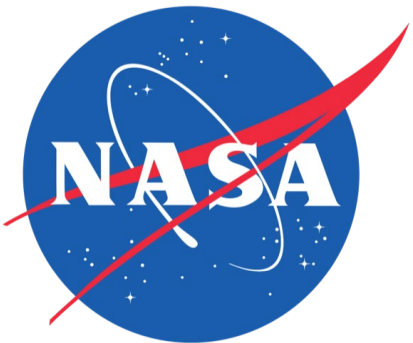

A Temporal Differential Dynamic Logic Formal Embedding

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Overview

- **dTL²**: Differential Temporal Dynamic Logic [1]
- **PVS**: Interactive theorem prover [2]

Result: Embedding of dTL² in PVS

- Extend the functionality of the embedding of **dL** in **PVS**, **Plaidypvs** [3] by allowing for temporal reasoning
- Fully operational in **PVS**



Plaidypvs: Properly Assured
Implementation of Differential
Dynamic Logic for Hybrid Program
Verification and Specification

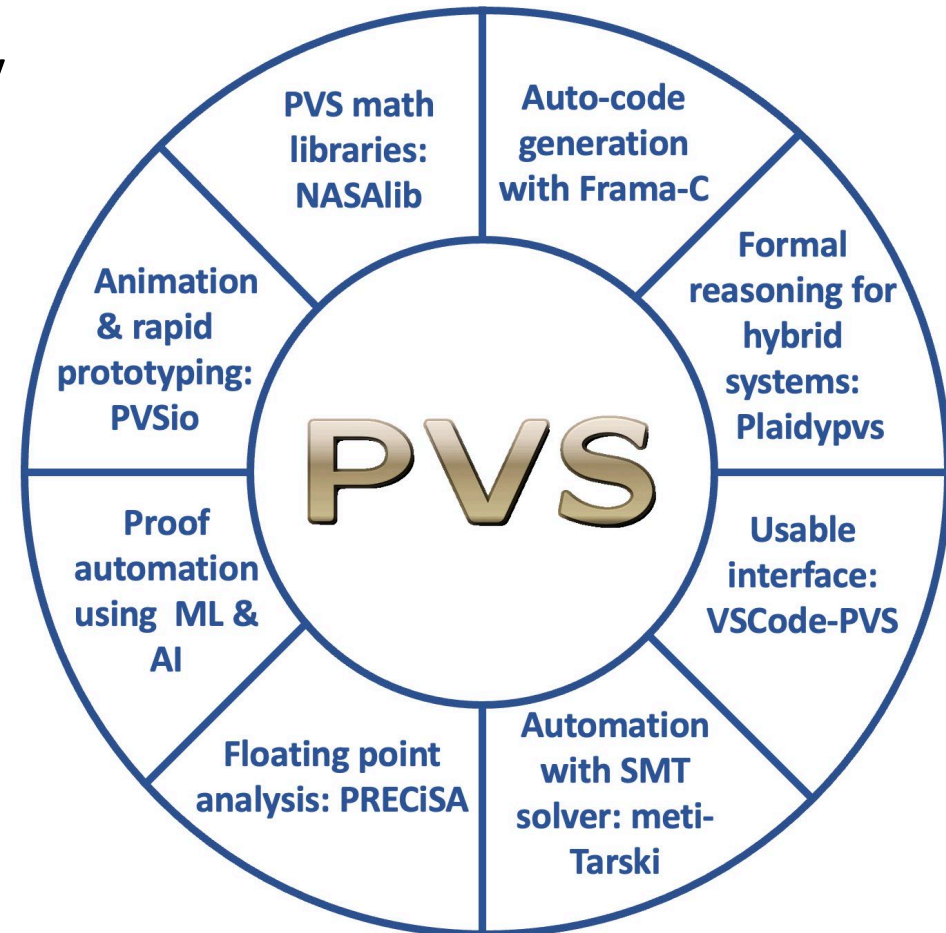
[1] Jeannin & Platzer. 2014. "dTL²: differential temporal dynamic logic with nested temporalities for hybrid systems." Springer International.

