



## Determining the Most Influencing Medical Conditions in MEDPRAT's SIN Directed Graph

Crew Health and Performance
Probabilistic Risk Assessment Project
(CHP-PRA)

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Background and Introduction

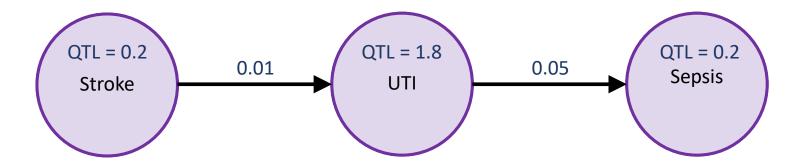


## SIN Graph



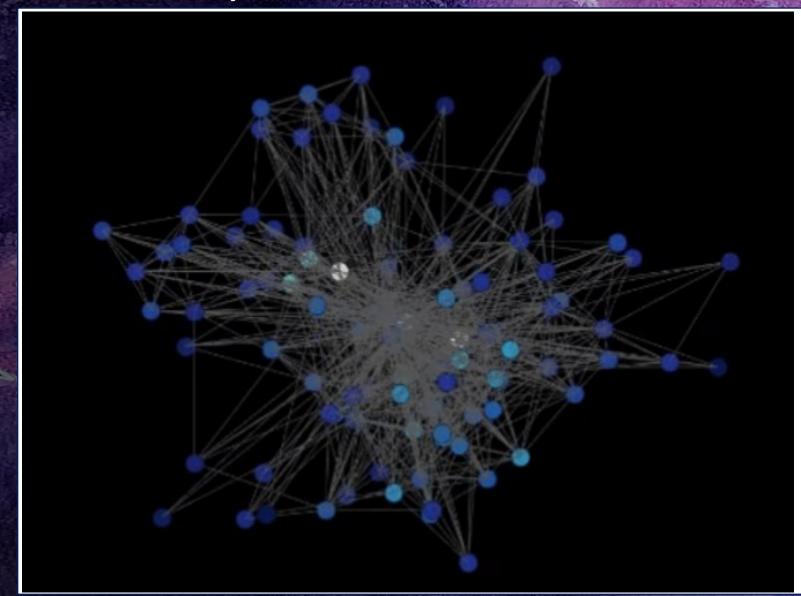
#### The Susceptibility Inference Network (SIN)

- Part of MEDPRAT, an event driven, time dependent, stochastic tool
- Subject matter expert informed prototype
- Mathematical data structure to quantify relationships between medical conditions:
  - o nodes: medical conditions
  - o edges: directed, progression from one condition to the next
  - o node weights: a statistic to measure the stand-alone risk, typically Quality Time Loss (QTL) or Loss Of Crew Life (LOCL)
  - o edge weights: probability of progression to the next condition



## SIN Graph





- 99 medical conditions
- 1078 direct connections
- Anxiety, Atrial Fibrillation/ Flutter, Medication Overdose/Adverse Reaction responsible for 273 of the connections



