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DIGITAL SYSTEMS DESIGN LANGUAGE

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First Annual Technical Report
October 1979
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for

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DESIGN SYNTHESSES OF DIGITAL SYSTEMS
George C. Marshall Space Flight Center
Alabama, 35812
FOREWORD

This is a technical summary of the research work conducted during October 1, 1978 to September 30, 1979 by The University of Alabama in Huntsville towards the fulfillment of the Contract NAS8-33096 from the George C. Marshall Space Flight Center, Alabama. The NASA technical officer for this contract is Mr. Robert E. Jones.

The author gratefully acknowledges the numerous discussions with and helpful comments of Mr. John M. Gould during this research work, and thanks Professor Donald Dietmeyer of the University of Wisconsin-Madison for providing the DDL Software.
Digital Systems Design Language (DDL) has been implemented on the SEL-32 Computer Systems of the Electronics and Controls Laboratory. This document provides the details of the language; the translator and the simulator programs. Several example descriptions and a tutorial on hardware description languages are provided, to guide the user.
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VI
1. INTRODUCTION

Hardware Description Languages (HDL) provide a convenient medium of inputting the design details into a design automation system. This report gives the details of one such language, Digital Systems Design Language (DDL), selected for integration into the Computer Aided Design and Test System (CADAT) of the Electronics and Controls Laboratory.

Chapter 2 provides the language details, Chapter 3 discusses the translator program and Chapter 4 discusses the Simulator Program. Some example descriptions are provided in Chapter 5. A tutorial on Hardware Description Languages is provided in the Appendix. An exhaustive bibliography for some of the literature in this area is also provided in the Appendix. Readers not familiar with any HDL are referred to the Appendix before reading the rest of the report.

The Simulator and Translator Programs are currently being tested on SEL-32 Computer System and hence, the complete deck set up details for the use of these programs is not included in this manual.
2. THE LANGUAGE [31]*

DDL was introduced in 1967 by Duley and Dietmeyer [33]. A translator and a simulator are written for a subset of this language in IFTRAN, an extended version of FORTRAN [35,36]. These programs are implemented in FORTRAN on SEL 32 Computer System. The translator (DDLTRN) translates a DDL description into a set of Boolean equations and register-transfer statements. The simulator (DDLSIM) enables the system designer to verify his design. The output of the translator is an input to the simulator. Simulation parameters are to be input by the designer. In DDL the structural elements are explicitly declared. At the lower level of description, functional and structural elements correspond directly to the actual elements of the system. DDL is highly suitable for describing the system at the gate, register transfer and major combinational block level.

The logical statements can be formed using the available primitive operators. The functional specification of the system consists of these logical statements, in blocks. The statements describe the state transitions of a finite state machine controlling the processes of the intended algorithm. The block then appears as an automaton.

Parallel operations are permitted. Synchronous behavior is described by either identifying the pulses or by including delay elements described in terms of multiples of clock pulses. Asynchronous behavior is modelled by using conditional statements. Data paths can be explicitly declared by using terminal declarations.

*The numbers in brackets point to the references listed in the Appendix.
DDL is a "block-oriented" language; the blocks of a DDL description usually correspond to natural divisions (blocks) of the hardware being described. Thus a computer may have a major block called an "ALU," which contains a block called "adder," which consists of interconnected logic blocks called "full-adders." This nested view of the hardware can be directly reflected in the DDL description of the computer.

Both facility declarations and operations can appear within the body of the more complex declarations that have a heading part. Identifiers declared within such complex declarations are said to be local facilities of that declaration, and are global facilities of complex declarations that appear in the body of the encompassing declaration. Other complex declarations that parallel the encompassing declaration cannot control or sense such facilities. Operations can reference only facilities that are local or global to the block in which they appear. Thus the same identifier may be declared in more than one parallel block without ambiguity.