[2] PVS website, SRI International: <https://pvs.csl.sri.com>

[3] The embedding of dL in PVS, NASA's PVS library: <https://github.com/nasa/pvslib/dL>

PVS

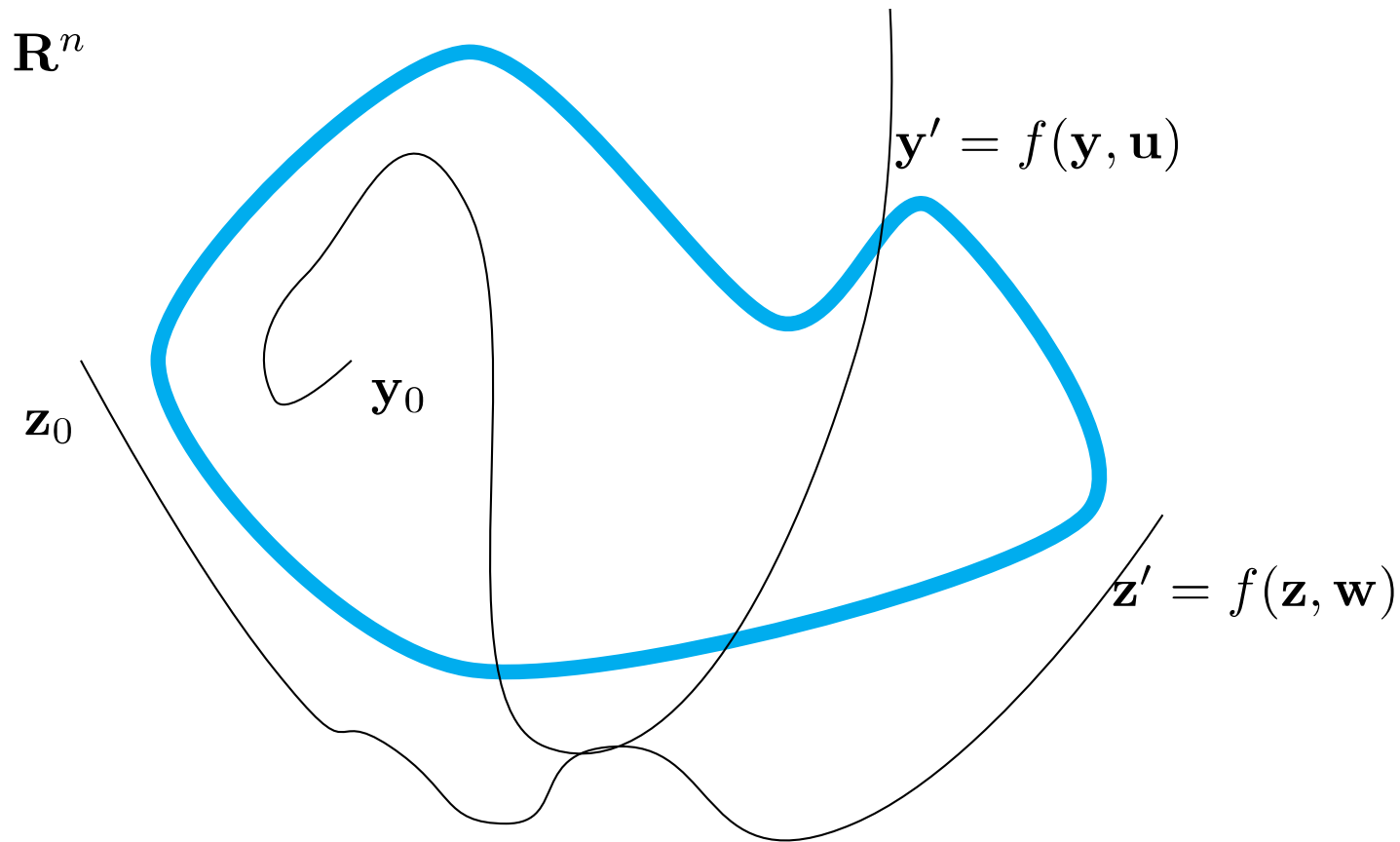
- “Prototype Verification System” developed by SRI International
- Interactive theorem prover
 - Higher order logic
 - Completely typed, dependent types
- Automation
 - Customizable tactics and strategies
- PVSio animation and rapid prototyping
- NASA PVS library [4]
 - 58 libraries with over 35,000 lemmas
- Visual studio code extension [5]



[4] NASALib, maintained by NASA Langley Formal Methods Group: <https://github.com/nasa/pvslib>

[5] VSCoDe-PVS, Paolo Masci: <https://github.com/nasa/vscode-pvs>

Plaidypvs and Hybrid Systems



- **Hybrid system:** dynamical system that exhibits
 - Continuous behavior
 - Discrete behavior