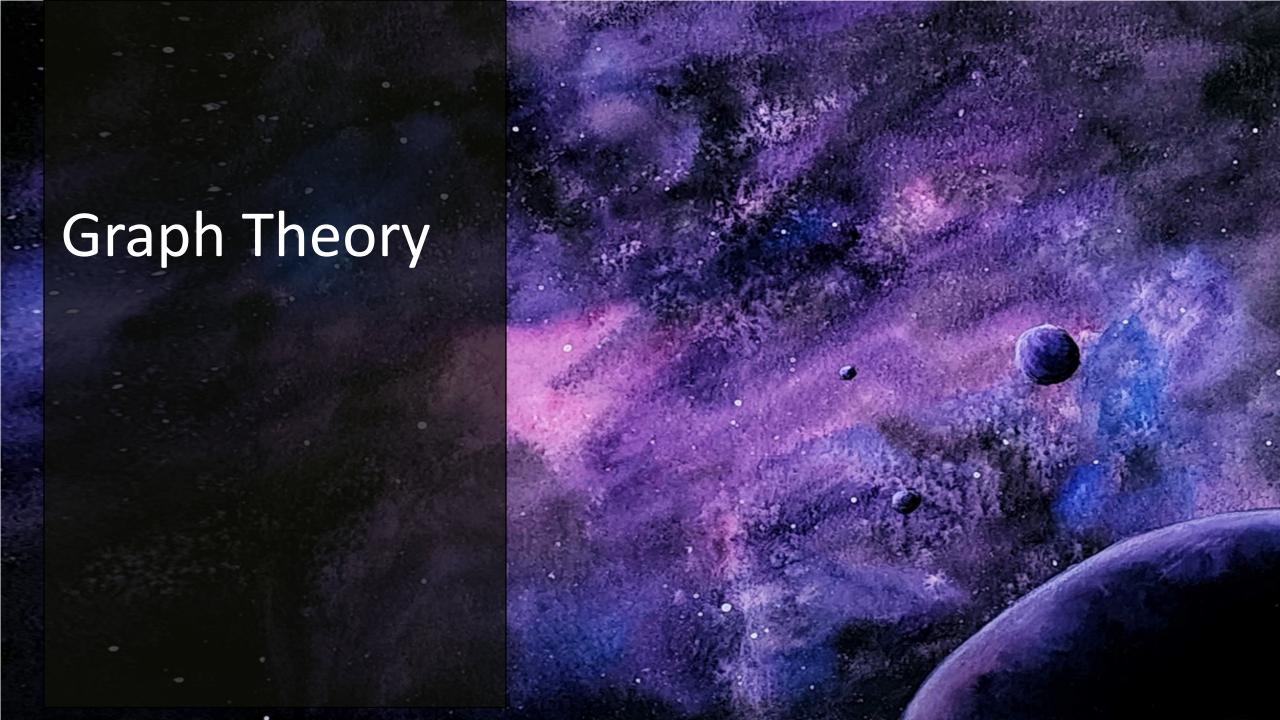
## Centrality Measures



• Our goal 1: compare how much each medical condition can progress through the network

- Measure of importance of a node
  - Local measures: degree centrality
  - Global measures: betweenness, eigenvector, Katz, PageRank Centrality

We use: Katz Centrality



## SIN Adjacency Matrix

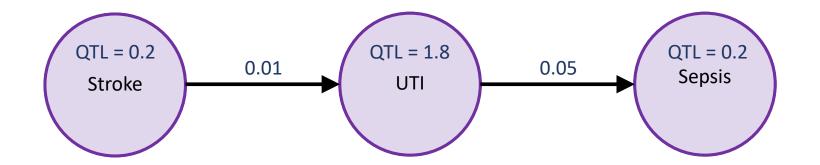


#### SIN adjacency matrix

Stand-alone risks

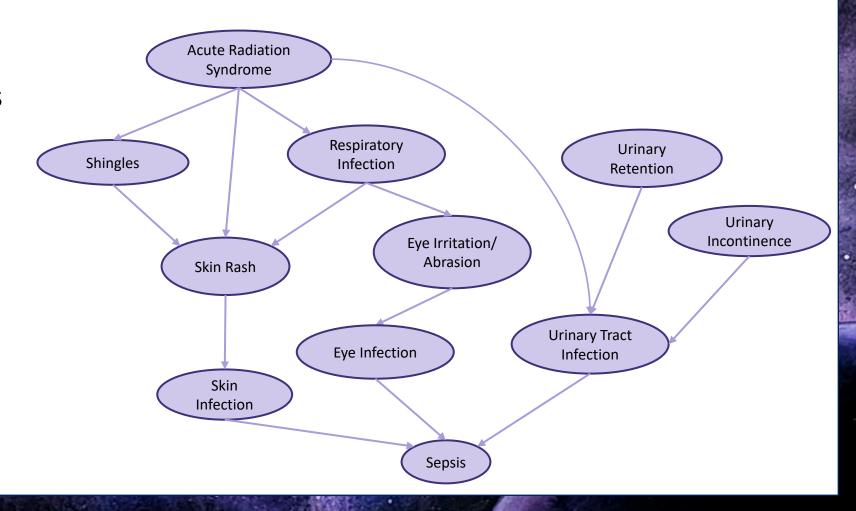
Stroke 
$$\begin{bmatrix} 0.2 \\ \end{bmatrix}$$

$$W = \begin{bmatrix} 0.2 \\ \end{bmatrix}$$
Sepsis  $\begin{bmatrix} 0.2 \\ \end{bmatrix}$ 