Figure 2-1 illustrates some of the possibilities. Facilities A, B, and C are declared facilities of the overall block named SYSTEM. These facilities are global to all blocks within SYSTEM; any or all of these blocks may control or sense the states of facilities A, B, and C. Hence A, B, and C are said to be public facilities. Facilities D and E are local to SUBSYSTEM 1, global to PART 1 and PART 2. SUBSYSTEM 2 and its inner blocks are not aware that facilities D and E exist; no reference to D and E may appear in the description of SUBSYSTEM 2.

Facilities H and I are local to PART 1; no other block of Fig. 2-1 may control or sense these facilities. PART 2 has its own facility I which may be of a very different hardware nature than facility I of
PART 1. PART 1 and PART 2 each control and sense their own facility I.

Similarly, PART 2 controls and senses its local facility J as does ASSEMBLY for its local facility J, which is global to CARD and hence can be controlled and sensed by CARD. References to K within CARD pertain to the most locally declared facility K, e.g., the one declared within CARD.

Permitting the same identifier to be used in parallel blocks allows designers working in parallel on the blocks to select without restriction names that appeal to them. If parallel blocks must communicate, facilities global to them must be established and assigned unique names. The designers of the parallel blocks must know and use these global names. Thus in Fig. 2-1 SUBSYSTEM 1 and SUBSYSTEM 2 may communicate via A, B, or C. PART 1 and PART 2 may communicate via D or E, or via A, B, or C.
2.1 SYNTAX RULES

VARIABLES:
Variable name may contain 1 to 8 characters, the first of which must be alphabetic. The remaining characters must be letters or digits.

Examples: MULT
          SYS1
          COMPLMNT

CONSTANTS:
Constants take the general form nRk. n is the number in base R (R=D for decimal, 0 for octal). k is the number of bits required for the representation, k ≤ 32. k is decimal.

Examples:

<table>
<thead>
<tr>
<th>Representation</th>
<th>Binary equivalent</th>
</tr>
</thead>
<tbody>
<tr>
<td>1D2</td>
<td>01</td>
</tr>
<tr>
<td>5D4</td>
<td>0101</td>
</tr>
<tr>
<td>20D5</td>
<td>10100</td>
</tr>
<tr>
<td>203</td>
<td>010</td>
</tr>
<tr>
<td>2006</td>
<td>010000</td>
</tr>
<tr>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
</tr>
</tbody>
</table>

2.2 DECLARATION STATEMENTS

The general format of a declaration statement is <DT> body.

The declaration type (device) is enclosed in angle brackets and the period terminates the declaration. Body consists of a list of items separated by commas. Following devices are allowed:
Terminal  Sets of wires
Registers  Sets of synchronized flip-flops
Memory   Sets of synchronized flip-flops
Latches  Sets of asynchronous latches
Time    Clock
Delay   Delay elements
Boolean  Combinational logic
Element  Off-the-shelf components

The device type can be abbreviated to the first two characters.

Examples:

<TE> X, Y(4), Z(0:2), W(3,4:1), A(12) = B + C(0:10) identifies a single wire X, four wires Y₁, Y₂, Y₃, Y₄ with Y₁ on the left, 3 wires Z₀, Z₁, Z₂ and 12 wires corresponding to W, placed in 3 rows, ith row of wires numbered W₁₄, W₁₃, W₁₂, W₁₁. The subscripts always have a left to right interpretation. A single subscript n indicates the range 1 to n while a range m:n indicates n to m left to right. In the above declaration, A₁ is also named B, A(2:2) are named C(0:10). + is the concatenation operator. The concatenation of B and C is a 12 bit terminal A with the most significant bit same as that of B and the least significant 11 bits same as those of C.

Register and Latch Declarations

<RE> IR(16) = OP(0:3) + IX(1 3) + ADRS(9), X(12). declares a 16 bit register IR and a 12 bit register X.
IR is identified with 3 subregisters OP, IX and ADRS.

<LA> BUF(4).
declares a set of 4 latches BUF.
<RE> A(8).
declares an 8 bit register, bits numbered from 1 to 8, left to right.

