



# A NASA Perspective on Quantum Computing: Algorithmic opportunities and challenges

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**NASA QuAIL mandate:** *Determine the potential for quantum computation to enable more ambitious and safer NASA missions in the future*

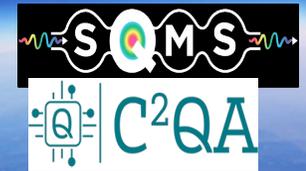
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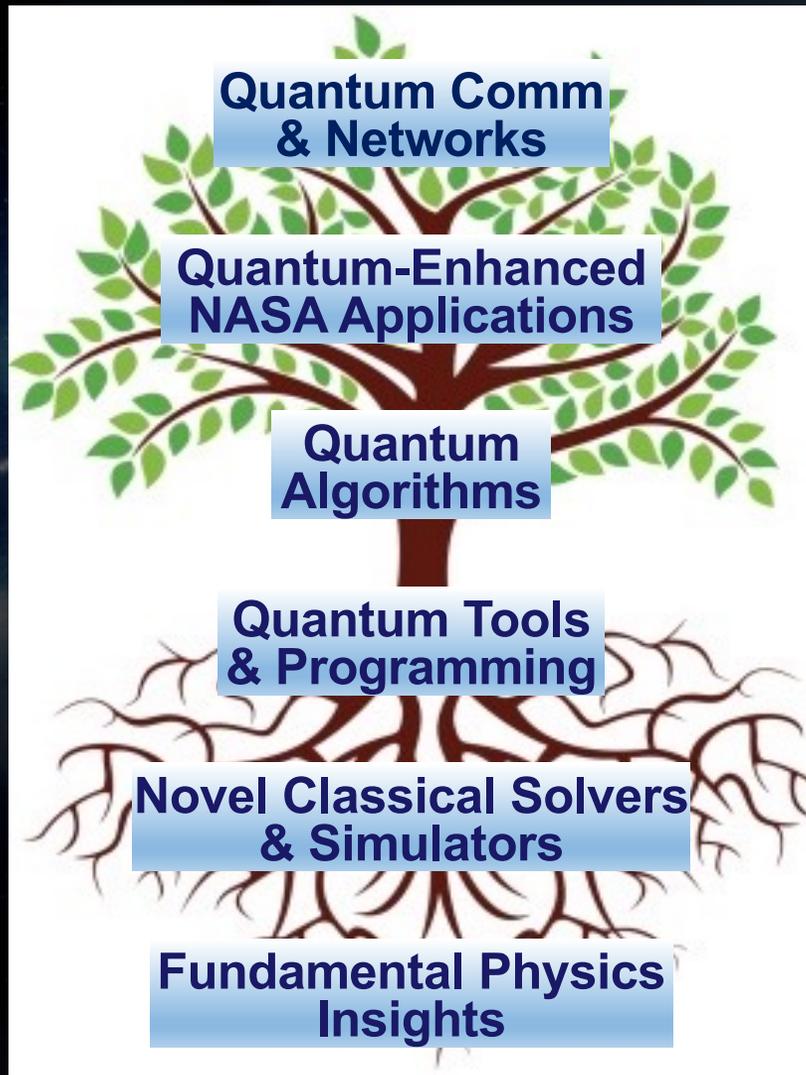
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+ active year-round intern program



# Quantum Computing R&D at NASA Ames



## Communication & Networks

Quantum networking

Distributed QC

## Application Focus Areas

Planning and scheduling

Material science

Fault diagnosis

Machine learning

Recently: Computational fluid dynamics (CFDs)

## Software Tools & Algorithms

Quantum algorithm design

Compiling to hardware

Mapping, parameter setting, error mitigation

Hybrid quantum-classical approaches

## Solvers & Simulators

Physics-inspired classical solvers

HPC quantum circuit simulators

## Physics Insights

Co-design quantum hardware

E. Rieffel *et al.* (2019), From Ansätze to Z-gates: A NASA view of quantum computing, *Adv. in Parallel Computing* **34**, 133–160

R. Biswas *et al.* (2017), A NASA perspective on quantum computing: Opportunities and challenges, *Parallel Computing* **64**, 81–98

# Exploration of Quantum Computing for NASA



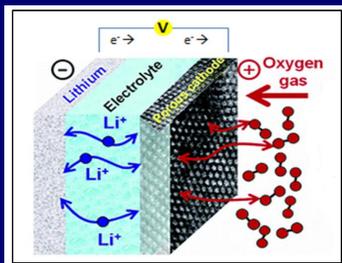
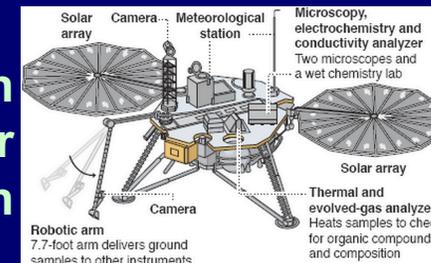
**Quantum ML  
for Earth  
Science Data  
Analysis**

**Superposition**



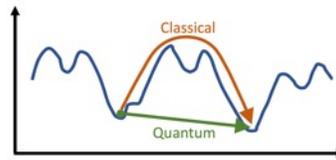
$$\psi_{cat} = \frac{1}{\sqrt{2}}(\psi_{alive} + \psi_{dead})$$

**Resource Allocation  
and Scheduling for  
Space Exploration**

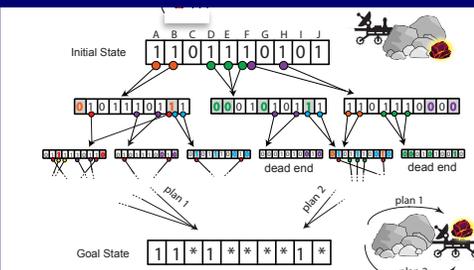


**Quantum  
Simulations for  
Aerospace Materials**

**Tunneling**



**Secure  
Airspace  
Communication**

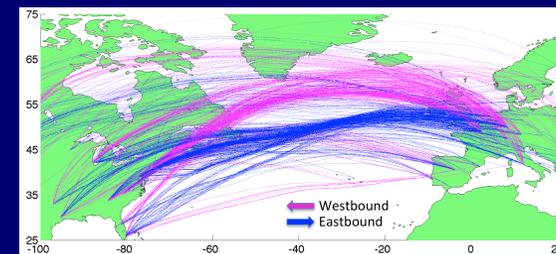


**Quantum  
Optimization  
for Mission  
Planning &  
Coordination**

**Entanglement**



**Air Traffic  
Management**



**Objective: Find "BETTER" solution  
faster time-to-solution OR more precise solution OR using less energy OR not found by classical  
methods**



# What is quantum computing (in one slide)?

The power of quantum computation comes from encoding information in a non-classical way