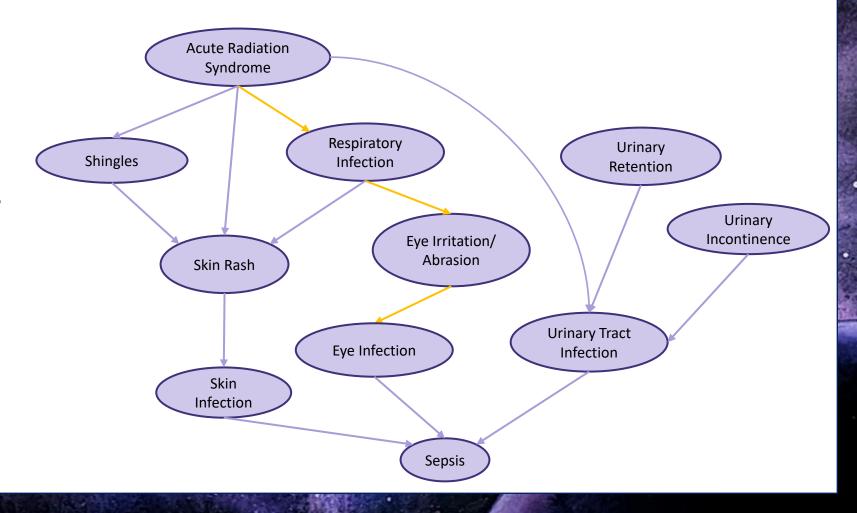


- ullet Walk of length k through k edges
- Weight of walk= product of edge weights