Plaidypvs:

- **Formal specification** of hybrid systems
- **Formal reasoning** of hybrid systems

Hybrid Programs

Hybrid programs allow **formal specification** of hybrid systems:

- Discrete jump set:
 $(x_1 := \theta_1, \dots, x_n := \theta_n)$
- Differential equations:
 $\{x'_1 := \theta_1, \dots, x'_n := \theta_n \ \& \ \chi\}$
 - $\{x_i\}_{i=1}^n$ variables
 - $\{\theta_i\}_{i=1}^n$ assignments (ex. – functions of past variable assignments)
 - χ first-order formula that describes the domain

For hybrid programs Hp_1, Hp_2 , first-order formula χ :

- Union
 $(Hp_1 \cup Hp_2)$
- Sequence
 $(Hp_1; Hp_2)$
- Repeat
 $(Hp_1)^*$
- Test
 $(? \chi)$

Hybrid Programs (continued)

Example: $\alpha := ((? (v < V_0); \{v' = M \ \&(v \leq V_0 + \tau)\}) \cup (? (v \geq V_0); \{v' = -M \ \&(v \geq V_0 - \tau)\}))^*$

The hybrid program α models a simple cruise controller.

When the velocity is less than V_0 , set the acceleration to M and when the velocity is greater than or equal to V_0 , set the acceleration to $-M$. The velocity is allowed to be between $V_0 - \tau$ and $V_0 + \tau$.

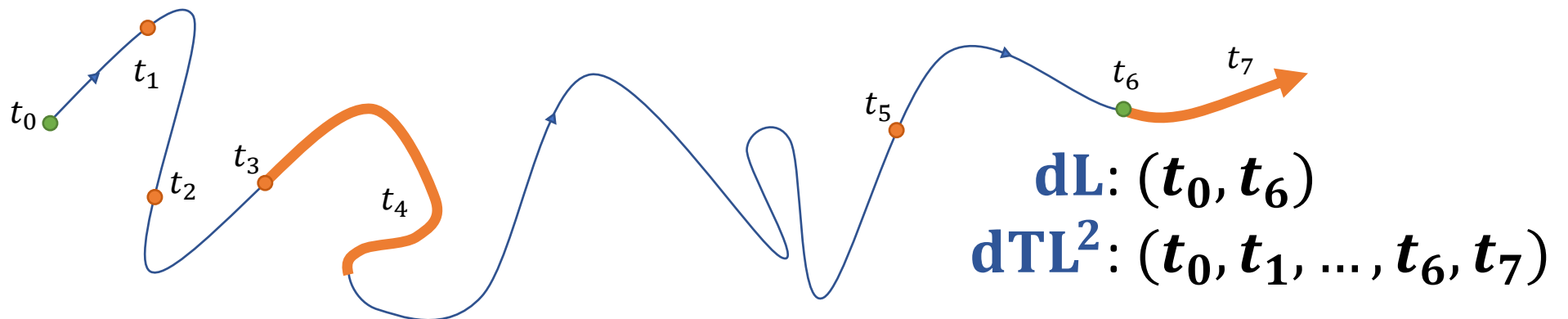
Semantics of **dL** vs. **dTL²**

In **dL** :

- Points in space defined as Environment types
 $\mathcal{E} : [\mathbb{N} \rightarrow \mathbb{R}]$
- Semantically relates input/output states of an HP.

In **dTL²** :

- Traces defined as lists with elements:
 1. Environment type
 2. $f : [D \rightarrow \mathcal{E}]$ where D is finite or infinite
 3. Error
- Elements of the trace are semantically related in chronological order.



dL and the dTL² extension

dL allows **formal reasoning** of the input/output (io) of hybrid programs:

- For hybrid program Hp and state predicate P
 - Allruns, Someruns for io
 $[Hp]_{st}P, \quad \langle Hp \rangle_{st}P$

dTL² extends **dL** by also allowing for **formal reasoning** of traces:

- For a trace predicate ϕ
 - Allruns, Someruns for traces
 $[Hp]_{tr}\phi, \quad \langle Hp \rangle_{tr}\phi$
 - Globally, Eventually
 $\Box\phi, \quad \Diamond\phi$