Quantum computers take advantage of purely quantum effects that are not available classically

quantum interference; quantum tunneling; quantum entanglement; quantum measurement, quantum many-body delocalization, quantum sampling

These effects can provide more efficient computation and higher levels of security than is available classically

What Shor's factoring algorithm can compute in days, would take a supercomputer longer than the age of the universe

Emerging quantum hardware enables empirical investigation of quantum optimization for myriad applications

# New Era for Quantum Computing

## Quantum supremacy achieved

- Perform computations not possible on even largest supercomputers in reasonable time
- Google – NASA – ORNL collaboration



*F. Arute et al. (2019), Quantum supremacy using a programmable superconducting processor, Nature 574, 505-510*

## 2023 Update

*A. Morvan, B. Villalonga, X. Mi, S. Mandrà, et al., (2023) Phase transition in Random Circuit Sampling, arXiv:2304.11119*

| Exp.             | 1 amp.               | 1 million noisy samples |                      |           |
|------------------|----------------------|-------------------------|----------------------|-----------|
|                  | FLOPs                | FLOPs                   | XEB fid.             | Time      |
| SYC-53 [9]       | $6.44 \cdot 10^{17}$ | $2.60 \cdot 10^{17}$    | $2.24 \cdot 10^{-3}$ | 6.18 s    |
| ZCZ-56 [10]      | $6.24 \cdot 10^{19}$ | $6.40 \cdot 10^{19}$    | $6.62 \cdot 10^{-4}$ | 25.3 min  |
| ZCZ-60 [11]      | $1.32 \cdot 10^{21}$ | $1.41 \cdot 10^{23}$    | $3.66 \cdot 10^{-4}$ | 38.7 days |
| <b>This work</b> | $4.74 \cdot 10^{23}$ | $6.27 \cdot 10^{25}$    | $1.68 \cdot 10^{-3}$ | 47.2 yr   |

## ... but so far only for a toy problem

- Quantum hardware currently too small for solving practical problems intractable on classical supercomputers
- These devices need to scale up and become more reliable

## So what to do in the interim?

- Unprecedented opportunity to invent, explore, and evaluate quantum algorithms empirically

## NASA QuAIL Focus

- **Algorithms and applications** to enable safer, more ambitious, and greater time- and energy-efficient missions
- **Tools** for advancing quantum computing, from quantum circuit simulation, noise characterization, error correction, compilation to realistic hardware

# Status of Quantum Algorithms

Quantum computing can do everything a classical computer can do *and*

Provable quantum advantage known for a few dozen quantum algorithms

Unknown quantum advantage for everything else

Status of classical algorithms

- Provable bounds hard to obtain
  - Analysis is just too difficult
- Best classical algorithm not known for most problems
- Empirical evaluation required
- Ongoing development of classical heuristic approaches
  - Analyzed empirically: ran and see what happens
  - E.g. SAT, planning, machine learning, etc. competitions
- **NISQ era supports unprecedented means for empirical analysis of quantum algorithms**
  - Quantum heuristics come into their own

**A handful of proven limitations on quantum computing**

**Conjecture: Quantum Heuristics will significantly broaden applications of quantum computing**

# Certainty and Randomness in Quantum Computation



Any computation a classical computer can do, a quantum computer can do with roughly the same efficiency

With the same probability of the outcome

If the classical computation is non-probabilistic, so is the quantum one

$O(\log n)$  overhead: solely due to making computation reversible

Like classical algorithms, some quantum algorithms are inherently probabilistic and others are not

First quantum algorithm was not probabilistic

- E.g. Deutsch-Jozsa algorithm solves problem with certainty that classical algorithms, of equivalent efficiency, could solve only with high probability

Shor's algorithms are probabilistic

Grover's is not intrinsically probabilistic

initial search algorithm was probabilistic, but

- slight variants, which preserve the speed up, are non-probabilistic



# Some closely related Quantum Optimization Algs: AQO, QA, QAOA

Phase separation operator based on the cost function

Usually  $H_P = -\sum C(z)|z\rangle\langle z|$  + (optionally) other terms, e.g. "penalty terms" to enforce constraints

Simple Driver/Mixing operator

Most frequently  $H_M = \sum_j X_j$ , though we will shortly see other mixers

Ground state easy to obtain

## AQO (special case of AQC)

- Evolution under  $H(t) = a(t)H_P + b(t)H_M$
- Slowly enough to stay in the ground subspace

## QA

- Evolution under  $H(t) = a(t)H_P + b(t)H_M$
- Many quick runs, thermal effect contribute

## QAOA

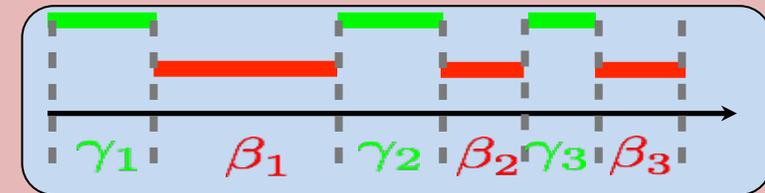
- Alternate application of  $H_P$  and  $H_M$
- For  $p$  alterations, the parameters are  $2p$  times/angles  $\gamma_1, \beta_1, \dots, \gamma_p, \beta_p$

# Quantum Approximate Optimization Algorithm



Gate model algorithm due to Farhi et al.

- Alternates between two Hamiltonians,  $p$  times
  - Phase separation (cost function dep.)
  - Mixing
  - $2p$  parameters: amount of time each Hamiltonian is applied
    - Parameter search means often a hybrid quantum-classical algorithm
    - Relation with Variational Quantum Eigensolvers (VQE)
- Aim: Provable approximation ratio



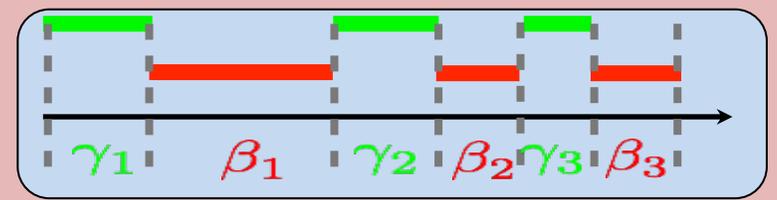
## Early results by Farhi and co-authors

- $p \rightarrow \infty$ : from AQO
  - Converges to optimum for  $p \rightarrow \infty$
- $p = 1$ : proofs modified from proofs for IQP circuits
  - Provably hard to sample output efficiently classically (up to standard complexity theory conjectures)
  - Briefly beat existing classical approx. ratio on MaxE3Lin2, but inspired better classical algorithm

# Quantum Alternating Operator Ansatz

Heuristic based on structure Quantum Approximate Optimization Algorithm of Farhi et al.

