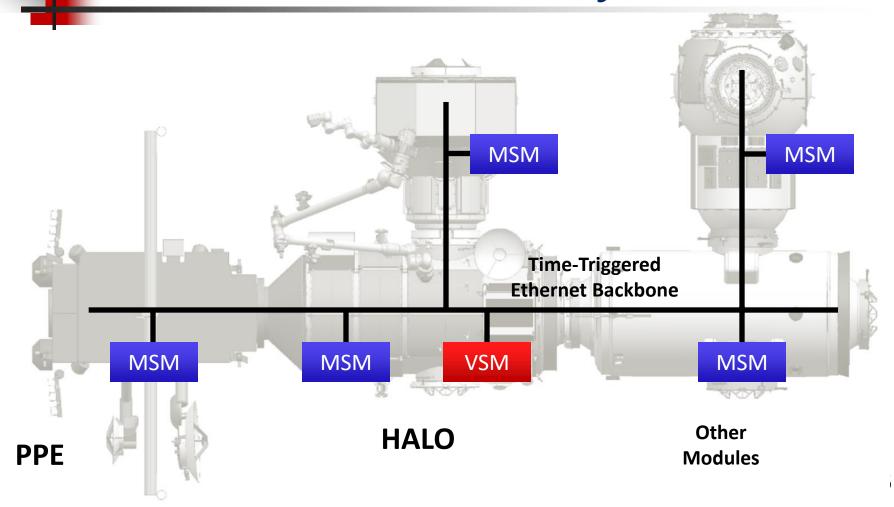




- Visiting vehicles:
  - Orion, Gateway Logistic Services, Human Landing System (HLS)







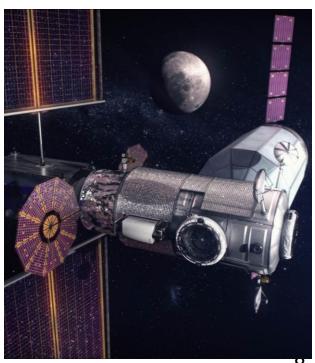


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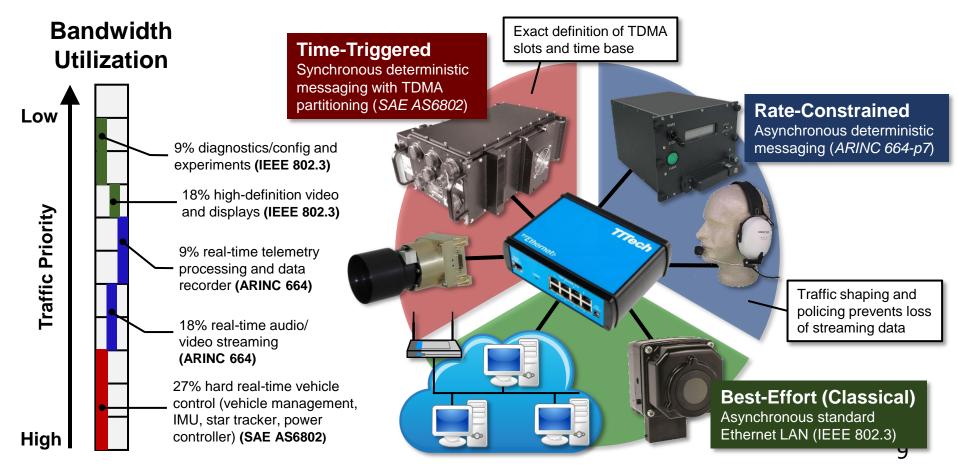
### **Competing Requirements**

- A vehicle like GW puts conflicting requirements on the network
  - It is both a safety-critical vehicle and a research platform
- Control-centric systems need networks with:
  - High integrity and availability
  - Worst case bounded latency and jitter
- Science-centric systems need networks with:
  - Compatibility with COTS devices
  - High throughput
  - Flexibility and expandability
- Often the same computer needs both:
  - E.g. IMA functions with different criticalities on the same computer

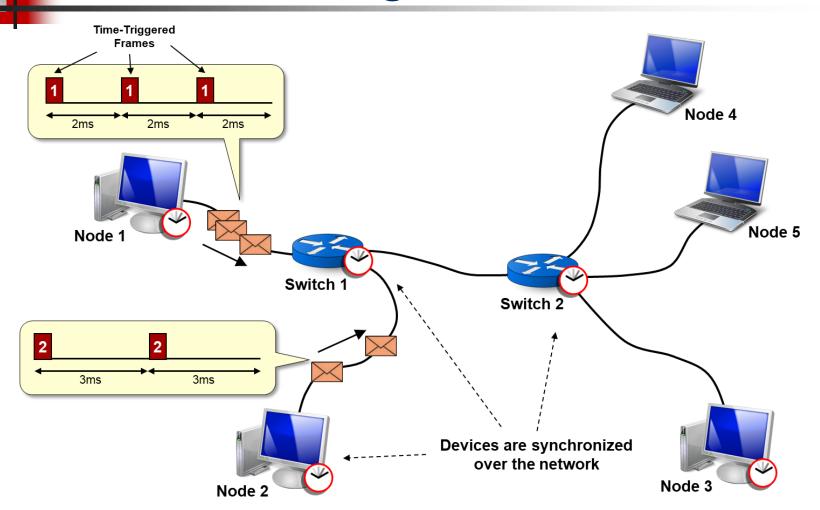




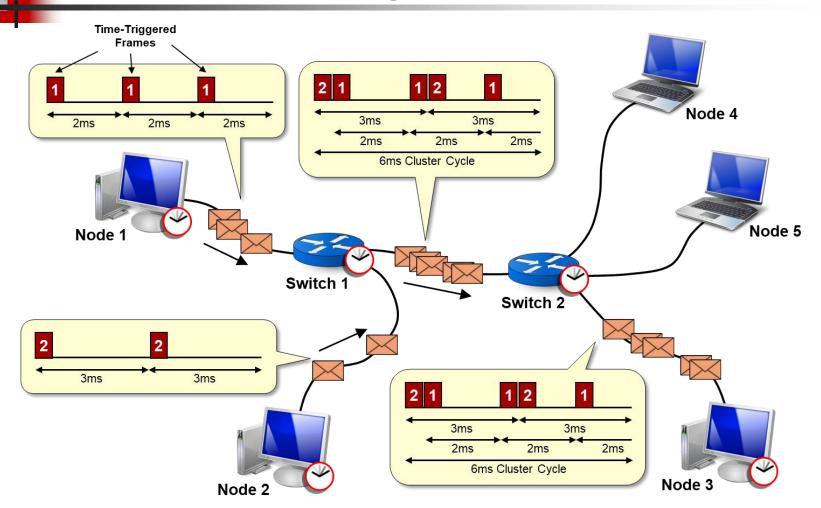
# Time-Triggered Ethernet (TTE)



