Recursion, Induction, and Iteration

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Program Validation and Verification in PVS
CADE 2021



Outline

Recursive Definitions

Induction Proofs

Induction-Free Induction

Recursive Judgements

Iterations

Inductive Definitions

Mutual Recursion and Higher-Order Recursion

Recursive Definitions in PVS

Suppose we want to define a function to sum the first n natural numbers:

$$\text{sum}(n) = \sum_{i=0}^{n} i.$$

In PVS:

```
sum(n): RECURSIVE nat =
      IF n = 0 THEN 0 ELSE n + sum(n - 1) ENDIF
MEASURE n
```

Recursive Definitions in PVS

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Functions in PVS are Total

Two Type Correctness Conditions(TCCs):

- The argument for the recursive call is a natural number.

```
%  Subtype TCC generated for n - 1
    %  expected type nat
sum_TCC1: OBLIGATION FORALL (n: nat):
    NOT n = 0 IMPLIES n - 1 >= 0;
```

- The recursion terminates.

```
%  Termination TCC generated for sum(n - 1)
sum_TCC2: OBLIGATION FORALL (n: nat):
    NOT n = 0 IMPLIES n - 1 < n;
```

Functions in PVS are Total

Two Type Correctness Conditions(TCCs):

- The argument for the recursive call is a natural number.

```
%  Subtype TCC generated for n - 1
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%  Termination TCC generated for sum(n - 1)
sum_TCC2: OBLIGATION FORALL (n: nat):
    NOT n = 0 IMPLIES n - 1 < n;
```

A Simple Property of Sum

We would like to prove the following closed form solution to sum:

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}.$$

In PVS:

```
closed_form: THEOREM
    sum(n) = (n * (n + 1)) / 2
```

Induction Proofs

```
>> (help "induct")
(induct$ var & optional (fnum 1) name) :
    Selects an induction scheme according to the type of VAR
    in FNUM and uses formula FNUM to formulate an induction
    predicate, then simplifies yielding base and induction
    cases. The induction scheme can be explicitly supplied
    as the optional NAME argument.
```

Induction Schemes from the Prelude

```
%  Weak induction on naturals.
nat_induction: LEMMA
    (p(0) AND (FORALL j: p(j) IMPLIES p(j+1)))
      IMPLIES (FORALL i: p(i))

%  Strong induction on naturals.
NAT_induction: LEMMA
    (FORALL j: (FORALL k: k < j IMPLIES p(k)) IMPLIES p(j))
      IMPLIES (FORALL i: p(i))
```

Proof by Induction

```
closed_form :

  |-------
{1} FORALL (n: nat): sum(n) = (n * (n + 1)) / 2

>> (induct "n")
Inducting on n on formula 1,
this yields 2 subgoals:
```

Proof by Induction

```
closed_form :

  |-------
{1} FORALL (n: nat): sum(n) = (n * (n + 1)) / 2

>> (induct "n")
Inducting on n on formula 1,
this yields 2 subgoals:
```

Base Case

```
closed_form.1 :

  |-------
{1} sum(0) = (0 * (0 + 1)) / 2

>> (grind)
 Rewriting with sum
 Trying repeated skolemization, instantiation, and if-lifting,

 This completes the proof of closed_form.1.
```

Base Case

```
closed_form.1 :

  |-------
{1} sum(0) = (0 * (0 + 1)) / 2

>> (grind)
 Rewriting with sum
 Trying repeated skolemization, instantiation, and if-lifting,

 This completes the proof of closed_form.1.
```

```
closed_form.2 :

  |-------
{1} FORALL j:
    sum(j) = (j * (j + 1)) / 2 IMPLIES
      sum(j + 1) = ((j + 1) * (j + 1 + 1)) / 2

>> (skeep)

Skolemizing with the names of the bound variables,
this simplifies to:
closed_form.2 :

{-1} sum(j) = (j * (j + 1)) / 2
  |-------
{1} sum(j + 1) = ((j + 1) * (j + 1 + 1)) / 2
```

```
closed_form.2 :

  |-------
{1} FORALL j:
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```