- Alternates between two Hamiltonians,  $p$  times
  - Phase separation (cost function dep.)
  - Mixing
  - $2p$  parameters: amount of time each Hamiltonian is applied
  - Aim: ~~Provable approximation ratio~~
  - Aim: Good typical performance
  - Better support for enforcing constraints, informed by compilation to hardware



## Early results

- Alternative algorithm for Grover's unstructured search problem
  - achieves  $\sqrt{N}$  query complexity by different means

# Quantum Alternating Operator Ansatz

Generalization of Quantum Approximation Optimization algorithm

$$Q_p(\beta, \gamma) = U_M(\beta_p)U_P(\gamma_p) \cdots U_M(\beta_1)U_P(\gamma_1)$$

$$e^{-i\gamma H_P}$$

$$\prod_j e^{-i\beta_k H_j}$$

$$|\beta, \gamma\rangle = Q_p(\beta, \gamma) |s\rangle$$

## Phase separator:

unitary for which

- The energy spectrum of  $H_P$  encodes the problem's objective function
- $H_P = -\sum C(z) |z\rangle\langle z|$

## Mixer: unitary which:

- Preserves the feasible subspace
- Provides nonzero transitions between all feasible states
- Not necessarily time evolution of a single local Hamiltonian
- $\beta_k$  depends on the level  $1 \leq k \leq p$ , but independent of  $H_j$

## Initial state $|s\rangle$ which:

- is a superposition of one or more solutions in the feasible subspace
- can be prepared efficiently

# Example: QAOA for Max-k-Colorable Subgraph

Problem: Given a graph  $G = (V, E)$ , and  $k$  colors  $1, \dots, k$ , find a color assignment maximizing the # of properly colored edges

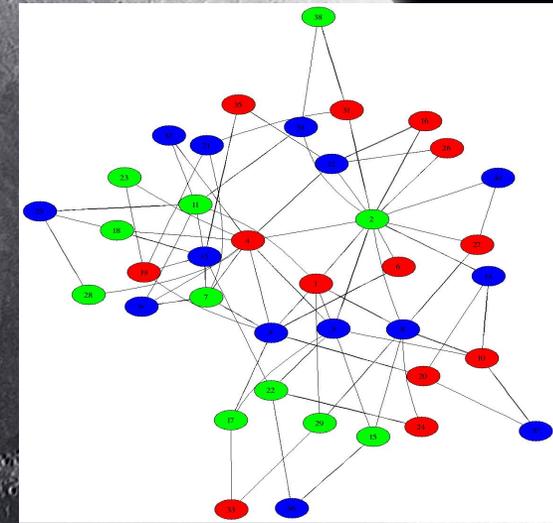
- Properly colored edge means endpoint vertices have been assigned different colors

“One-hot” Encoding:  $nk$  variables

$$x_{uj} = 1 \text{ iff vertex } u \text{ is colored color } j$$

Optimization: Write cost function as

$$C(x) = m - \sum_{(uv) \in E} \sum_{j=1}^k x_{uj} x_{vj}$$



**Must avoid invalid colorings**

e.g. if a vertex is labeled as both red and blue, or not colored at all

**Requires  $n$  constraints: one for each vertex  $u$**

$$\sum_{j=1}^k x_{uj} = 1$$

# Example: QAOA for Max-k-Colorable Subgraph



Could add constraints to the cost function to enforce penalties

- standard approach in quantum annealing

**Better: design mixer to keep evolution in feasible subspace (constant Hamming weight in colors for each vertex)**

**Feasible subspace is exponentially smaller search space than entire Hilbert space**

While still exponentially large

**Initial state choice**

Any classical feasible state

e.g. all colored red

Any superposition of feasible states

e.g. superposition of all colors (W state)

**Use a swap or XY-mixer (XX+YY) on the colors rather than bit flip mixer:**

instead of  $\sum_j X_j$ , use sum of swap operators,

$$|00\rangle\langle 00| + |10\rangle\langle 01| + |01\rangle\langle 10| + |11\rangle\langle 11|$$

between colors at a vertex

**Ring mixer:**

order the colors, and apply swaps to adjacent colors only

**Complete mixer:**

swap for every pair of colors

**Complete mixer mixes more quickly, but has higher circuit depth, especially when compiled to realistic hardware**

# Example: QAOA for Max- $\kappa$ -Colorable Subgraph



Partitioned Mixers: Products of  $U_{Bv} = e^{-iBv}$ . Don't commute, so different orders give different mixers

$$U_{\text{parity}}(\beta) = U_{\text{last}}(\beta)U_{\text{even}}(\beta)U_{\text{odd}}(\beta),$$

where

$$U_{\text{odd}}(\beta) = \prod_{a \text{ odd}, a \neq n} e^{-i\beta(X_a X_{a+1} + Y_a Y_{a+1})},$$

$$U_{\text{even}}(\beta) = \prod_{a \text{ even}} e^{-i\beta(X_a X_{a+1} + Y_a Y_{a+1})},$$

$$U_{\text{last}}(\beta) = \begin{cases} e^{-i\beta(X_d X_1 + Y_d Y_1)}, & \kappa \text{ odd,} \\ I, & \kappa \text{ even.} \end{cases}$$



# Many variants of QAOA

Relation between parameter setting in QAOA and annealing schedule choice in quantum annealing

Close ties to sampling, e.g. for ML

## Developing General Theory of Iterative Quantum Algorithms

### Components of an Iterative Quantum Algorithm

*Preparation Rule* – Run Quantum (or Classical) Algorithm to get state

*Selection Rule* – Rank features in the system based on the prepared state

*Reduction Rule* – Eliminate a feature of the system based on ranking

Iterative Quantum Algorithms can be designed to guarantee enforcement of constraints

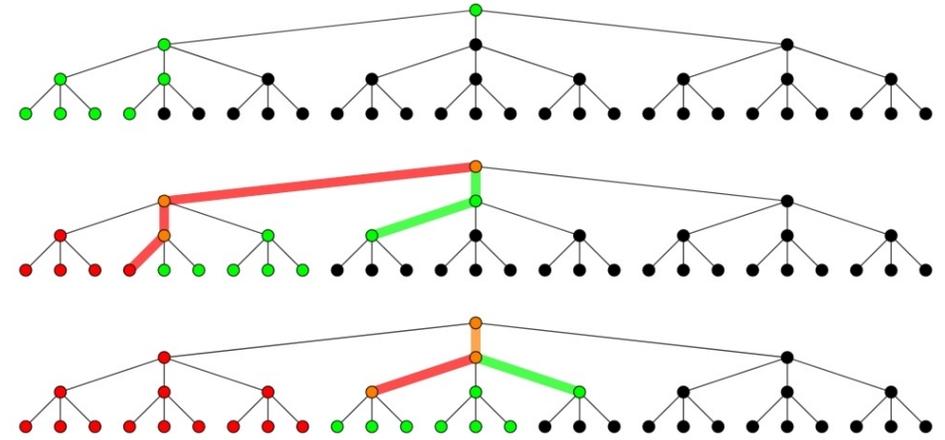
Classical updates resemble classical greedy algorithms and can pull classical proof techniques