Memory DECLARATION

<ME> M(X:Y).
declares X words (numbered from 0 to X-1) of Y bits each (numbered 1 through Y).

<ME> MP(256:8).
declares a 256 word memory, 8 bits/word.

References to the memory must be of the form M(MAR) where MAR is the same register in all references to M. MAR is declared in a RE declaration. Only full words may be accessed from memories.

Time DECLARATION

<TI> A(1E-6), Q(20E-9) \$2s\$.
declares a single phase clock A with a 1 microsecond period and a two-phase clock Q with 20 nanosecond period.

<TI> P.
declares a single phase clock with an arbitrary time period (unit).

Delay DECLARATION

<DE> P(10E-9), Q(5E-7).
declares two delays P with 10 nanoseconds and Q with .5 microsecond.
The context in which the Delay element is referenced determines whether its input or output terminal is used.

Boolean DECLARATION

<BO> Identifier = Boolean expression.

Examples:

<TE> A, B(5), C(0:4), D(6, 5:1).

<BO> D(4) = B+C, D(5) = A*B.
declares that the fourth row of D is formed by ORing terminals B and C i.e. \((D_45 = B_1 + C_0\) etc.) bit by bit; the fifth row of D is a bit by bit AND of A and B. Since A is 1 wire and B is a set of 5 wires, A is fanned out to combine with each bit of B.

**Element Declaration**

Enables the description of an element in the system whose logical specifications are unknown or impertinent.

For example,

```
<EL> JKFF (Q1,NQ1: C, J1, K1), COUNT (K(5:1), ZERO: UPDWN, CLK).
```

declares an element JKFF with 3 inputs C, J1, K1 and two output Q1 and NQ1; and an element COUNT with two inputs and 6 outputs. The only information available on these black boxes is the input/output terminals.

**2.3 Operations**

Table 2.1(a) shows the operations allowed and their hierarchy; Table 2.1(b) shows three special operators. "-" is used to show the connections while \(<-\) indicates a data transfer from one facility to the other \(\rightarrow\) is equivalent to a "GOTO", used to show a state transition.

The extension operator "$\$" creates k copies of the terminal or terminal set offered as its left operand.

The selection operator \(','\), selectively complements, or not complements the bits of the facility (left hand operand) depending on the value of the corresponding bit in kDn is a 0 or 1.
For example, A' 10D5 is equivalent to

```
1
2
A
3
A'01010
4
5
```

The operator preceding the reduction operator (/) determines the nature of the reduction on the right hand operand of /. Six types of reductions are possible. For example, given a signal A,

*/A implies

\[ A \]

If A is a 3 bit signal,

*/A' 2D3 implies

\[ \text{Selection} \]

\[ \begin{array}{c}
\text{A} \\
\text{Reduction}
\end{array} \]
Boolean expressions (Be) can be formed by using the operators and variables in the usual manner. Parenthesis could be used where there is an ambiguity. The expressions are evaluated from left to right following the operator hierarchy.

Conditional operations have the format

\[ !BE! \text{ OP \_1} \]
\[ !BE! \text{ OP \_1 ; OP \_2} \]

The first form implies: If the value of BE is 1, perform OP\_1; the second form implies: If BE is 1, perform OP\_1 else perform OP\_2. "If ... then" operations can be nested:

\[ !A! \text{ ! B! OP \_1 ; ! C! OP \_2} \]

2.4 IF - VALUE CLAUSE

"\" is used for "IF" and "\#va" is used for the value in an IF-value clause. For example;

\[ B = \text{ C \#0 DO \#1 D1 \#2 D2.} \]

implies that DO is connected to B if the value of C is 0, D1 is connected to B if the value of C is 1, etc.

As another example,

\[ X \text{ \#OD2 A<-B \#1D2 A<-C \#2D2 A<-AB \#3D2 A<-AC.} \]

describes a 4 way conditional transfer operation into A depending on the
value of X.