```
{-1} sum(j) = (j * (j + 1)) / 2
  |-------
{1} sum(j + 1) = ((j + 1) * (j + 1 + 1)) / 2

>> (expand "sum" +)
Expanding the definition of sum,
this simplifies to:
closed_form.2 :

[-1] sum(j) = (j * (j + 1)) / 2
  |-------
{1} 1 + sum(j) + j = (2 + j + (j * j + 2 * j)) / 2
```

```
{-1} sum(j) = (j * (j + 1)) / 2
  |-------
{1} sum(j + 1) = ((j + 1) * (j + 1 + 1)) / 2

>> (expand "sum" +)
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```

```
[-1] sum(j) = (j * (j + 1)) / 2
  |-------
{1} 1 + sum(j) + j = (2 + j + (j * j + 2 * j)) / 2

>> (assert)
Simplifying, rewriting, and recording with decision
procedures,

This completes the proof of closed_form.2.

Q.E.D.
```

```
[-1] sum(j) = (j * (j + 1)) / 2
  |-------
{1} 1 + sum(j) + j = (2 + j + (j * j + 2 * j)) / 2

>> (assert)
Simplifying, rewriting, and recording with decision
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Q.E.D.
```

Automated Simple Induction Proofs

```
  |-------
{1} FORALL (n: nat): sum(n) = (n * (n + 1)) / 2

>> (induct-and-simplify "n")
Rewriting with sum
Rewriting with sum
By induction on n, and by repeatedly rewriting and
simplifying,
Q.E.D.
```

Automated Simple Induction Proofs

```
  |-------
{1} FORALL (n: nat): sum(n) = (n * (n + 1)) / 2

>> (induct-and-simplify "n")
Rewriting with sum
Rewriting with sum
By induction on n, and by repeatedly rewriting and
simplifying,
Q.E.D.
```

Limitations of automation

Consider the nth factorial:

$$n! = \begin{cases} 1, & \text{if } n = 0 \\ n(n-1)!, & \text{otherwise.} \end{cases}$$

In the NASA PVS theory `ints@factorial`:

```
factorial(n:nat): RECURSIVE posnat =
    IF n = 0 THEN 1 ELSE n * factorial(n - 1) ENDIF
MEASURE n
```

A Simple Property of Factorial

$$\forall n : n! \geq n$$

In PVS:

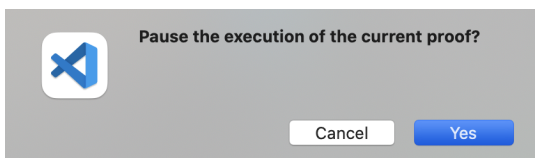
```
factorial_ge : LEMMA
    FORALL (n:nat): factorial(n) >= n
```

A Series of Unfortunate Events . . .

```
>> (induct-and-simplify "n")
Rewriting with factorial
Rewriting with factorial
Rewriting with factorial
Warning: Rewriting depth = 50; Rewriting with factorial
Warning: Rewriting depth = 100; Rewriting with factorial
. . .
```

Whenever the theorem prover falls into an infinite loop, `Ctrl+c` will automatically interrupt PVS and restore to the previous state

```
Executing (induct-and-simplify "n")
To interrupt the execution, press Ctrl+c
```



Factorial in C

Consider a common implementation of the n-th factorial in an imperative programming language:

```
/* Pre: n >= 0 */

int a = 1;
for (int i=0;i < n;i++) {
  /* Inv: a = i!  */
  a = a*(i+1);
}

/* Post: a = n! */
```

Factorial in C

Consider a common implementation of the n-th factorial in an imperative programming language:

```
/* Pre: n >= 0 */

int a = 1;
for (int i=0;i < n;i++) {
  /* Inv: a = i!  */
  a = a*(i+1);
}

/* Post: a = n! */
```

In PVS . . .

```
fact_it(n:nat,i:upto(n),a:posnat) : RECURSIVE posnat =
    IF i = n THEN a
    ELSE fact_it(n,i+1,a*(i+1))
    ENDIF
    MEASURE n-i

fact_it_correctness : THEOREM
    fact_it(n,0,1) = factorial(n)
```

Proving fact_it_correctness

```
  |-------
{1} FORALL (n: nat): fact_it(n, 0, 1) = factorial(n)

>> (induct-and-simplify "n")
this simplifies to:
fact_it_correctness :

{-1} fact_it(j!1, 0, 1) = factorial(j!1)
  |-------
{1} fact_it(1 + j!1, 1, 1) =
    factorial(j!1) + factorial(j!1) * j!1
```

The proof by (explicit) induction requires an inductive proof of an auxiliary lemma.