- Special case: IQA for Max Independent Set

L. T. Brady & S. Hadfield, "Iterative Quantum Algorithms for Maximum Independent Set: A Tale of Low-Depth Quantum Algorithms" arXiv:2309.13110  
S. Hadfield, L. T. Brady, et al., "Quantum-Enhanced Classical Algorithms" (in preparation)

# Brief Glimpse: Quantum-Accelerated Constraint Programming

- In constraint programming (CP), problems are solved with backtracking tree search augmented by logical inference
- Quantum algorithms can accelerate the inference process being performed at each node in the tree
- These quantum inference algorithms can then be integrated within classical, fully-quantum, or partially-quantum backtracking tree search schemes
- Partially quantum backtracking schemes yield speedups for smaller sections of the tree, intended for early, more resource-constrained quantum devices



**Other good target state-of-the-art classical algorithms for quantum acceleration?**

Booth, Kyle EC, Bryan O’Gorman, Jeffrey Marshall, Stuart Hadfield, and Eleanor Rieffel. *Quantum-accelerated global constraint filtering*. In *International Conference on Principles and Practice of Constraint Programming*, pp. 72-89. 2020  
Booth, Kyle EC, Bryan O’Gorman, Jeffrey Marshall, Stuart Hadfield, and Eleanor Rieffel. *Quantum-accelerated constraint programming*. *Quantum* 5 (2021): 550.

# Quantum Distributed Algorithms for Approximate Steiner Trees and Directed Minimum Spanning Trees

Joint work with Phillip Kerger, David Bernal Neira, Zoe Gonzales Izquierdo

- We provide quantum distributed algorithms to tackle challenging graph problems
  - Approximate Steiner Tree Problem
  - Directed Minimum Spanning Tree (Arborescence)
- These quantum algorithms provide an asymptotic improvement with respect to the current best known classical algorithm in terms of computational rounds in the CONGEST CLIQUE model
- We provided detailed analysis for the main algorithmic step: finding the all-pairs shortest paths
- We obtained complexity results realizing impractical scales where quantum counterparts become better than classical

Phillip A. Kerger, David E. Bernal Neira, Zoe Gonzalez Izquierdo, Eleanor G. Rieffel, "Mind the  $\tilde{O}$ : asymptotically better, but still impractical quantum distributed algorithms," *Algorithms* 16(7), 332, 2023. [arXiv:2304.02825](https://arxiv.org/abs/2304.02825)





# Main results



Quantum distributed algorithms to tackle challenging graph problems

- Approximate Steiner Tree Problem
- Directed Minimum Spanning Tree Problem (Arborescence Problem)

Asymptotic improvement over current best known classical algorithm in terms of computational rounds in CONGEST CLIQUE model

Detailed analysis for the main algorithmic steps

Non-asymptotic complexity results mean both prior classical distributed algorithms and our quantum algorithm only have advantage over simpler schemes at impractically large graph sizes

*P Kerger, DE Bernal Neira, Z Gonzalez Izquierdo, EG Rieffel, **Mind the  $\tilde{O}$ : Asymptotically better, but still impractical, quantum distributed algorithms**, Algorithms 16 (7), 332, 2023*

New classical distributed algorithm for the broad class of Survivable Network Design Problems (SNDPs) in CONGEST CLIQUE model

New quantum distributed algorithm for SNDPs in QUANTUM CONGEST CLIQUE model

Main ingredients:

- Building on prior distributed all-pair shortest path (APSP) algorithm
- Added routing table computation
- Detailed analysis of constant and log factors

*P Kerger, DE Bernal Neira, Z Gonzalez Izquierdo, EG Rieffel, **Classical and Quantum Distributed Algorithms for the Survivable Network Design Problem**, arXiv:2404.10748*

# Background



INTELLIGENT  
SYSTEMS  
DIVISION

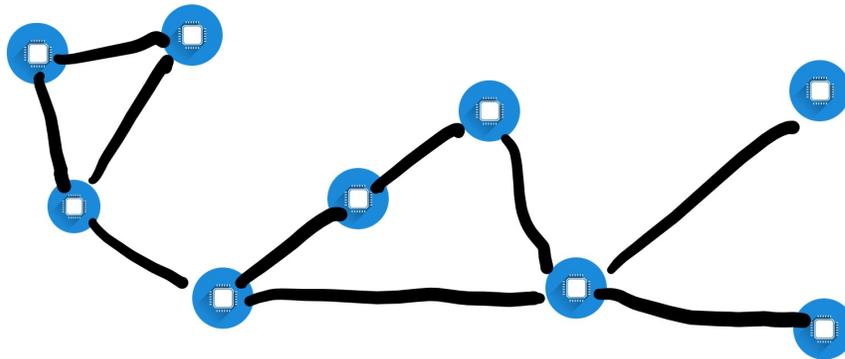


# Introduction: Algorithms on Distributed Data



Algorithms on Distributed Data: Network of Multiple Processors that communicate

- Model as graph where nodes are processors
- Each node its own starting information
  - For graph problems, this is often a list of neighbor nodes



- Example: Network of spacecraft, satellites, or control stations that each have computing power and can communicate
- **Goal:** Answer some question about that distributed information, through communication and computation among the processors



# Introduction: Two Classical Models



Graph  $G = (V, E, W)$  with  $n = |V|$  number of nodes,  $m = |E|$  number of edges, and  $W$  the weights on the edges

**CONGEST Model:** Aim is to minimize the number of rounds

Computation happens in rounds (compute, communicate, compute, communicate, ...)

Congested: Communication limited by message size: each node can send to each of its neighbors

$O(\log(n))$  bits each round

—  $\log(n)$  is length of a node id

Unlimited local computation at each node

Nodes can communicate only with their neighbors

**CONGEST-CLIQUE Model:**

1., 2., 3. are the same as for Congest Model

4. All nodes can communicate with each other

Key difference: communication graph distinct from graph  $G$

Initial conditions: Each node knows

- its own ID
- the ID's of its neighbors

assuming ID's are 1 to  $n \rightarrow \log(n)$  bits to encode

Aim: Answer a question about graph in as few rounds as possible

- Ex: Spanning trees, subgraph detection, shortest paths...



## Core Research Question



What problems can benefit from a distributed *quantum* approach?



# Introduction: Models



Graph  $G = (V, E, W)$  with  $n = |V|$  number of nodes,  $m = |E|$  number of edges, and  $W$  the weights on the edges

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Ex: Spanning trees, subgraph detection, shortest

**Quantum versions:** Take CONGEST or CONGEST-CLIQUE, but allow messages to consist of  $O(\log(n))$  qubits



# CONGEST: Negative Results



Main reference: Elkin et al 2012,

## “Can Quantum Communication Speed Up Distributed Computation?”

- Proved limitations for quantum CONGEST model
- Quantum communication does NOT provide an improvement for many fundamental problems: **Shortest paths**, Minimum Spanning Tree, **Steiner Tree**, Min Cut, Hamiltonian Cycle...
- **Intuition:** In CONGEST, a significant bottleneck can be communicating between “distant” parts of the network – qubits don’t help with that!