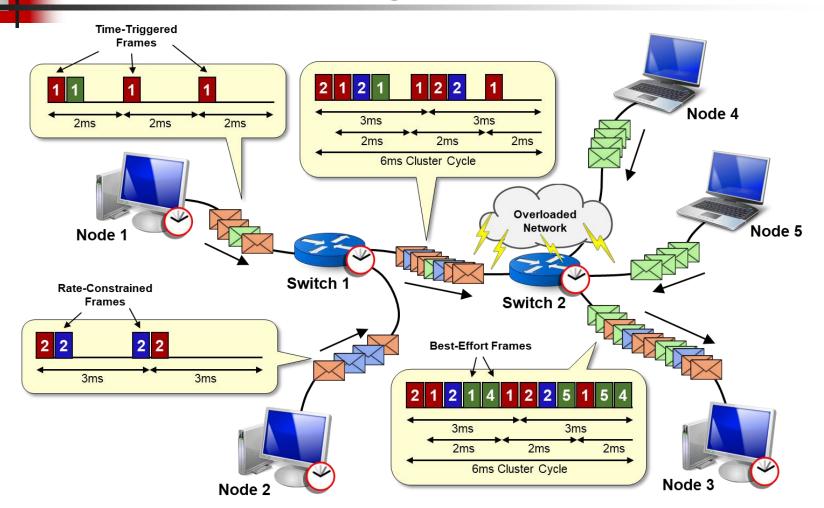
# TTE Traffic Integration



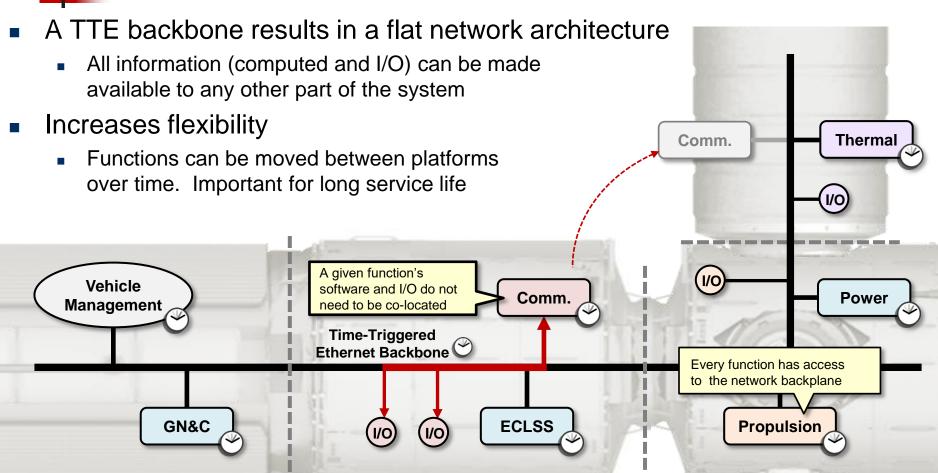
# TTE Traffic Integration



# TTE Traffic Integration









Time-Triggered
Ethernet Backbone

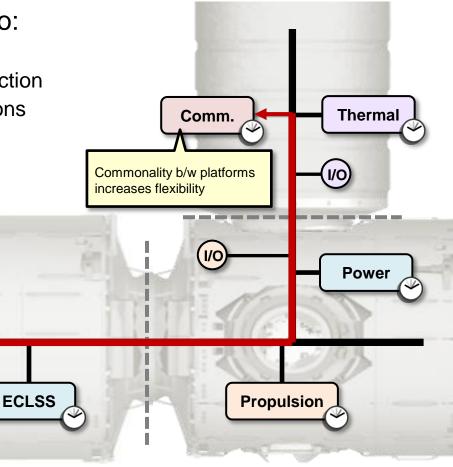
- A table-driven approach can be used to:
  - Assign functions to different computers
  - Assign CPU/memory resources to each function
  - Configure messaging paths between functions
- Integration with cFS