$t \models \Box\phi :=$ For all suffixes s of t , $s \models \phi$

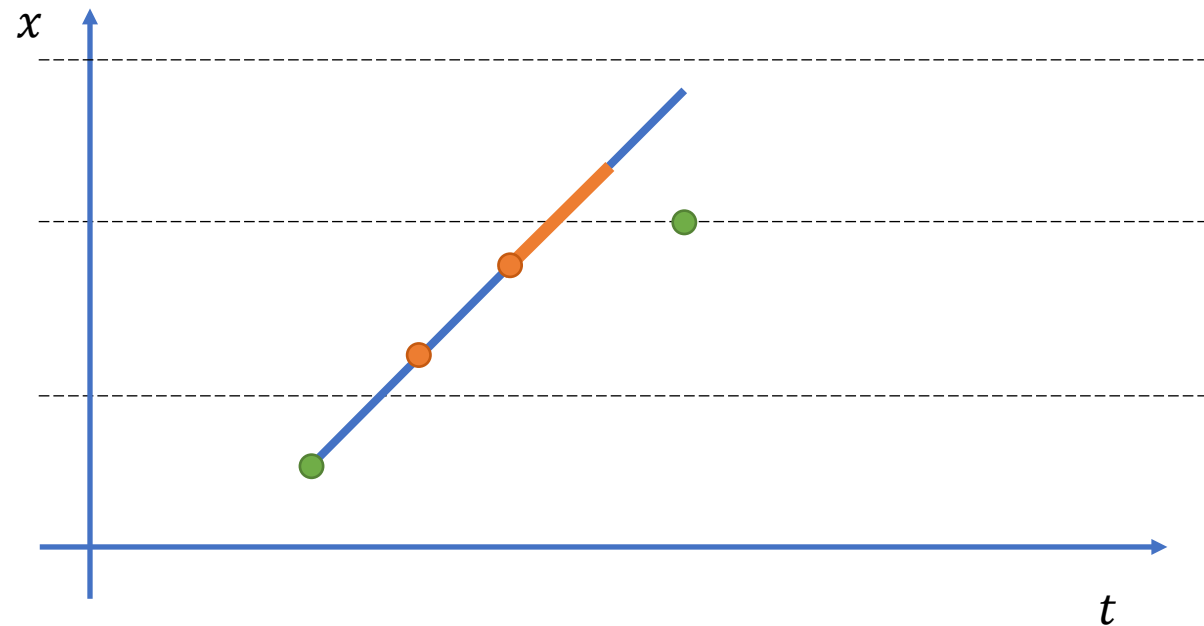
$t \models \Diamond\phi :=$ Exists a suffix s of t , $s \models \phi$

Solving Problems in **dL** and **dTL²**

Example:

$$\alpha := \{x' = 6\}; (x := \lfloor x \rfloor)$$

TRUE: $[\alpha]_{st}(x \in \mathbb{Z})$ and $\langle \alpha \rangle_{tr} \diamond (x \notin \mathbb{Z})$



Elimination of Error Traces

Example:

$$\alpha := x := 0; (? x < 10; x := x + 1)^*; ? x \geq 10$$

Semantics with errors

$$\llbracket \alpha \rrbracket = \{ (x = 0) \cdot \text{ERROR}, \dots, \\ (x = 0) \cdot (x = 1) \cdots (x = 9) \cdot \text{ERROR}, \\ (x = 0) \cdot (x = 1) \cdots (x = 9) \cdot (x = 10) \}$$

Semantics without errors

$$\llbracket \alpha \rrbracket = \{ (x = 0) \cdot (x = 1) \cdots (x = 9) \cdot (x = 10) \}$$

After the elimination of error traces...

$$[\alpha]_{tr} \diamond (x = 10)$$

dTL²: Differential Temporal Dynamic Logic – Rule Schema

Assignment:

$$\frac{\phi \vee [x := \theta]_{st}(\psi \vee \phi)}{[x := \theta]_{tr}\psi \sqcup \diamond\phi} \quad ([:=]\sqcup)$$

Differential Equation:

$$\frac{\neg\chi \vee (\phi \wedge [\bar{x} = \bar{\theta} \& \chi]_{st}(\psi \wedge \phi))}{[\bar{x} = \bar{\theta} \& \chi]_{tr}\psi \sqcap \square\phi} \quad ([']\sqcap)$$

Repetition:

$$\frac{[\alpha]_{tr}([\beta]_{tr}(\phi \sqcup \diamond\psi)) \sqcup \diamond\psi}{[\alpha; \beta]_{tr}(\phi \sqcup \diamond\psi)} \quad ([:;]\sqcup)$$

Test:

$$\frac{\neg\chi \vee \phi}{[?\chi]_{tr}\phi \blacktriangleleft \diamond\square\psi} \quad ([?]\blacktriangleleft \diamond)$$

..and more!