• Elkin, M., Klauck, H., Nanongkai, D., & Pandurangan, G. (2014, July). Can quantum communication speed up distributed computation?. In *Proceedings of the 2014 ACM symposium on Principles of distributed computing* (pp. 166-175).



# CONGEST CLIQUE



- Elkin et al.'s results and analysis **do not** carry over to the CONGEST CLIQUE
- So, can quantum communication help in this model?

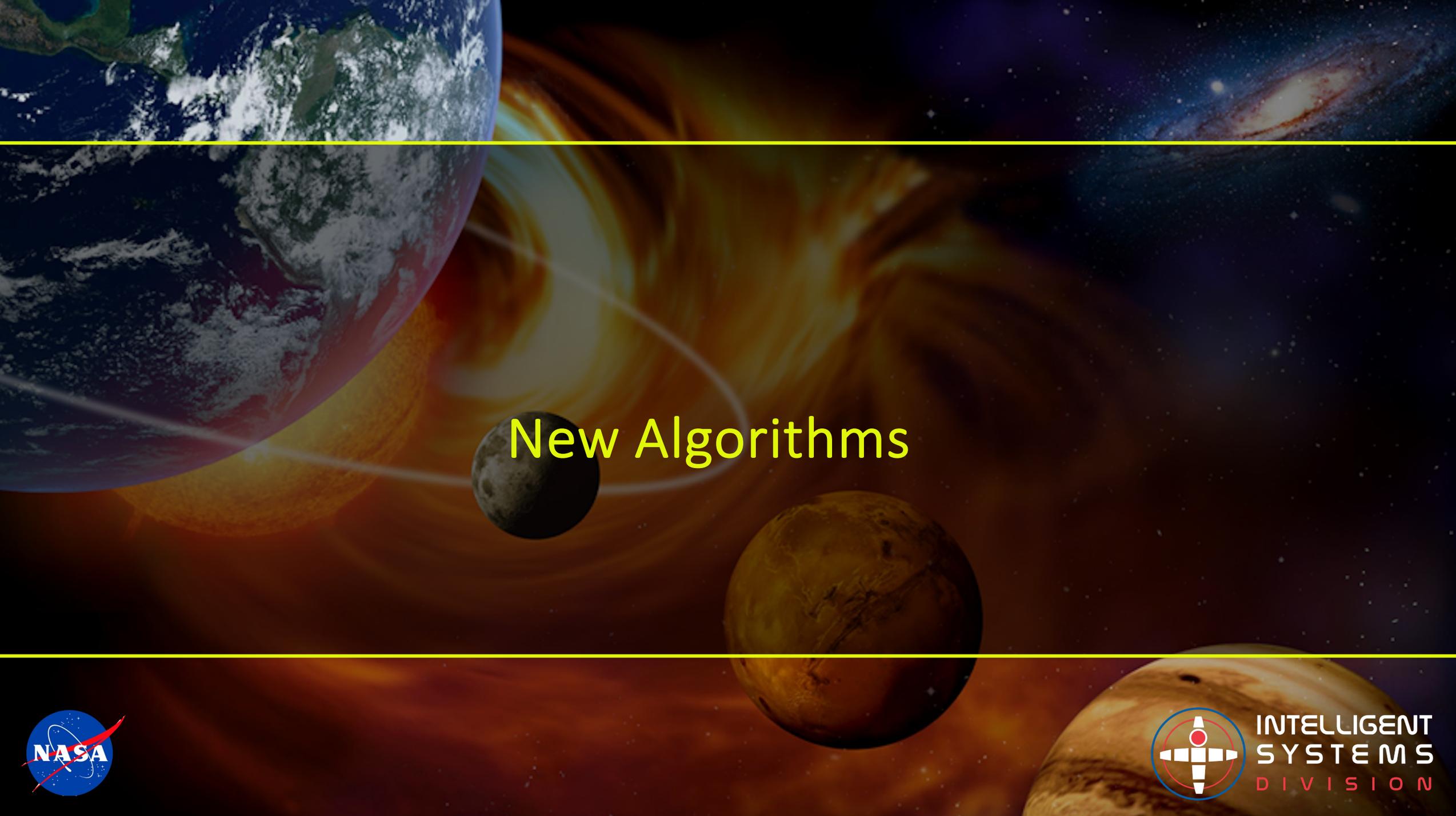
Surprising **positive results** in in Quantum CONGEST CLIQUE Model (qCCM):

- Faster Triangle Detection, Izumi & Le Gall 2019
- Faster All-Pairs **Shortest-Paths** (APSP), Izumi & Le Gall 2020
  - $\tilde{O}(n^{1/4})$  in quantum versus  $\tilde{O}(n^{1/3})$  in classical
  - This was in Elkin's list of problems not admitting speedups!

$$\begin{aligned} f(n) \in \tilde{O}(g(n)) \\ \text{if} \\ \exists k: f(n) \in \mathcal{O}(g(n) \log^k n) \end{aligned}$$

**For which other problems can we exhibit improvements in the quantum CONGEST CLIQUE model?**

• Elkin, M., Klauck, H., Nanongkai, D., & Pandurangan, G. (2014, July). Can quantum communication speed up distributed computation?. In *Proceedings of the 2014 ACM symposium on Principles of distributed computing* (pp. 166-175).  
• Izumi, T., & Le Gall, F. (2017, July). Triangle finding and listing in CONGEST networks. In *Proceedings of the ACM Symposium on Principles of Distributed Computing* (pp. 381-389).  
• Izumi, T., & Le Gall, F. (2019, July). Quantum distributed algorithm for the All-Pairs Shortest Path problem in the CONGEST-CLIQUE model. In *Proceedings of the 2019 ACM Symposium on Principles of Distributed Computing* (pp. 84-93).

A composite image of space featuring Earth in the top left, the Moon in the center, Mars in the bottom center, Jupiter in the bottom right, and a galaxy in the top right. A yellow horizontal line runs across the middle of the image.