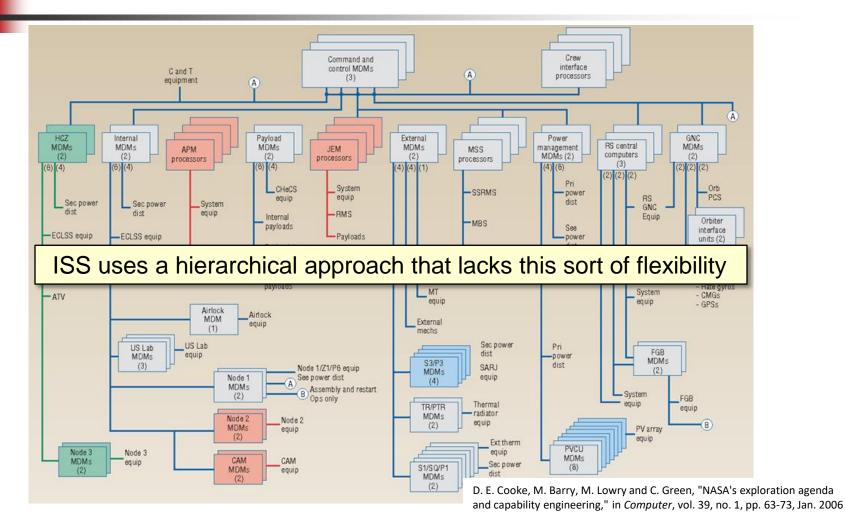
**GN&C** 

**Vehicle** 

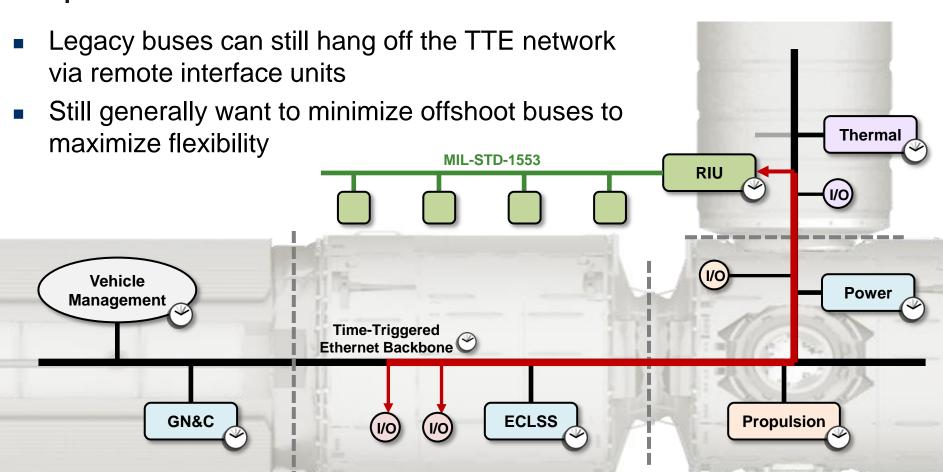
Management

Couple cFS tables and TTE network tables







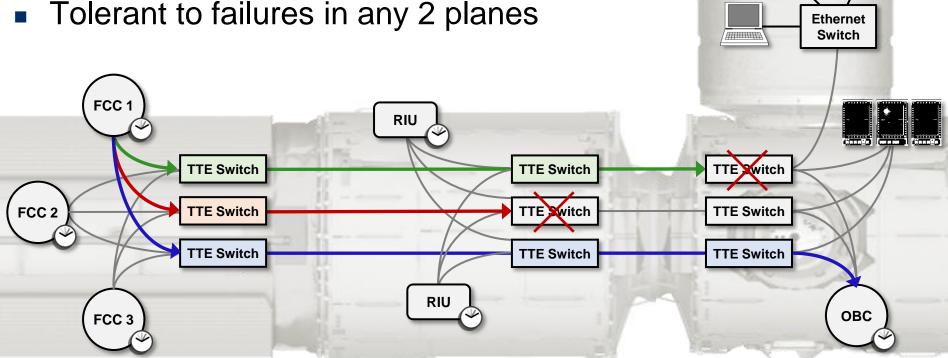


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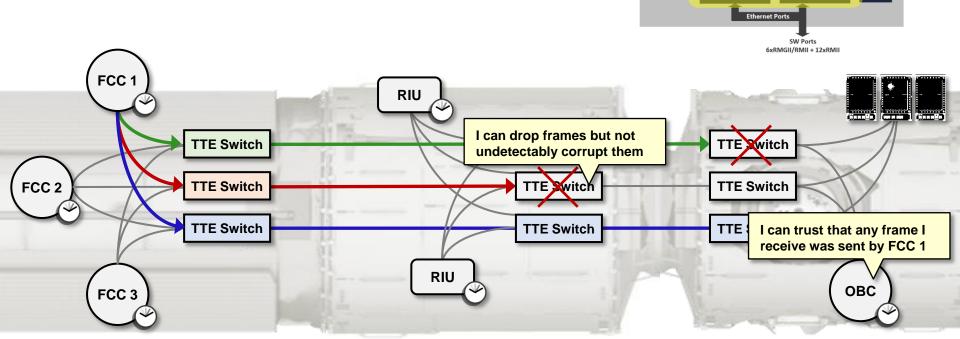
# TTE Network Availability

- Network availability comes from 3 planes
- Traffic goes over all planes simultaneously
- Tolerant to failures in any 2 planes



# TTE Network Integrity

- High-integrity switch design prevents undetectable frame corruption
- Uses COM/MON with two switch IPs



RGU

CGU

IOMU

MCU

Management

Functions

TTE-Switch

Ethernet Port

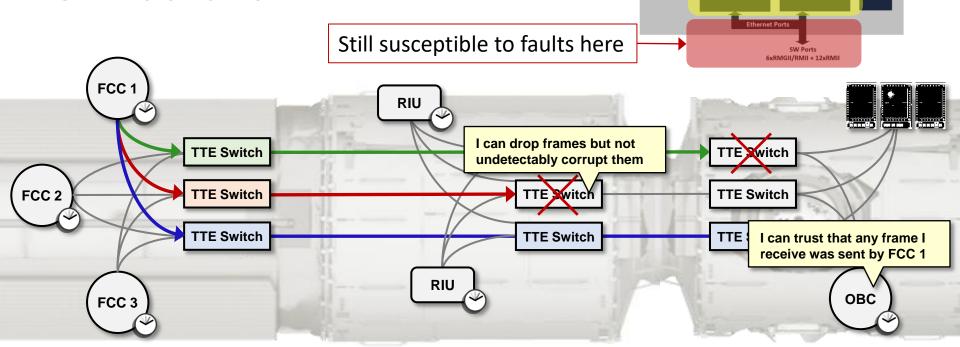
TTE-End System

Controller

Config

# TTE Network Integrity

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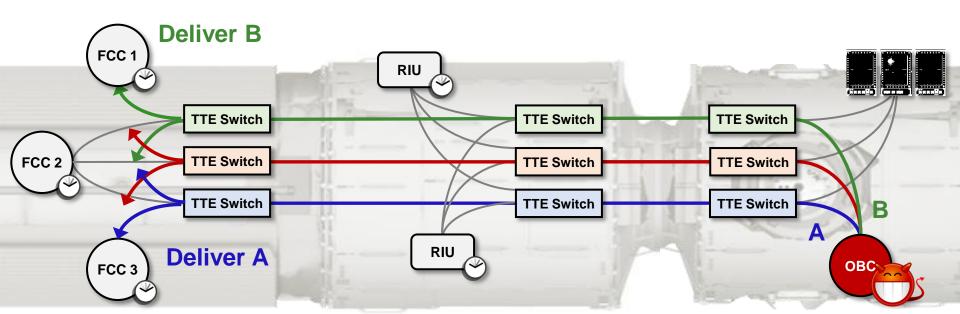
Controller

Config &



### TTE Asymmetric Transmissions

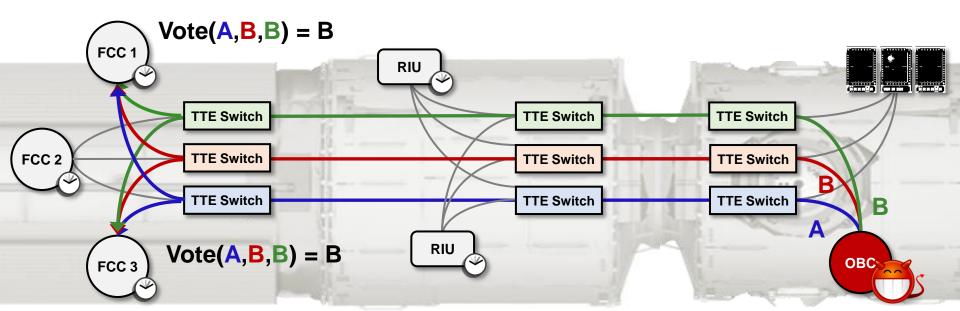
- A faulty ES can send different frames to different planes
- Receivers deliver the first valid frame
- Means different receivers can deliver different frames!



# 4

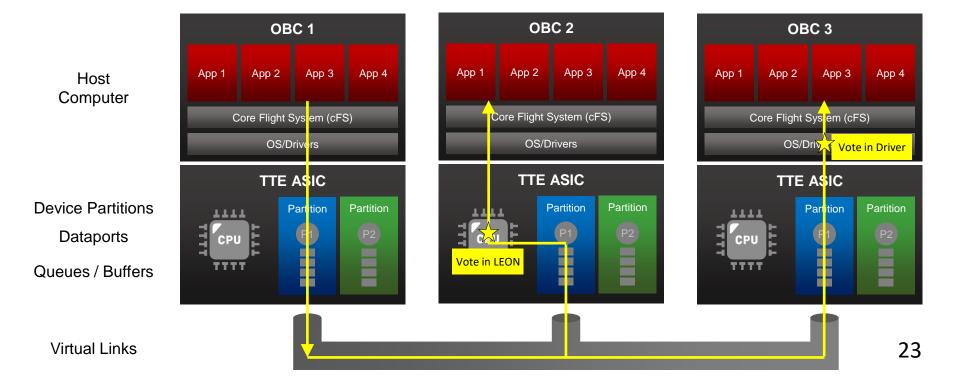
### TTE Asymmetric Transmissions

- Return all redundant frames to the host
- Hybrid majority vote (exclude manifest faulty frames)
- Guarantees all correct receivers deliver the same message