Proving fact_it_correctness

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{1} FORALL (n: nat): fact_it(n, 0, 1) = factorial(n)

>> (induct-and-simplify "n")
this simplifies to:
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```

The proof by (explicit) induction requires an inductive proof of an auxiliary lemma.

Induction-Free Induction By Predicate Subtyping

```
fact_it(n:nat,i:upto(n),(a:posnat|a=factorial(i))) :
    RECURSIVE {b:posnat | b=factorial(n)} =
    IF i = n THEN a
    ELSE fact_it(n,i+1,a*(i+1))
    ENDIF
    MEASURE n-i

n : VAR nat

fact_it_correctness : LEMMA
    fact_it(n,0,1) = factorial(n)
% |- fact_t_correctness : PROOF (skeep) (assert) QED
```

Induction-Free Induction By Predicate Subtyping

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```

There is No Free Lunch

```
fact_it_TCC4 :
  |-------
{1} FORALL (n: nat, i: upto(n),
    (a: nat | a = factorial(i))):
    NOT i = n IMPLIES a * (i + 1) = factorial(1 + i)
>> (skeep :preds? t)

fact_it_TCC4 :
{-1} n >= 0
{-2} i <= n
{-3} a = factorial(i)
  |-------
{1} i = n
{2} a * (i + 1) = factorial(1 + i)
```

There is No Free Lunch

```
fact_it_TCC4 :
  |-------
{1} FORALL (n: nat, i: upto(n),
    (a: nat | a = factorial(i))):
    NOT i = n IMPLIES a * (i + 1) = factorial(1 + i)
>> (skeep :preds? t)

fact_it_TCC4 :
{-1} n >= 0
{-2} i <= n
{-3} a = factorial(i)
  |-------
{1} i = n
{2} a * (i + 1) = factorial(1 + i)
```

```
>> (expand "factorial" 2)
fact_it_TCC4 :

[-1] n >= 0
[-2] i <= n
[-3] a = factorial(i)
  |-------
[1] i = n
{2} a * i + a = factorial(i) + factorial(i) * i

>> (assert)
Q.E.D.
```

```
>> (expand "factorial" 2)
fact_it_TCC4 :

[-1] n >= 0
[-2] i <= n
[-3] a = factorial(i)
  |-------
[1] i = n
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>> (assert)
Q.E.D.
```

```
>> (expand "factorial" 2)
fact_it_TCC4 :

[-1] n >= 0
[-2] i <= n
[-3] a = factorial(i)
  |-------
[1] i = n
{2} a * i + a = factorial(i) + factorial(i) * i

>> (assert)
Q.E.D.
```

You Can Also Pay at the Exit

```
fact_it2(n:nat,i:upto(n),a:posnat) : RECURSIVE
    {b:posnat | b = a*factorial(n)/factorial(i)} =
    IF i = n THEN a
    ELSE fact_it2(n,i+1,a*(i+1))
    ENDIF
    MEASURE n-i

fact_it2_correctness : LEMMA
    fact_it2(n,0,1) = factorial(n)
```