# New Algorithms



INTELLIGENT  
SYSTEMS  
DIVISION



# Algorithmic Recipe



Approach: Make use of previous techniques such as

1. Distributed Grover Search
2. Triangle Finding
3. Distance Products
4. Shortest Paths and Routing Tables

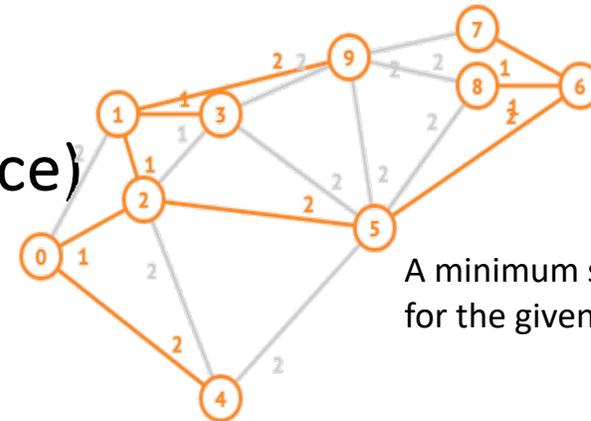
New algorithms in Quantum CONGEST-CLIQUE Model (qCCM) that succeed with high probability for

- (approximately optimal) Steiner Trees
- Directed Minimum Spanning Trees (Arborescence)

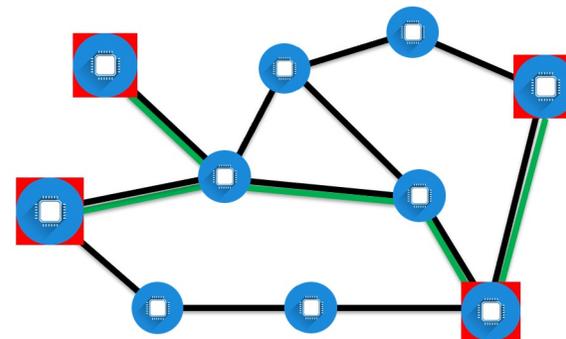
in asymptotically fewer rounds required than for any known classical algorithm

$$\rightarrow \tilde{O}(n^{1/4}) \text{ versus } \tilde{O}(n^{1/3})$$

Exact complexity analysis of quantum and classical algorithms reveals improvements needed for both to become practical!



A minimum spanning tree (orange) for the given graph (grey)



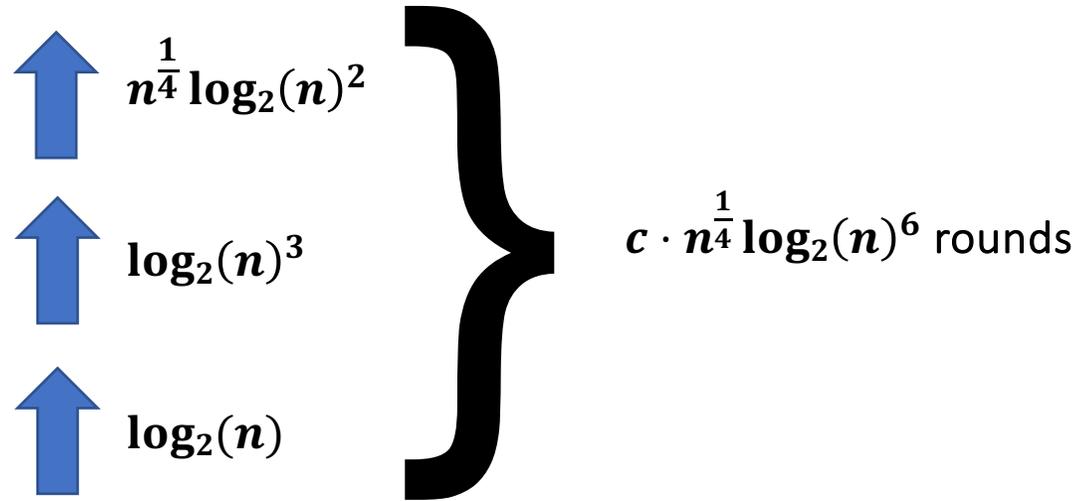
Steiner tree (green) for graph with marked terminal nodes (red)



# Algorithmic Recipe & Complexity



1. Distributed Grover Search helps with...
2. Triangle Finding helps with...
3. Distance Products helps with...
4. Shortest Paths and Routing Tables helps with...
5. Steiner and Directed Minimum Spanning Trees!



In CONGEST CLIQUE, can solve anything in  $n$  rounds:

To be practical, need roughly

$$3200 \cdot n^{\frac{1}{4}} \log_2(n)^6 < n$$

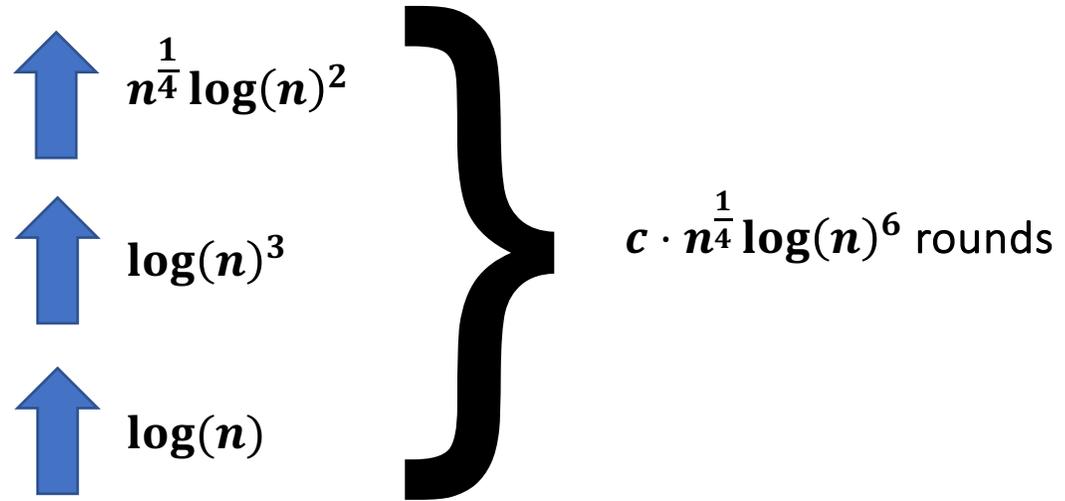
The asymptotic results are exciting!  
 But more work is needed to bring these algorithms into a practical realm



# Algorithmic Recipe & Complexity



1. Distributed Grover Search helps with...
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5. Steiner and Directed Minimum Spanning Trees!



In CONGEST CLIQUE, can solve anything in  $n$  rounds:

To be practical, need roughly

$$3200 \cdot n^{\frac{1}{4}} \log(n)^6 < n$$

for which

$$n > 10^{18}$$

is required!!  $10^{11}$  for classical  $\tilde{O}\left(n^{\frac{1}{3}}\right)$  counterpart.

**The asymptotic results are exciting!  
But more work is needed to bring these algorithms into a practical realm**



# Distributed quantum algorithms: Future directions



What further improvements be made to bring the asymptotic speedup closer to practical?

Other problems that for which these methods can demonstrate advantages in the quantum setting?

There are other distributed computing models and approaches. What can be shown in quantum versions of these modes?