# Voting Realization

- Voting is configurable on a VL-by-VL basis groups redundant frames by arrival time
- Can be performed in the Leon CPU in TTE ASIC, or in the driver software





#### TTE Broadcast Channel

Lets you realize a broadcast channel abstraction

- Given 1 faulty end system:
  - If the sender is correct and broadcasts v, all correct receivers get v
  - All correct receivers deliver the same value
- Given 1 faulty end system + 1 faulty switch:
  - If the sender is correct and broadcasts v, all correct receivers get v
  - All correct receivers <u>that deliver a value</u> deliver the same value



#### TTE Broadcast Channel

Lets you realize a broadcast channel abstraction

Given 1 faulty end system:

- If the sender is correbyzantines Agreement correct receivers
- All correct receivers deliver the same value.
- Given 1 faulty end system + 1 faulty switch:
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  - All correct receivers that deliver a value deliver the same value

nice degradation



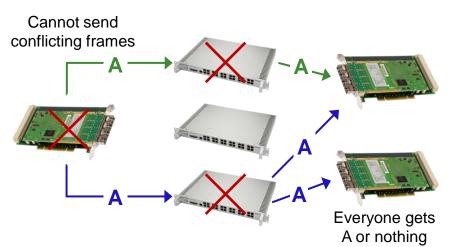
### Integrity/Availability Tradeoff

#### Orion (Phoenix IP)



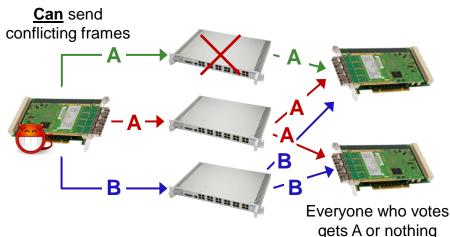


- Has <u>high-integrity</u> end systems
- Crusader broadcast channel with faulty ES and two faulty switches





- Only **standard-integrity** end systems
- Crusader broadcast channel with faulty ES and one faulty switch





### **Building on Broadcast Channels**

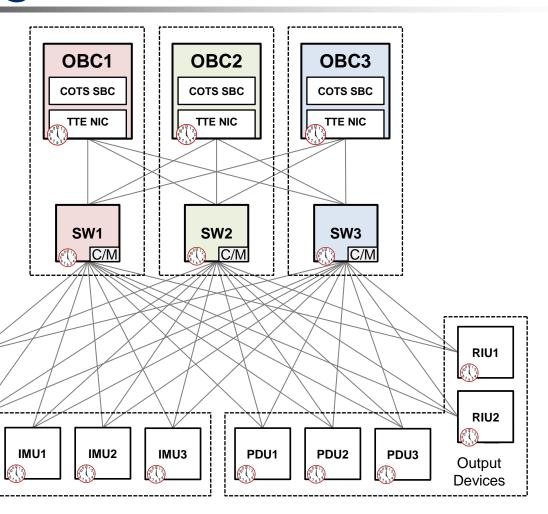
**Example**: 1 Byzantine-resilient fault-tolerant switched triplex

COM<sub>1</sub>

COM<sub>2</sub>

Input

**Devices** 



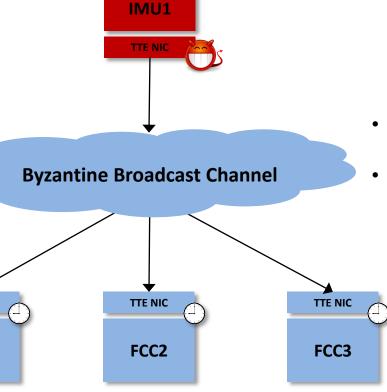


### Agreeing on External Data

**Key tenet of SMR:** Replicas must go through same state transitions

**How to accomplish?**: Replicas run Byzantine Agreement protocols

on all input data



 Makes agreement on inputs (and state) trivial

 No explicit 2 round exchange needed

Must be consistent -

**Deliver: A** 

TTE NIC

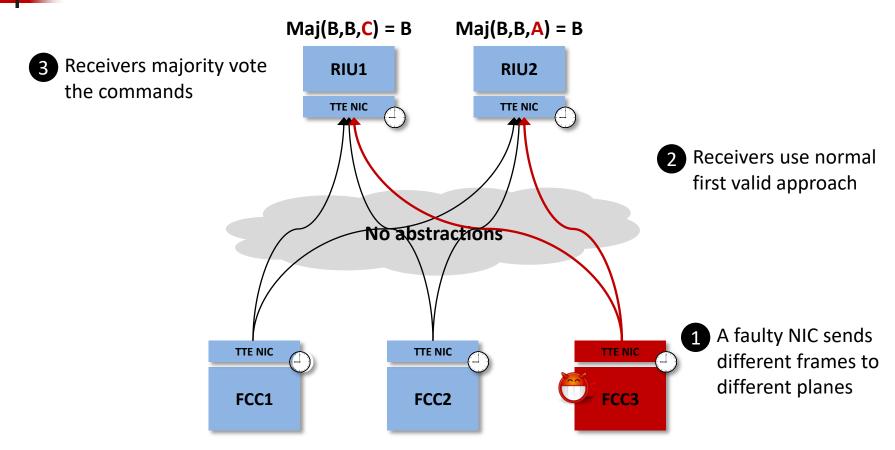
FCC1

**Deliver: A** 

**Deliver: A** 

# Comma

### **Commanding Actuators**





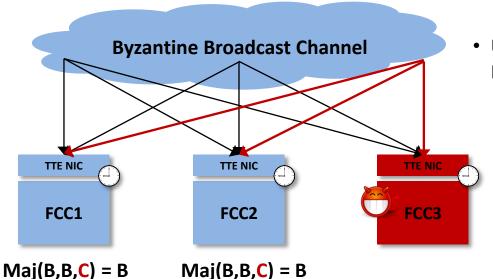
### Fault Diagnosis

**Happening Simultaneously ...** 



1 Outputs are reflected back via broadcast channel

2 Replicas vote the consistent outputs from the replicas. Any replicathat disagrees is faulty.

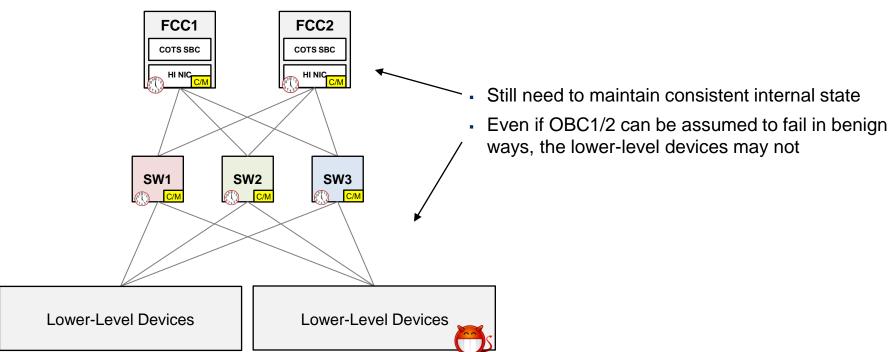


 Uses same VLs as RIUs are reading



### **Building on Broadcast Channels**

- Broadcast channels are applicable to many fault models and fault tolerance approaches
- E.g. A primary-backup system constrained to symmetric transmissive failures