You Can Also Pay at the Exit

```
fact_it2(n:nat,i:upto(n),a:posnat) : RECURSIVE
    {b:posnat | b = a*factorial(n)/factorial(i)} =
    IF i = n THEN a
    ELSE fact_it2(n,i+1,a*(i+1))
    ENDIF
    MEASURE n-i

fact_it2_correctness : LEMMA
    fact_it2(n,0,1) = factorial(n)
```

```
  |-------
{1} FORALL (n: nat): fact_it2(n, 0, 1) = factorial(n)

>> (skeep)

  |-------
{1} fact_it2(n, 0, 1) = factorial(n)

>> (typepred "fact_it2(n,0,1)")

{-1} fact_it2(n, 0, 1) > 0
{-2} fact_it2(n, 0, 1) = 1 * factorial(n) / factorial(0)
  |-------
[1] fact_it2(n, 0, 1) = factorial(n)
```

```
  |-------
{1} FORALL (n: nat): fact_it2(n, 0, 1) = factorial(n)

>> (skeep)

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  |-------
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```
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  |-------
[1] fact_it2(n, 0, 1) = factorial(n)
```

```
>> (expand "factorial" -2 2)

[-1] fact_it2(n, 0, 1) > 0
{-2} fact_it2(n, 0, 1) = 1 * factorial(n) / 1
  |-------
[1] fact_it2(n, 0, 1) = factorial(n)

>> (assert)
Q.E.D.
```

```
>> (expand "factorial" -2 2)

[-1] fact_it2(n, 0, 1) > 0
{-2} fact_it2(n, 0, 1) = 1 * factorial(n) / 1
  |-------
[1] fact_it2(n, 0, 1) = factorial(n)

>> (assert)
Q.E.D.
```

But The Price is Higher

```
fact_it2_TCC5: OBLIGATION
    FORALL (n: nat, i: upto(n),
      v:
        [d1: z: [n: nat, upto(n), posnat] |
          z‘1 - z‘2 < n - i ->
        b: posnat | b = d1‘3 * factorial(d1‘1) /
        factorial(d1‘2)],
      a: posnat):
    NOT i = n IMPLIES
      v(n, i + 1, a * (i + 1)) =
      a * factorial(n) / factorial(i);
```

```
>> (skeep :preds? t)

Skolemizing with the names of the bound variables,
this simplifies to:
fact_it2_TCC5 :

{-1} n >= 0
{-2} i <= n
{-3} a > 0
  |-------
{1} i = n
{2} v(n, i + 1, a * (i + 1)) = a * factorial(n) / factorial(i)
```

```
>> (skeep :preds? t)

Skolemizing with the names of the bound variables,
this simplifies to:
fact_it2_TCC5 :

{-1} n >= 0
{-2} i <= n
{-3} a > 0
  |-------
{1} i = n
{2} v(n, i + 1, a * (i + 1)) = a * factorial(n) / factorial(i)
```

```
>> (name-replace "HI" "v(n, i + 1, a * (i + 1))")
Using HI to name and replace v(n, i + 1, a * (i + 1)),
this yields 2 subgoals:
fact_it2_TCC5.1 :


[-1] n >= 0
[-2] i <= n
[-3] a > 0
  |-------
[1] i = n
{2} HI = a * factorial(n) / factorial(i)
```

```
>> (name-replace "HI" "v(n, i + 1, a * (i + 1))")
Using HI to name and replace v(n, i + 1, a * (i + 1)),
this yields 2 subgoals:
fact_it2_TCC5.1 :

[-1] n >= 0
[-2] i <= n
[-3] a > 0
  |-------
[1] i = n
{2} HI = a * factorial(n) / factorial(i)
```

```
>> (typepred "HI")
Adding type constraints for HI,
this simplifies to:
fact_it2_TCC5.1 :

{-1} HI > 0
{-2} HI = (factorial(n) * a +
    factorial(n) * a * i) / factorial(1 + i)
[-3] n >= 0
[-4] i <= n
[-5] a > 0
  |-------
[1] i = n
[2] HI = a * factorial(n) / factorial(i)
```

```
>> (typepred "HI")
Adding type constraints for HI,
this simplifies to:
fact_it2_TCC5.1 :

{-1} HI > 0
{-2} HI = (factorial(n) * a +
    factorial(n) * a * i) / factorial(1 + i)
[-3] n >= 0
[-4] i <= n
[-5] a > 0
  |-------
[1] i = n
[2] HI = a * factorial(n) / factorial(i)
```

```
>> (expand "factorial" -2 3)
Expanding the definition of factorial,
this simplifies to:
fact_it2_TCC5.1 :


[-1] HI > 0
{-2} HI =
    (factorial(n) * a + factorial(n) * a * i) /
    (factorial(i) + factorial(i) * i)
[-3] n >= 0
[-4] i <= n
[-5] a > 0
  |-------
[1] i = n
[2] HI = a * factorial(n) / factorial(i)
```

```
>> (expand "factorial" -2 3)
Expanding the definition of factorial,
this simplifies to:
fact_it2_TCC5.1 :

[-1] HI > 0
{-2} HI =
    (factorial(n) * a + factorial(n) * a * i) /
    (factorial(i) + factorial(i) * i)
[-3] n >= 0
[-4] i <= n
[-5] a > 0
  |-------
[1] i = n
[2] HI = a * factorial(n) / factorial(i)
```

```
>> (replaces -2)
 Iterating REPLACE,

this simplifies to:
fact_it2_TCC5.1 :

{-1} (factorial(n) * a + factorial(n) * a * i) /
    (factorial(i) + factorial(i) * i)
    > 0
{-2} n >= 0
{-3} i <= n
{-4} a > 0
  |-------
{1} i = n
{2} (factorial(n) * a + factorial(n) * a * i) /
    (factorial(i) + factorial(i) * i)
    = a * factorial(n) / factorial(i)
```

```
>> (replaces -2)
 Iterating REPLACE,

this simplifies to:
fact_it2_TCC5.1 :

{-1} (factorial(n) * a + factorial(n) * a * i) /
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    > 0
{-2} n >= 0
{-3} i <= n
{-4} a > 0
  |-------
{1} i = n
{2} (factorial(n) * a + factorial(n) * a * i) /
    (factorial(i) + factorial(i) * i)
    = a * factorial(n) / factorial(i)
```

```
>> (grind-reals)
Rewriting with pos_div_gt
Rewriting with cross_mult

Applying GRIND-REALS,

This completes the proof of fact_it2_TCC5.1.
```