Pay attention to quantities hidden in  $\tilde{O}(N)$  notation!

- Constants and log factors can be important in both the near and long term

Distributed quantum computing is in its infancy, with few results and many open directions

Many opportunities for classical computing to inform quantum computing and to work with or as part of quantum computing



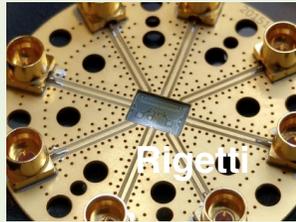
# Other Topics

# Status of Quantum Hardware



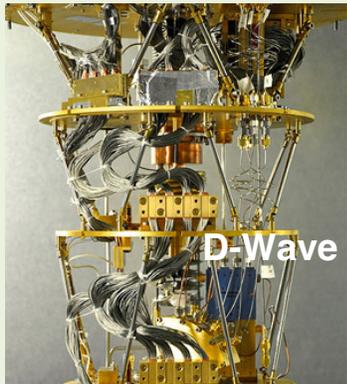
## General Purpose:

Universal quantum processors

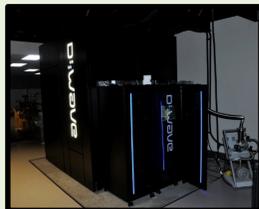


## Special Purpose:

E.g. Quantum annealers



**Noisy  
Intermediate-  
Scale  
Quantum  
(NISQ)  
devices**



Many different underlying physical substrates for quantum processors:

- Superconducting
- Trapped ion
- Photonic
- Other
  - - Electron spins in silicon
  - - Neutral atom, cold atom
  - - Topological, anyon based quantum computing

All quantum hardware is small and non-robust

Special purpose vs general purpose processors

- Algorithm/hardware codesign

Number of qubits alone is not a good measure

- Analogy: billions of switches do not a classical computer make

Other key factors

- precision, speed, and generality of the control
  - particularly operations involving multiple qubits
- how long quantum coherence can be maintained
- stability over time
- speed with which processors can be calibrated



# Future quantum computers

Application scale quantum computers will resemble supercomputers

Many quantum processing units (QPUs), and classical processing units

Quantum and classical communication

Robustness:

Quantum error correction and fault tolerance are mature areas, with continuous breakthroughs

Tens to thousands of logical qubits per QPUs

Rule of thumb: ~1000 physical qubits per logical qubits

Synergies and differences between quantum computer and supercomputer architectures

2D local structures, higher dimensional comm at larger scales

No cloning severely limits duplication

E.g. Two ways to move quantum info

- local quantum physical links
- teleportation across arbitrary distances
  - requires prior set up of entanglement through local links, and classical comm as part of teleportation

# Quantum-inspired classical algorithms and hardware



Quantum Monte Carlo

Improved classical techniques for simulating quantum systems

De-quantized quantum algorithms

e.g. for E3Lin2

e.g. certain sampling and quantum ML algorithms

Quantum proofs for classical theorems (Survey: Drucker & Wolf arXiv:0910.3376)

DARPA's Quantum-Inspired Classical Computing (QuICC) program

# PySA: Suite of State-Of-Art classical optimization algs



## Features and state-of-the-art implementations:

- Modern C++17 with template metaprogramming for high level of abstraction
- Compile time optimization for improved performance

## Algorithms:

- Parallel Tempering
- Ergodic and non-ergodic Isoenergetic cluster moves
- Approximate solution using mean-field theory

## Recent augmentations:

- Improved Python interface

*We continuously update PySA with optimized code for state-of-the-art classical optimization, including physics inspired approaches we have developed*

**Open Source Code: <https://github.com/nasa/pysa>**



# Brief Glimpse: Qubit Routing for Quantum Circuits

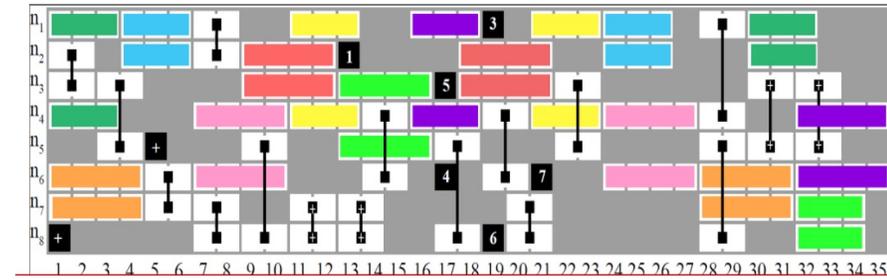


Quantum algorithms must be compiled before they can be run on quantum hardware

- Gate synthesis: rewrite only in terms of native gates
- Qubit routing: move information around to where two-qubit gates can act, given connectivity constraints

Can be viewed as a temporal planning problem

- Applied state of the art temporal planners
- Minimizes “makespan,” the time it takes to carry out the computation
- There is now a quantum circuit compilation domain in the International Planning Competitions temporal planning track
- Combining CP with temporal planning is advantageous



D Venturelli, M Do, E. Rieffel, J Frank, Compiling quantum circuits to realistic hardware architectures using temporal planners. Quantum Science and Technology 3 (2), 2018  
Venturelli, Davide, Minh Do, Bryan O’Gorman, Jeremy Frank, Eleanor Rieffel, Kyle EC Booth, Thanh Nguyen, Parvathi Narayan, and Sasha Nanda. "Quantum circuit compilation: An emerging application for automated reasoning." (2019).

Booth, Kyle EC, Minh Do, J. Christopher Beck, Eleanor Rieffel, Davide Venturelli, and Jeremy Frank. "Comparing and integrating constraint programming and temporal planning for quantum circuit compilation." In Twenty-Eighth international conference on automated planning and scheduling. 2018



# Quantum Error Correction

Quantum error correction initially thought impossible!

No cloning principle: an unknown quantum state cannot be copied reliably without destroying the original

Quantum information theory was just too interesting

Steane and Shor & Calderbank saw a way to finesse what had seemed insurmountable barriers to quantum error correction

Now quantum error correction is one of the most developed areas of quantum computing

Uses properties of quantum measurement and entanglement to its advantage

Stabilizer code formulation most common

Surface code

Remains active area of research

Subsystem codes

Dynamical/Floquet codes

Quantum Low-Density Parity Check (LDPC) codes

Decoders



# QuAIL simulation software and theory



## Quantum Circuit simulation software

Google – NASA – ORNL collaboration  
F. Arute *et al.* (2019),  
Quantum supremacy using a programmable superconducting processor,  
*Nature* **574**, 505-510



Recent NASA collaboration with Google AI

A. Morvan, B. Villalonga, X. Mi, S. Mandrà, et al., **Phase transition in Random Circuit Sampling**, arXiv:2304.11119, April 24, 2023

Experimental results that are significantly harder to simulate than the 2018 ones