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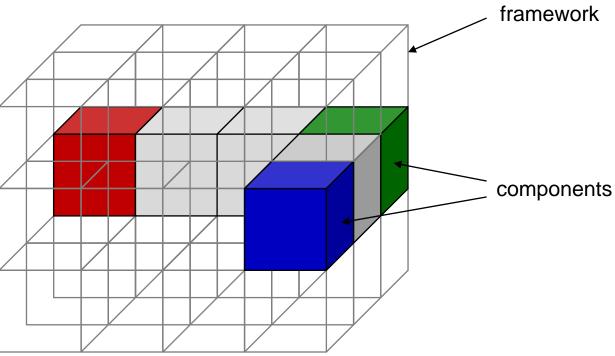


### TTE: It's a Framework

TTE is not just an interconnect, it is an integration framework

Framework provides the supporting structure for integrating components into the system

- Components can rely on the framework
- Framework doesn't rely on the components





### TTE: It's a Framework

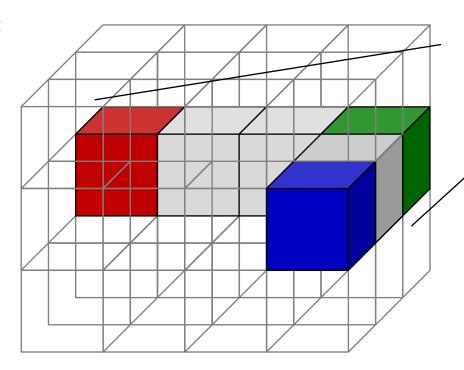
TTE is not just an interconnect, it is an integration framework

TTE provides multiple formally verified services:

- Clock synchronization
- Clique detection
- Group membership
- Startup and Restart

... Which in combination provide services to the components (e.g. SW, subsystems):

- · Partitioned Messaging
- Scheduled Execution



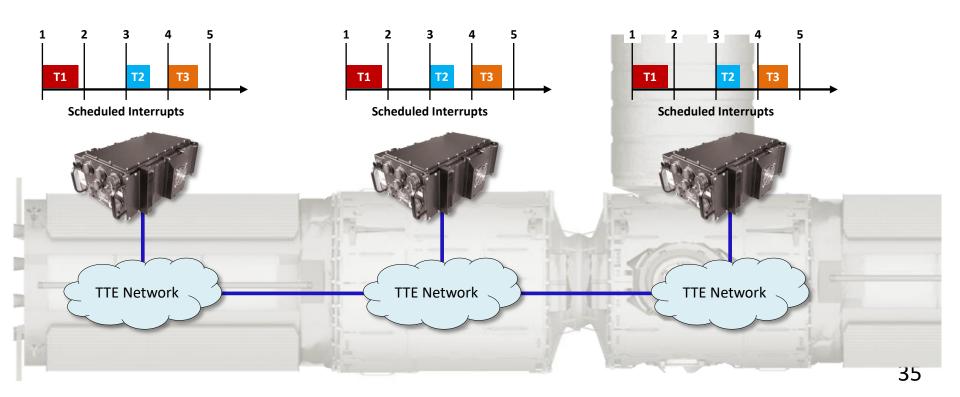
... That guarantee certain properties about the integrated system:

- 1. Composability
- 2. Compositionality



### **Scheduled Execution**

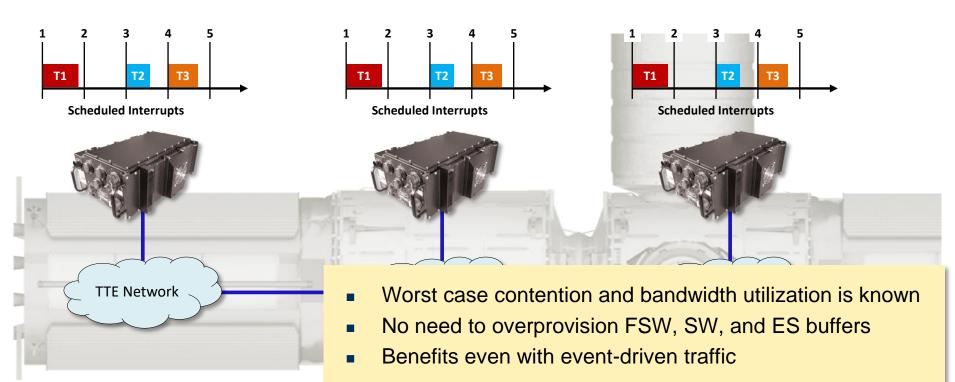
- Schedule interrupts based on global synchronized clock
- Cross-platform alignment between task scheduling and TT network scheduling



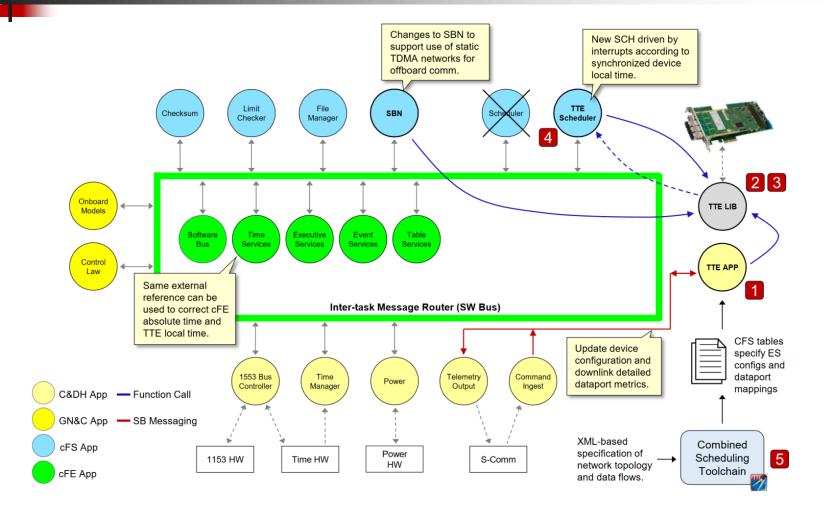


### **Scheduled Execution**

- Schedule interrupts based on global synchronized clock
- Cross-platform alignment between task scheduling and TT network scheduling



### Scheduled Execution: cFS

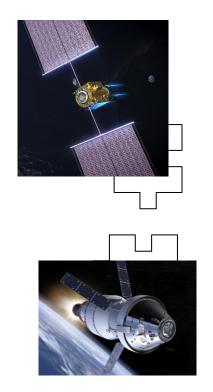


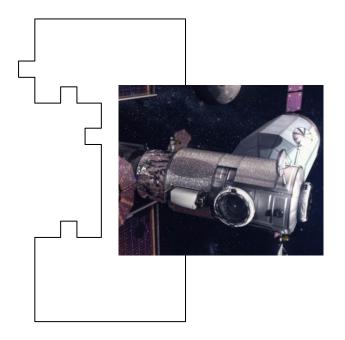


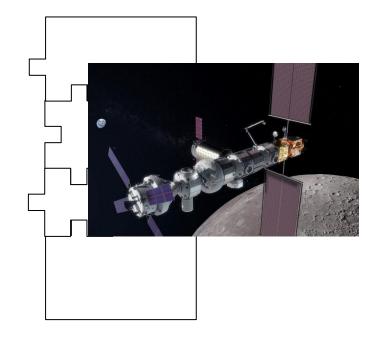
# Composable Systems

Because we know how these behave in isolation ...