- All the other subgoals are discharged by `(assert)`.

Induction-Free Induction

+ Induction scheme based on the recursive definition of the function not on the measure!
+ Proofs exploit type-checker power.
- Some TCCs look scary (but they are easy to tame).
- If you modify the definitions, the TCCs get re-arranged (be careful or you can lose your proof).
? Can this method be used when the recursive function was not originally typed that way?

Recursive Judgments

Consider the Ackermann function:

$$A(m.n) = \begin{cases} n+1, \text{ if } m = 0 \\ A(m-1,1), \text{ if } m > 0 \text{ and } n = 0 \\ A(m-1, A(m, n-1)), \text{ otherwise.} \end{cases}$$

In PVS:

```
ack(m,n) : RECURSIVE nat =
IF m = 0 THEN n+1
ELSIF n = 0 THEN ack(m-1,1)
ELSE ack(m-1,ack(m,n-1))
ENDIF
MEASURE ?
```

Recursive Judgments

Consider the Ackermann function:

$$A(m.n) = \begin{cases} n+1, \text{ if } m = 0 \\ A(m-1,1), \text{ if } m > 0 \text{ and } n = 0 \\ A(m-1, A(m, n-1)), \text{ otherwise.} \end{cases}$$

In PVS:

```
ack(m,n) : RECURSIVE nat =
IF m = 0 THEN n+1
ELSIF n = 0 THEN ack(m-1,1)
ELSE ack(m-1,ack(m,n-1))
ENDIF
MEASURE lex2(m,n)
```

Ackermann

Proving this fact:

$$\forall m, n : A(m, n) > m + n$$

by regular induction is not trivial: you may need two nested inductions!