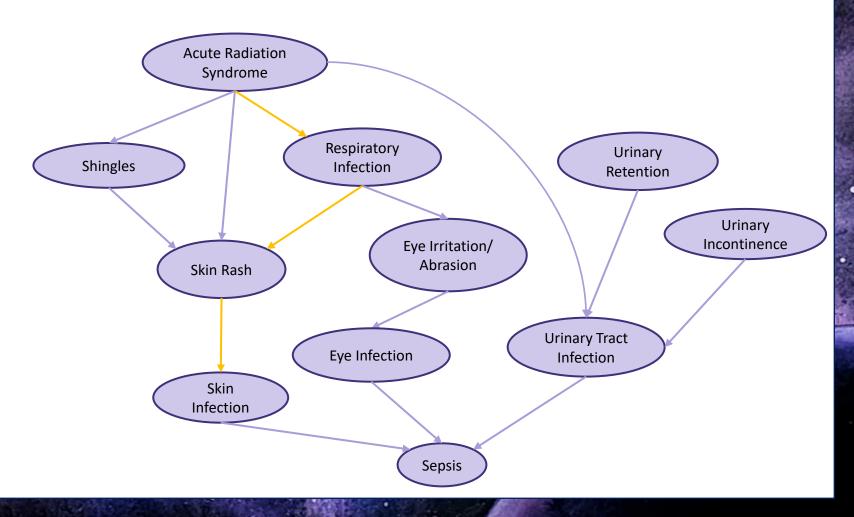


- Walk of length 3
- Weight of walk= product of edge weights



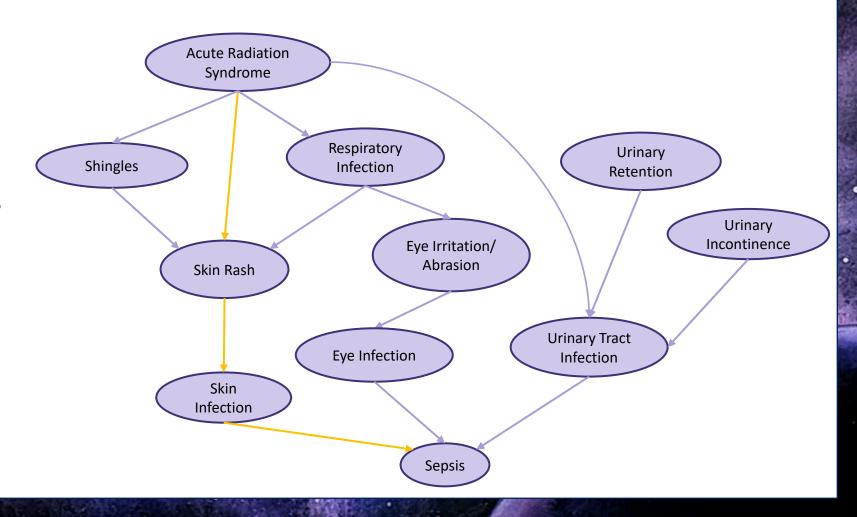


- Walk of length 3
- Weight of walk= product of edge weights



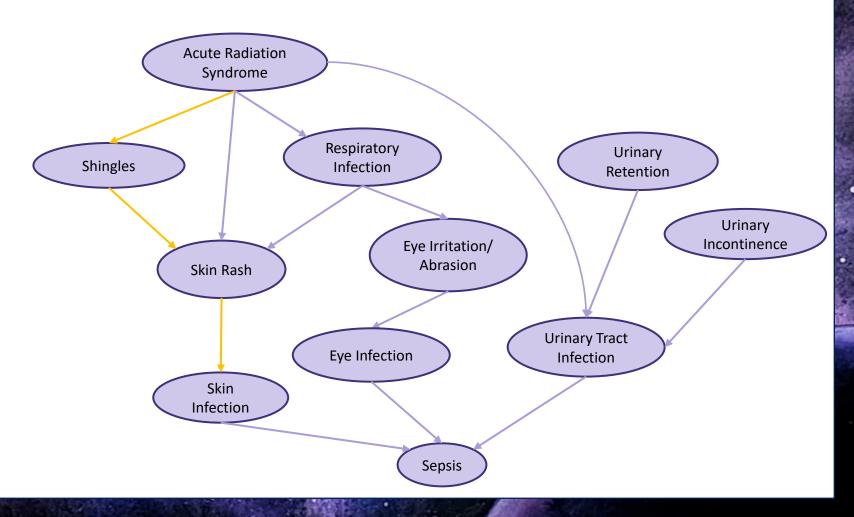


- Walk of length 3
- Weight of walkproduct of edge weights



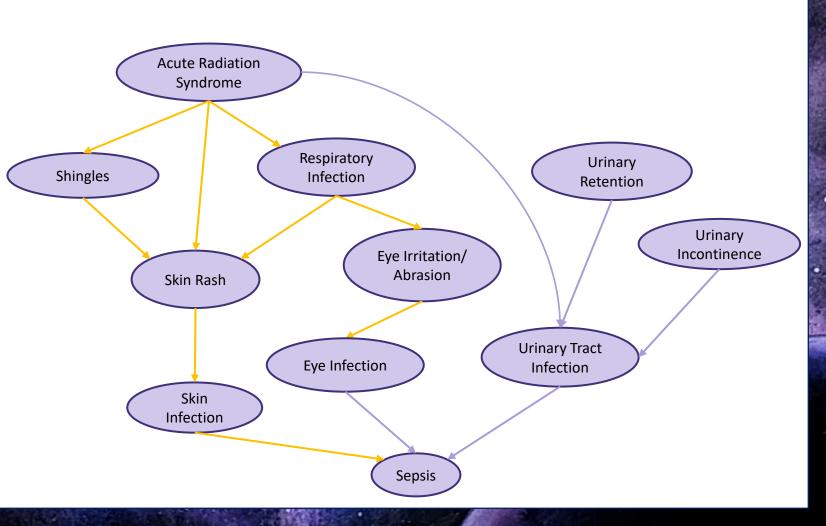


- Walk of length 3
- Weight of walk= product of edge weights





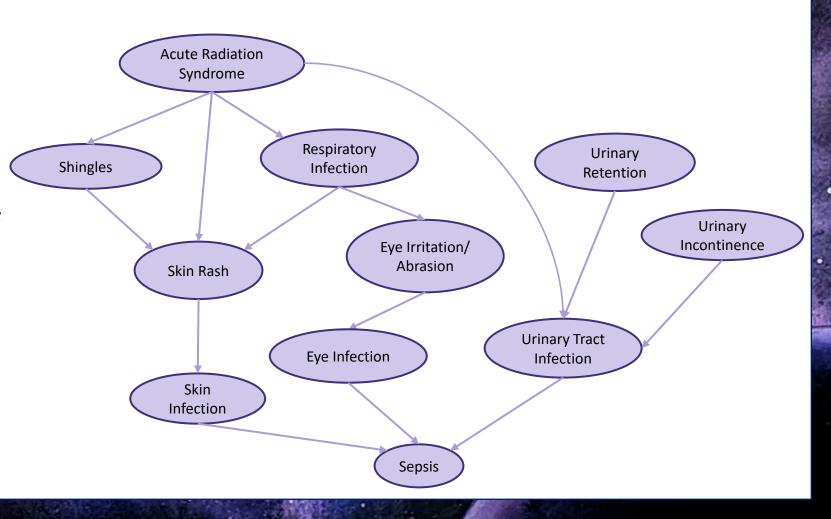
- All walks of length 3
- Weighted number of walks of length 3
  - = sum of weighted number of all walk of length 3





#### For adjacency matrix A

- Weighted number of walks of length k= $A^k \cdot 1$
- Sum of weighted number of walks of any length  $=(A + A^2 + A^3 + \cdots) \cdot 1$



### Katz Centrality



ullet The Katz Centrality of vertex  $oldsymbol{i}$  in a graph with adjacency matrix A:

$$C(\alpha)_{i} = (\alpha A + \alpha^{2} A^{2} + \alpha^{3} A^{3} + \cdots) \cdot W$$
$$= \left( (I - \alpha A)^{-1} - I \right) \cdot W$$

• Katz parameter  $\alpha$  to penalize longer walks