We know how this behaves when fully integrated





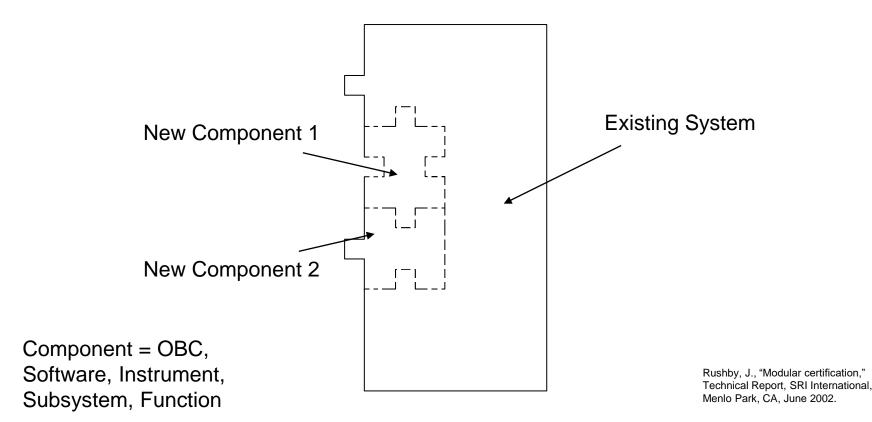


Module Integration



## Composable Systems

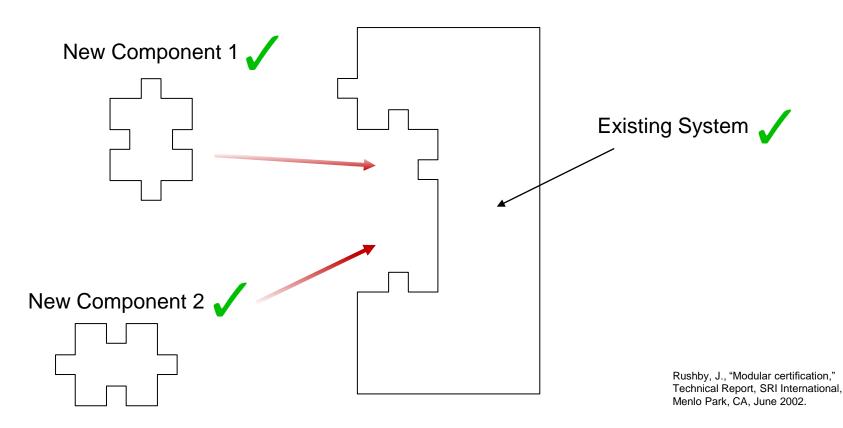
Certification traditionally requires us to view the system as an indivisible whole





## Composable Systems

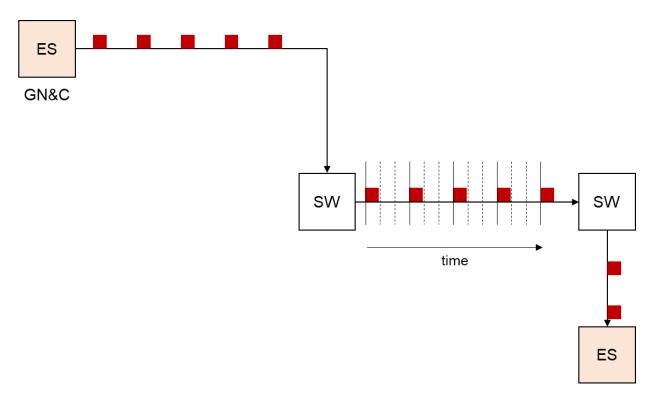
Desirable to certify the whole by integrating smaller certified components





## Why is it Hard?

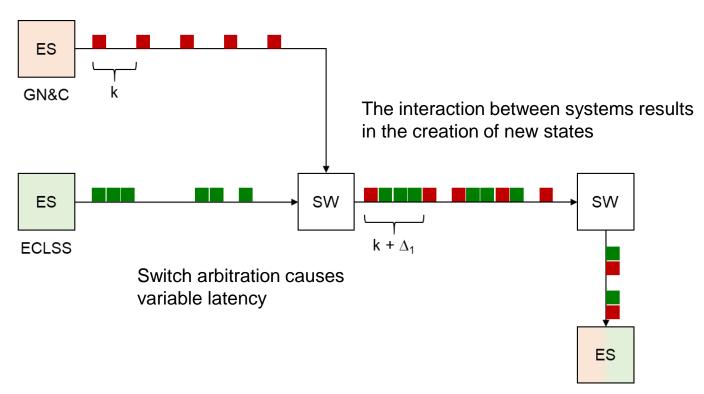
- Network traffic between modules can interact in unanticipated ways
- Makes it impossible to consider modules in isolation





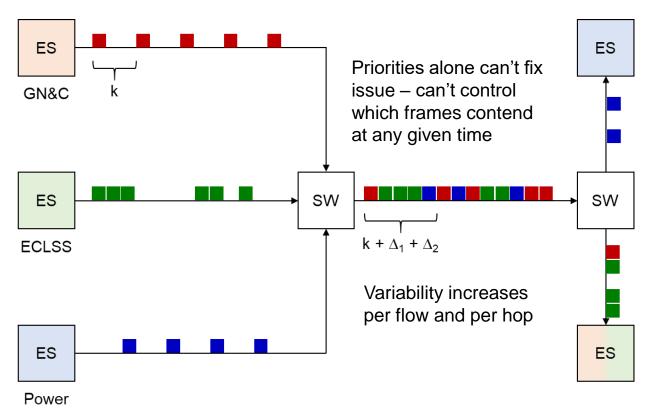
## Why is it Hard?

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# Why is it Hard?