Recursive Judgements

```
ack_gt_m_n : RECURSIVE JUDGEMENT ack(m,n) HAS_TYPE above(m+n)
```

The type checker generates TCCs corresponding to the recursive definition of the type-restricted version of ack, e.g.,

```
ack_gt_m_n_TCC1: OBLIGATION FORALL (m, n: nat): m=0
    IMPLIES n+1 > m+n;

ack_gt_m_n_TCC3: OBLIGATION
    FORALL (v: [d: [nat, nat] -> above(d`1+d`2)], m, n: nat):
    n=0 AND NOT m=0 IMPLIES v(m-1, 1) > m+n;

ack_gt_m_n_TCC7: OBLIGATION
    FORALL (v: [d: [nat, nat] -> above(d`1+d`2)], m, n: nat):
    NOT n=0 AND NOT m=0 IMPLIES v(m-1, v(m, n-1)) > m+n;
```

Recursive Judgements

```
ack_gt_m_n : RECURSIVE JUDGEMENT ack(m,n) HAS_TYPE above(m+n)
```

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PVS Automatically Uses Judgements

Most of these TCCs are automatically discharged by the type checker (in this case, all of them). Furthermore, the theorem prover automatically uses judgements:

```
ack_simple_property :


  |-------
{1} FORALL (m, n): ack(m, n) > max(m, n)

>> (grind)
Rewriting with max
Trying repeated skolemization, instantiation, and if-lifting,
Q.E.D.
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Iterations

```
/* Pre: n >= 0  */
int a = 1;
for (int i=0;i < n;i++) {
  /* Inv: a = i!  */
  a = a*(i+1);
}
/* Post: a = n! */
```

In PVS:

```
IMPORTING structures@for_iterate

fact_for(n:nat) : real =
    for[real](0,n-1,1,LAMBDA(i:below(n),a:real): a*(i+1))
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Proving Correctness of Iterations

Consider the following implementation of factorial:

```
fact_for : THEOREM
fact_for(n) = factorial(n)

fact_for :


  |-------
{1} FORALL (n: nat): fact_for(n) = factorial(n)

>> (skeep)(expand "fact_for")
fact_for :


  |-------
{1} for[real](0,n-1,1,LAMBDA(i:below(n),a:real):a+a*i) =
    factorial(n)
```

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```

```
>> (lemma "for_induction[real]")
Applying for_induction[real]
this simplifies to:
fact_for :

{-1} FORALL (i, j: int, a: real, f: ForBody[real](i, j),
    inv: PRED[[UpTo[real](1 + j - i), real]]):
    (inv(0, a) AND
    (FORALL (k: subrange(0, j - i), ak: real):
    inv(k, ak) IMPLIES inv(k + 1, f(i + k, ak))))
    IMPLIES inv(j - i + 1, for(i, j, a, f))
  |-------
[1] for[real](0,n-1,1,LAMBDA (i:below(n),a:real):a+a*i) =
    factorial(n)
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>> (lemma "for_induction[real]")
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  |-------
[1] for[real](0,n-1,1,LAMBDA(i:below(n),a:real):a+a*i) =
    factorial(n)
```

```
>> (inst?)
Instantiating quantified variables,
this yields 2 subgoals:
fact_for.1 :

{-1} FORALL (inv:PRED[[UpTo[real](n)real]]):
    (inv(0,1) AND
    (FORALL (k:subrange(0,n-1),ak:real):
      inv(k,ak) IMPLIES inv(k+1,ak+ak*(0+k))))
      IMPLIES
        inv(n,
          for(0,n-1,1,LAMBDA(i:below(n),a:real):a+a*i))
  |-------
[1] for[real](0,n-1,1,LAMBDA (i:below(n),a:real):a+a*i) =
    factorial(n)
```

```
>> (inst?)
Instantiating quantified variables,
this yields 2 subgoals:
fact_for.1 :

{-1} FORALL (inv:PRED[[UpTo[real](n)real]]):
    (inv(0,1) AND
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      IMPLIES
        inv(n,
          for(0,n-1,1,LAMBDA(i:below(n),a:real):a+a*i))
  |-------
[1] for[real](0,n-1,1,LAMBDA (i:below(n),a:real):a+a*i) =
    factorial(n)
```

```
>> (inst -1 "LAMBDA(i:upto(n),a:real) : a = factorial(i)")
fact_for.1.1 :

{-1} . . .
  |-------
[1] for[real](0,n-1,1,LAMBDA(i:below(n),a:real):a+a*i) =
    factorial(n)
```

- The variable *i* in the invariant refers to the *i*th iteration.
- Remaining subgoals are discharged with `(grind)`.
  - See `Examples/Lecture-2.pvs`.