Hybrid Programs in PVS

Elements in a Trace list

[typecheck-file](#) | [evaluate-in-pvsio](#) | [view-as-markdown](#)

TraceState : DATATYPE

```
BEGIN
  IMPORTING hp_def,
  | | | | | hp_expr

  STATE(state : Environment) : state?
  INF_DIFF(s0:Environment, inf_behavior:{ib:[(hp(0)) -> Environment] | ib(0)=s0}) : inf_diff?
  STATE_DIFF(D:{D:(dd?) | EXISTS(b:posreal): D = closed_interval(0,b)}
  | | | | | , s0:Environment
  | | | | | , behavior:{b:[(D) -> Environment]|b(0)=s0}) : state_diff?
  ERROR : error?

END TraceState
```

Well founded Traces

```
wf_trace?(trace:list[TraceState]) : bool =
  cons?(trace) AND (NOT error?(car(trace))) AND
  FORALL (i:below(length(trace)-1)) :
  | state?(nth(trace,i)) OR state_diff?(nth(trace,i))
```

Trace semantics of ASSIGN

```
LET envi = car(trace),
  | | envo = nth(trace,1),
  | | l = assigns(hp) IN
state?(envi) AND
state?(envo) AND
  ((FORALL (i:below(length(l)))) :
  | LET (vari,rexpr) = nth(l,i) IN
  | state(envo)(vari) = rexpr(state(envi))) AND
  FORALL (varj:(not_in_map(l)) :
  | state(envo)(varj) = state(envi)(varj))
```

Formal Verification of Soundness of dTL²

Sequential Composition:

```
% [;]n
% [A] ([B] (P n □Q) n □Q)
% -----
% [A;B] (P n □Q)
prove | discharge-tccs | status-proofchain | show-prooflite
dltl_SEQcap: LEMMA
  ALLRUNS_tr(A, normDLGLOBALLY(ALLRUNS_tr(B, normDLGLOBALLY(P,Q)), Q)) |-
  ALLRUNS_tr(SEQ(A,B), normDLGLOBALLY(P,Q))
```

Test:

```
% [?]n
% ¬P v (Q ∧ R)
% =====
% [?P] (Q n □R)
prove | discharge-tccs | status-proofchain | show-prooflite
dltl_TESTcap_eq: LEMMA
DLOR(DLNOT(P),DLAND(Q,R)) = ALLRUNS_tr(TEST(P),normDLGLOBALLY(Q,R))
```

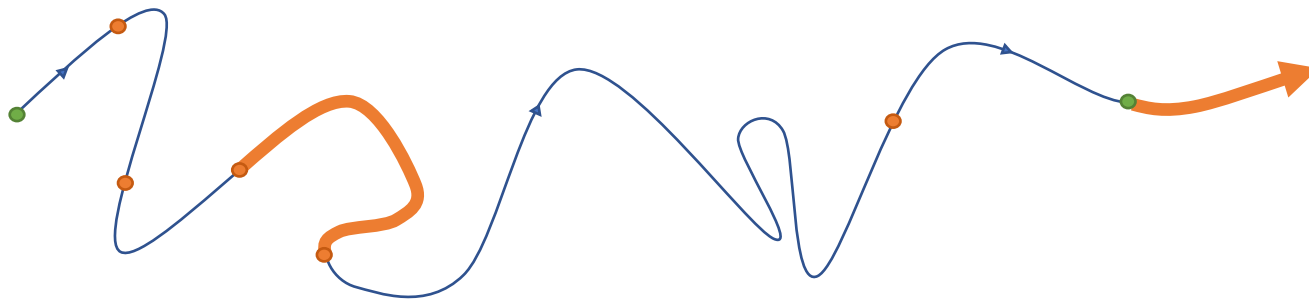
84 rules/axioms of dTL² in PVS

Summary

- **dTL²**: Differential Temporal Dynamic Logic for hybrid programs
- **PVS**: Interactive theorem prover

Result: Embedding of dTL² in PVS

- Formal verification of **dTL²**
- Fully operational embedding in **PVS**



Thank you for
listening!