## Open Quantum System Simulation

N Suri, J Barreto, S Hadfield, N Wiebe, F Wudarski, J Marshall, Two-Unitary Decomposition Algorithm and Open Quantum System Simulation, *Quantum* **7**, 1002 (2022)

- avoids classically expensive singular value decomposition (SVD)
- requires only a single call to state preparation oracle
- calls to the encoding oracle can also be reduced at the expense of an acceptable error in measurements

## Simulation of Photonic Quantum Systems

- Effect of **distinguishability** and **loss** errors in QIP
  - J Marshall, Distillation of Indistinguishable Photons *Phys. Rev. Lett.* **129** (21), 21360
- Efficient representations
  - J Marshall, N Anand Simulation of quantum optics by coherent state decomposition, arXiv:2305.17099



# HybridQ: A Hybrid Quantum Simulator for Large Scale Simulations



INTELLIGENT  
SYSTEMS  
DIVISION

Hardware agnostic quantum circuit simulator

Can run tensor contraction simulations, direct evolution simulation and Clifford+T simulations using the same syntax

Fully compatible with Python (3.8+)

Low-level optimization achieved by using C++ and Just-In-Time (JIT) compilation with JAX and Numba,

It can run seamlessly on CPU/GPU and TPU, either on single or multiple nodes (MPI) for large scale simulations, using the exact same syntax

User-friendly interface

Can run on supercomputers or laptop

Commutations rules are used to simplify circuits (useful for QAOA)

Expansion of density matrices as superpositions of Pauli strings accepts arbitrary non-Clifford gates,

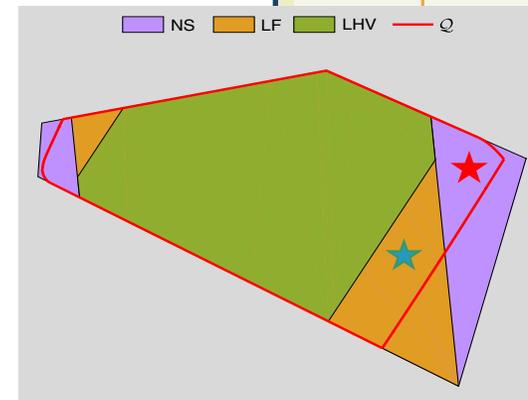
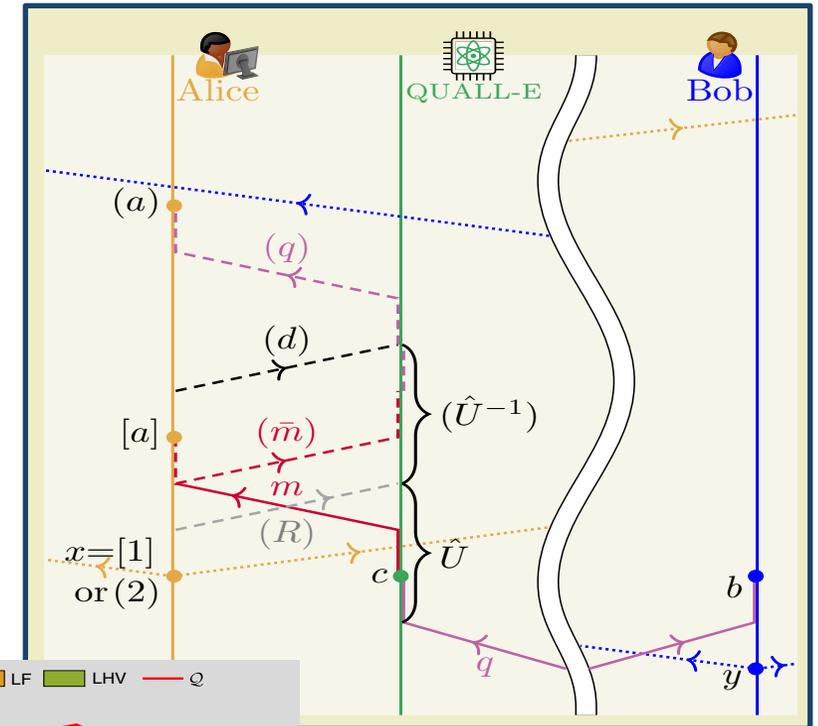
Open-source project with continuous-integration, multiple tests and easy installation using either `pip` or `conda`

Open source code available at <https://github.com/nasa/HybridQ>

*S. Mandrà, J. Marshall, E. G. Rieffel, R. Biswas, HybridQ: A Hybrid Simulator for Quantum Circuits, QCS 2021, arXiv:2111.06868*

# Wigner's friend inequalities & Experiments

- Wigner friend scenario recent work
  - new inequalities, with weaker assumptions than Bell's inequalities
  - Proof-of-principle experiments have been done
    - Single photon as friend
- Full experiment would combine Artificial Intelligence and Quantum Computing
  - QUALL-E
- Open research directions for experiments between proof-of-principle and full
  - Space-based experiments



**Polytope with 96 extreme points and 932 facets**

Bong, Kok-Wei, Aníbal Utreras-Alarcón, Farzad Ghafari, Yeong-Cherng Liang, Nora Tischler, Eric G. Cavalcanti, Geoff J. Pryde, and Howard M. Wiseman. "A strong no-go theorem on the Wigner's friend paradox." *Nature Physics* 16, 12 (2020)

H.M. Wiseman, E.G. Cavalcanti, E.G. Rieffel, A "thoughtful" Local Friendliness no-go theorem: a prospective experiment with new assumptions to suit, *arXiv:2209.08491 (Accepted to Quantum)*

# A Historical Perspective



Illiac IV – first massively parallel computer

- 64 64-bit FPUs and a single CPU
- 50 MFLOP peak, fastest computer at the time

Finding good problems and algorithms was challenging

Questions at the time:

- How broad will the applications be of massively parallel computing?
- Will computers ever be able to compete with wind tunnels?



NASA Ames director Hans Mark brought Illiac IV to NASA Ames in 1972



# For more info

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Eleanor G. Rieffel, Stuart Hadfield, Tad Hogg, Salvatore Mandrà, Jeffrey Marshall, Gianni Mossi, Bryan O'Gorman, Eugeniu Plamadeala, Norm M. Tubman, Davide Venturelli, Walter Vinci, Zihui Wang, Max Wilson, Filip Wudarski, Rupak Biswas, ***From Ansätze to Z-gates: a NASA View of Quantum Computing***, arXiv:1905.02860

Rupak Biswas, Zhang Jiang, Kostya Kechezhi, Sergey Knysh, Salvatore Mandrà, Bryan O'Gorman, Alejandro Perdomo-Ortiz, Andre Petukhov, John Realpe-Gómez, Eleanor Rieffel, Davide Venturelli, Fedir Vasko, Zihui Wang, ***A NASA Perspective on Quantum Computing: Opportunities and Challenges***, arXiv:1704.04836