• Criteria to choose  $\alpha$  not well defined in the literature

#### Katz Parameter



• Our goal 2: give guidance to the choice of  $\alpha$ 

 Based on the maximum path length of interest L when navigating the network

• Takeaway: for maximum path length L and error tolerance  $\epsilon$ , choose  $\alpha$  that satisfies

$$\left[\log_{\alpha\|A\|_{2}}\left(\frac{\epsilon}{\|C(\alpha)\|_{2}}\right)\right] = L$$

#### Our propositions and theorems:

**Proposition 1** Let  $C \mid and C'$  be two centrality measures on a network N. If  $||C - C'||_{\infty} < \epsilon$ ,

then C and C'  $\epsilon$ -agree.

**Lemma 2** (Absolute Error Tolerance) Let  $p \in \{1, 2, \infty\}$  and  $\alpha \in (0, 1/\rho)$ . Then

$$||C(\alpha) - C(\alpha, \ell)||_{\infty} \le (\alpha ||A||_p)^{\ell} ||C(\alpha)||_p := \epsilon_{\ell}.$$

**Lemma 3** (Error Tolerance Guarantee) Let  $p \in \{1, 2, \infty\}$ ,  $\alpha \in (0, 1/\|A\|_p)$  and  $\epsilon > 0$ . If

$$\ell > \log_{\alpha \|A\|_p} \left( \frac{\epsilon}{\|C(\alpha)\|_p} \right) := L$$

then  $||C(\alpha) - C(\alpha, \ell)||_{\infty} < \epsilon$ .