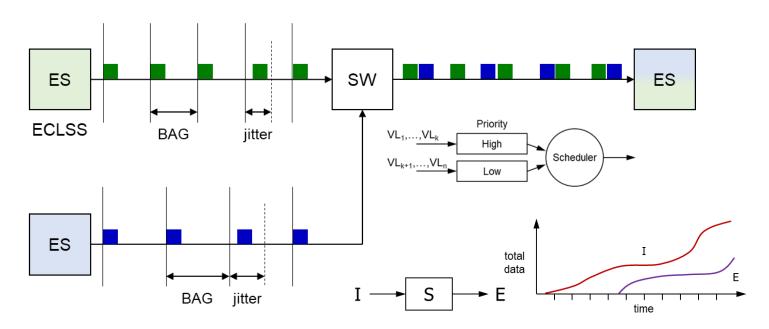
- Network traffic between modules can interact in unanticipated ways
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### One Approach: A664 P7

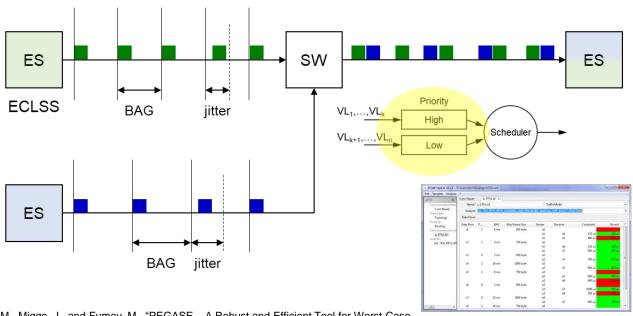
- **Bounding Variability:** ARINC 664 P7 can bound by mathematical proof:
  - 1) per VL end-to-end delay, 2) per port/VL jitter, 3) buffer sizes (from port waiting times)
- Use network calculus to calculate bounds on latency and queue sizes
  - Then use those metrics to determine VL priorities, BAGs, and frame sizes

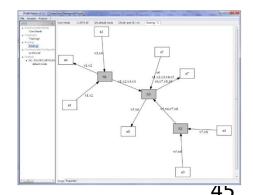




### One Approach: A664 P7

- But changing the priority or BAG of a given VL could have little impact, or huge impact
  - Depends not only on that VL, but <u>what it contends with</u>
- Therefore still requires consideration of the entire integrated system
- Increases complexity of maintaining a network that needs to grow/evolve over time

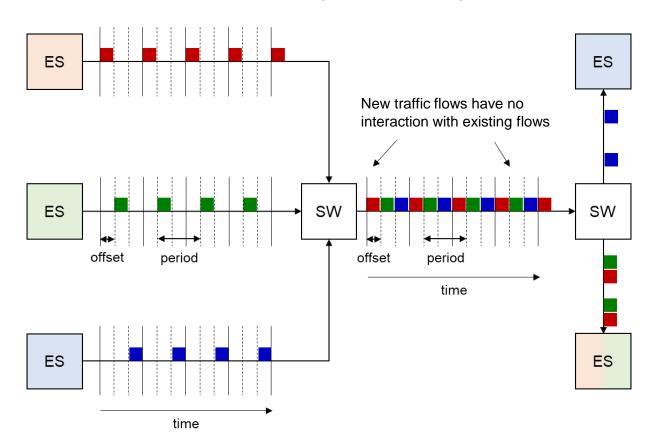






### Alternative Approach: TTE

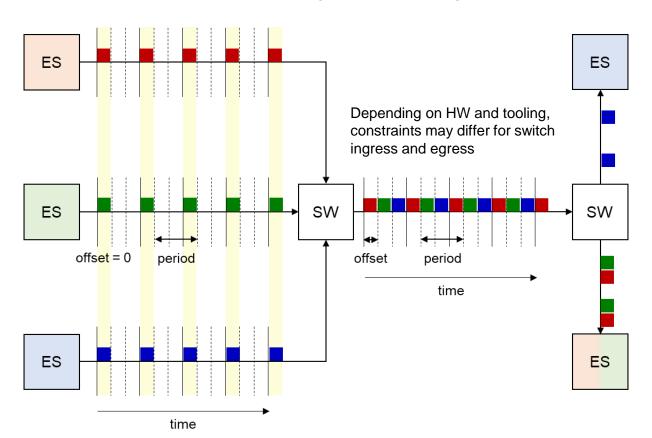
Prevent the contention that causes the timing variability (e.g. for buffers, ports)





# Alternative Approach: TTE

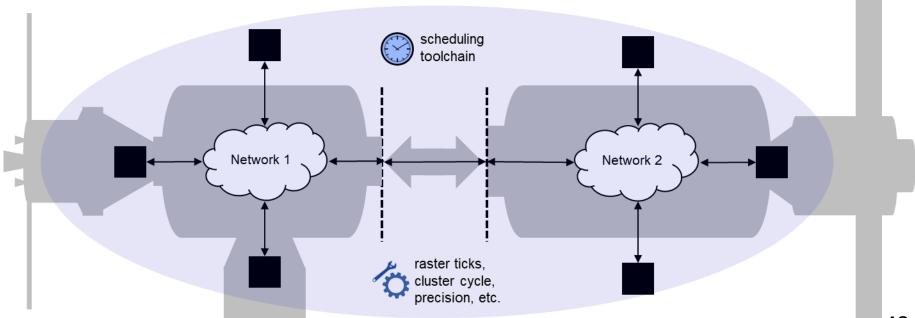
Prevent the contention that causes the timing variability (e.g. for buffers, ports)





# Incremental Buildup

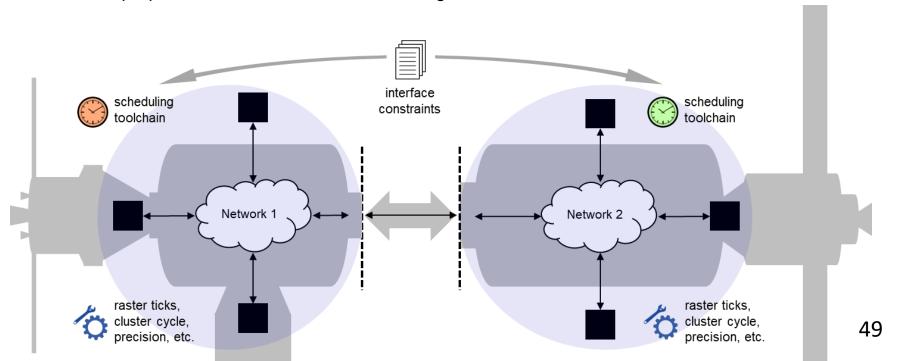
- One Option: Schedule the whole integrated vehicle at once
  - x Not composable schedule change in one module could necessitate change in another
  - x Requires the same tooling for all Gateway modules





# Incremental Buildup

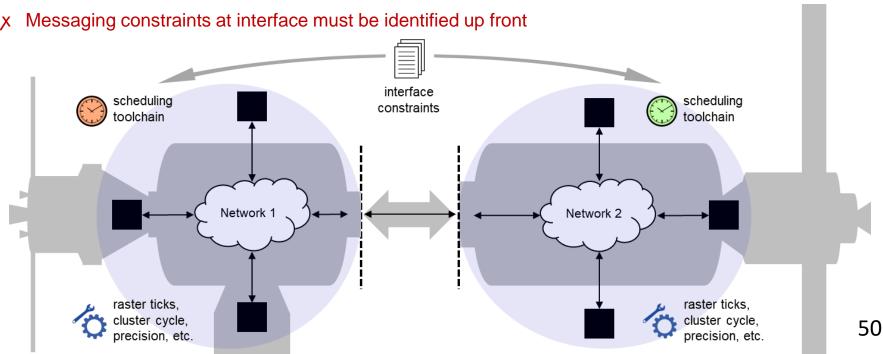
- Another Option: Vehicle has one network, scheduled piecewise with different tooling
  - Agree on consistent timing parameters for each module (e.g. raster granularity, cluster cycle, precision)
  - Define message properties at every module interface (e.g. direction, timing, max size)
  - Use these properties as constraints for scheduling each module





# Incremental Buildup

- Another Option: Vehicle has one network, scheduled piecewise with different tooling
  - ✓ Scheduling can be done separately per module (potentially with different tooling)
  - ✓ Scheduling within a module must only change if the interface constraints change
  - ✓ Allows different modules to be tested independently interactions are completely controlled.