Inductive Definitions

- An inductive definition gives rules for generating members of a set.
- An object is in the set, only if it has been generated according to the rules.
- An inductively defined set is the smallest set closed under the rules.
- PVS automatically generates weak and strong induction schemes that are used by the command `(rule-induct "<name>")`

Even and Odd

```
even(n:nat): INDUCTIVE bool =
    n = 0 OR (n > 1 AND even(n - 2))

odd(n:nat): INDUCTIVE bool =
    n = 1 OR (n > 1 AND odd(n - 2))
```

Induction Schemes

The definition of even generates the following induction schemes:

```
even_weak_induction: AXIOM
FORALL (P: [nat -> boolean]):
    (FORALL (n: nat): n = 0 OR (n > 1 AND P(n - 2))
      IMPLIES P(n))
        IMPLIES
          (FORALL (n: nat): even(n) IMPLIES P(n));

even_induction: AXIOM
    FORALL (P: [nat -> boolean]):
      (FORALL (n: nat):
        n = 0 OR (n > 1 AND even(n - 2) AND P(n - 2))
          IMPLIES P(n))
            IMPLIES (FORALL (n: nat): even(n) IMPLIES P(n));
```

Inductive Proof

```
even_odd :

  |-------
{1} FORALL (n: nat): even(n) => odd(n + 1)

>> (rule-induct "even")
Applying rule induction over even, this simplifies to:
even_odd :

  |-------
{1} FORALL (n: nat):
    n = 0 OR (n > 1 AND odd(n - 2 + 1)) IMPLIES odd(n + 1)
```

The proof can then be completed using

```
(skosimp*)(rewrite "odd" +)(ground)
```

Inductive Proof

```
even_odd :

  |-------
{1} FORALL (n: nat): even(n) => odd(n + 1)

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Applying rule induction over even, this simplifies to:
even_odd :

  |-------
{1} FORALL (n: nat):
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Inductive Proof

```
even_odd :

  |-------
{1} FORALL (n: nat): even(n) => odd(n + 1)

>> (rule-induct "even")
Applying rule induction over even, this simplifies to:
even_odd :

  |-------
{1} FORALL (n: nat):
    n = 0 OR (n > 1 AND odd(n - 2 + 1)) IMPLIES odd(n + 1)
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```

Mutual Recursion and Higher-Order Recursion

The predicates odd and even can be defined using a mutual-recursion:

$$
\begin{aligned}
\texttt{even?}(0) &= \texttt{true} \\
\texttt{odd?}(0) &= \texttt{false} \\
\texttt{odd?}(1) &= \texttt{true} \\
\texttt{even?}(n+1) &= \texttt{odd?}(n) \\
\texttt{odd?}(n+1) &= \texttt{even?}(n)
\end{aligned}
$$

In PVS . . .

```
my_even?(n) : INDUCTIVE bool =
n = 0 OR n > 0 AND my_odd?(n-1)

my_odd?(n) : INDUCTIVE bool =
n = 1 OR n > 1 AND my_even?(n-1)
```

- Theses definitions don't type-check. What is wrong with them?
- PVS does not (directly) support mutual recursion.

In PVS . . .

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my_even?(n) : INDUCTIVE bool =
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```

- Theses definitions don't type-check. What is wrong with them?
- PVS does not (directly) support mutual recursion.

Mutual Recursion via Higher-Order Recursion

```
even_f?(fodd:[nat->bool],n) : bool =
n = 0 OR
n > 0 AND fodd(n-1)

my_odd?(n) : INDUCTIVE bool =
n = 1 OR
n > 1 AND even_f?(my_odd?,n-1)

my_even?(n) : bool =
even_f?(my_odd?,n)
```

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```
even_f?(fodd:[nat->bool],n) : bool =
n = 0 OR
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