Corollary 4 (Relative Error) Let G be a graph, A the adjacency matrix of G,  $\alpha = \alpha_0/\|A\|_2$ , and  $\epsilon = \epsilon_0\|c_\alpha\|_2$  where  $\alpha_0, \epsilon_0 \in (0,1)$ . Then  $\alpha$ -Katz centrality and  $(\alpha, \ell)$ -Katz centrality are  $\epsilon$ -tolerant whenever  $\ell > \log_{\alpha_0}(\epsilon_0) = L$ .

Corollary 5 (Local Katz Centrality) Let G be a directed graph,  $\alpha \in (0, 1/\|A\|_2)$ ,  $\epsilon > 0$ , and L be as defined in Lemma 2. Let H be the subgraphs formed by the L-hop neighborhood of v in G. Then the  $\alpha$ -Katz centrality of  $v \in G$  is within  $\epsilon$  of the  $\alpha$ -Katz centrality of  $v \in H$ .

# Application to SIN



## **Example Mission**



#### Given short mission details:

- 1. 6 months long
- 2. L = 6
- 3. Error tolerance  $\epsilon_0 = 0.001$

Use calculated  $\alpha = 0.37$ 

Rank (QTL)	Condition	Rank (LOCL)	Condition
1	Acute Radiation Syndrome	1	Acute Radiation Syndrome
2	Anaphylaxis	2	Burns Secondary to Fire
3	Allergic Reaction	3	Appendicitis
4	Toxic Exposure	4	Urinary Tract Infection
5	Abdominal Injury	5	Sudden Cardiac Arrest

## **Example Mission**



#### Given long mission details:

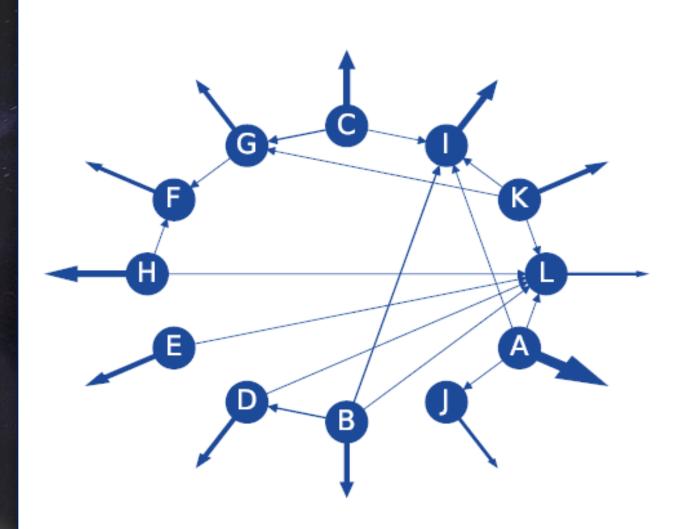
- 1. 3 years long
- 2. L = 15
- 3. Error tolerance  $\epsilon_0 = 0.001$

Use calculated  $\alpha = 0.80$ 

Rank (QTL)	Condition	Rank (LOCL)	Condition
1	Acute Radiation Syndrome	1	Acute Radiation Syndrome
2	Anaphylaxis	2	Burns Secondary to Fire
3	Toxic Exposure	3	Appendicitis
4	Allergic Reaction	4	Urinary Tract Infection
5	Abdominal Injury	5	Sudden Cardiac Arrest

#### Subnetwork of the SIN with Highest Centrality Nodes





For  $\alpha$ =**0**.37

- 12 top ranked conditions in the SIN in contributing to the QTL risk
- Thickness of edges pointing outwards reflecting total edge weight towards nodes in the remainder of the network
- Node A in first rank: Acute Radiation Syndrome

## Conclusions and Limitations



#### **Conclusions and Limitations**



#### **Conclusions**

- The nodes Katz centrality scores depend on Katz parameter  $\alpha$ .
- We recommend an  $\alpha$  based on the depth to navigate in the network.
- Medical conditions are ranked differently for different mission profiles, based on the number of progressions of concern during each mission.

#### **Limitations**

- $\alpha$  is bounded by the inverse of the spectral radius of the graph's adjacency matrix which limits the range of  $\alpha$  to consider.
- We can get smaller margin of error computationally than analytically.