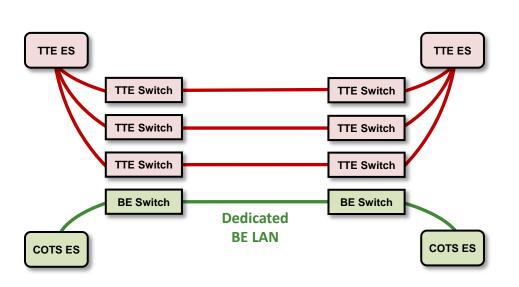
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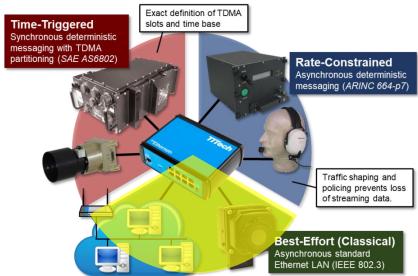
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#### **Best-Effort Ethernet**

- Gateway has no dedicated best-effort plane
  - Standard Ethernet/COTS devices connect to the same TTE network
  - Supported by TTE standard should be no problem

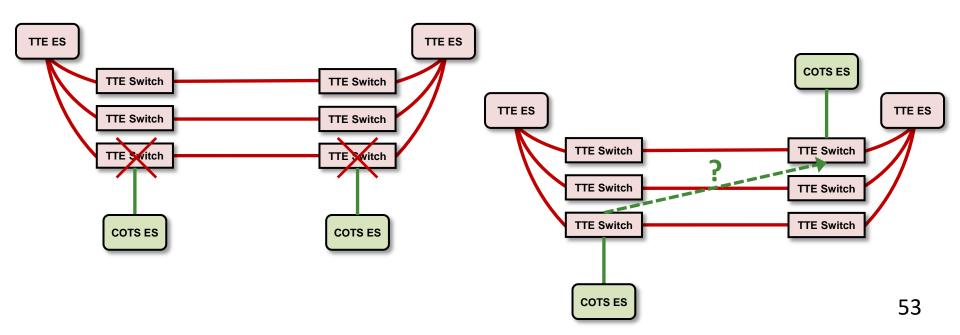






### **Best-Effort Ethernet**

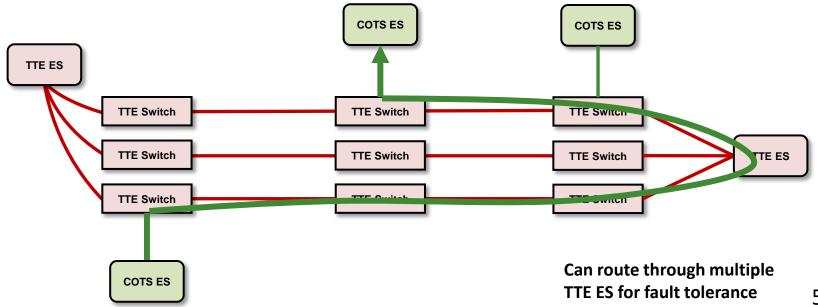
- A Challenge: Many COTS devices only connect to one plane
  - If all COTS devices are on same plane, you lose them after one fault
    - Some "non-critical" devices are critical. E.g. Cameras, crew laptops
  - But if COTS devices are on different planes, how do they talk?





### A Few Options

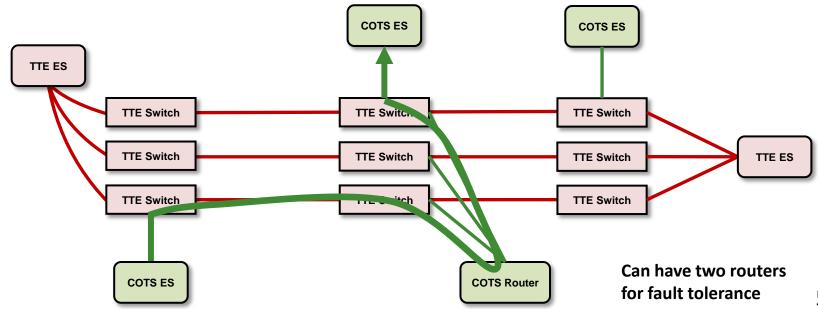
- Retransmit frames through a TTE end system
  - Pros: Adds no SWaP, needs no switch ports
  - Cons: Needs extra CPU time, SW changes, custom routing rules





### A Few Options

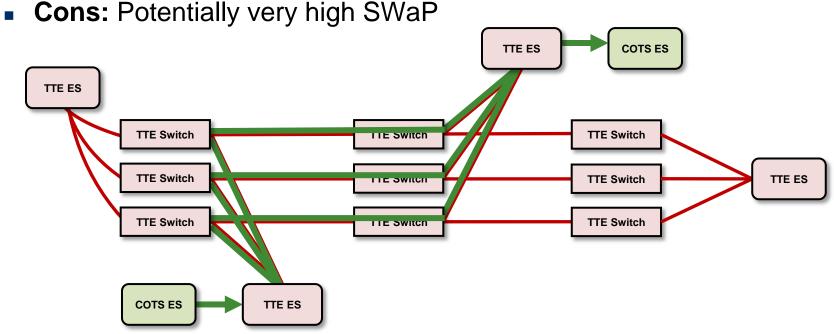
- Add a separate device to route frames
  - Pros: No SW changes, can use COTS routers
  - Cons: Adds SWaP, needs more switch ports, higher layers (L3+)





### A Few Options

- Connected devices through a remote interface unit (RIU)
  - Pros: Can convert BE to TT/RC, increases availability of BE, reduces switch ports on main backbone





#### **Best-Effort Ethernet**

- **Takeway:** No single best way to integrate COTS devices with a multi-plane TTE network
- Fact that TTE supports best-effort Ethernet is only <u>part</u> of the battle

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# Conclusion

- Gateway is based on a redundant TTE backbone
- TTE allows mixed-criticality traffic to exist on same network
- TTE can be used to realize a reliable broadcast abstraction
- TTE should be thought of as an integration framework
- Using COTS devices on TTE is not fool proof

# Sources

#### Content

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- https://ntrs.nasa.gov/archive/nasa/casi.ntrs.nasa.gov/20170001652.pdf
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#### Images

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