2.5 IDENTIFIER

Identifier declaration enables the naming of a group of operations so that they do not have to be written repeatedly (equivalent to MACROS). The general format of Identifier declaration is,

<ID> list.

where list takes the form

id = compound facility
id = (CSOP)

For example, <ID> X = C(2:10) &1. names the compound facility C(2:10)&1 to be X. Then, any reference to X is expanded into C(2:10)&1.

For example, S = R & X. is equivalent to S = R & C(2:10) &1.

(A compatible set of operations (CSOP) is a set of operations separated by commas. It must be possible for the hardware to perform all these operations simultaneously.)

The order in which the operations are listed is of no consequence.

For example,

<ID> A = (Y & X, Z & Z(2:5) & AZ(1)),
B = (Y & X, Z & Y).

names two CSOPS. Note that the operations Y & X and Z & Y in B are simultaneous and are compatible.

2.6 OPERATOR DECLARATION

Blocks of combinational circuitry can be defined with the Operator declaration. The body of the Operator declaration consists of a Boolean declaration and perhaps a Terminal declaration. Boolean equations in the body of the Boolean declaration include Boolean expressions which
may involve conditions and be relatively complex. References in these
Boolean equations may be made to (1) facilities global to the OPerator
declaration. (2) local terminals declared within the OPerator declara-
tion by a TErnal delcaration, and (3) terminals delcared and dimension-
ed in the head of the OPerator declaration. The TErnal declaration
may be used to define local terminals of the operator, and must be used
to dimension "dummy" identifiers listed in the heading, if any.

The head of the Operator declaration consists of one or a list
(separated by commas) of identifiers with or without an argument list
enclosed in $s, with or without parenthetic subscript ranges. Permitted
syntactic forms for heads are:

\[ \text{id}_1, \text{id}_2(i_2), \text{id}_3$ $\text{X}_1, \text{X}_2, \ldots \text{X}_k$, \text{id}_4 (i_4)$

\[ \text{X}_1, \text{X}_2 \ldots \text{X}_k$

where subscript ranges can also be placed within the parenthesis. The
identifiers name the combinational logic blocks and their output termi-
nals. Parenthetic integers dimension the output terminal sets with the
same syntax and semantics as in TErnal declarations. The arguments
are local dummy identifiers of input terminals of the combinational
blocks. Such dummy identifiers must be dimensioned via a local terminal
declaration within the OPerator body.

As an example of a time-shared operator block, ALU is declared
below. This combinational block is able to add two 16-bit binary
sequences presented to it on lines X and Y or form their bit-by-bit
EXCLUSIVE-OR. Input signal F determines which task is performed. The
carry into rightmost full-adder must also be presented to the unit.

\[ \text{OP} \quad \text{ALU}(16) \quad \text{S} \quad \text{X}, \text{Y}, \text{CIN}, \text{FS}$

\[ \text{TE} \quad \text{X}(16), \text{Y}(16), \text{CIN}, \text{F}, \text{C}(16) = \text{CXQCC}(15). \]
<BO> C=XY + CCE CIN (X+Y),
ALU = (!F! X@Y @ CCE CIN; X@Y) ...(end of BO, end of OP)

Note the inline comment capability of DDL (end of BO, end of OP).
Suppose the following declaration is global to ALU,
<RE> ACC(16), MBR(16), COUNT (12).
we can define several operations using ALU as following:
!LDA! ACC <- ALUSO,MBR,0,0$
!ADD! ACC <- ALUSACC,MBR,0,1$
!SUB! ACC <- ALUSACC,AMBR,1,1$
!KNT! COUNT <- ALU(5:16) $OD4$COUNT,0,1,1$
!XOR! ACC <- ALUSACC,MBR,0,0$

2.7 STATE DECLARATION

DDL views the operation sequencing (control) circuitry as a finite state machine. Each state (step) of the control circuitry is described by a State declaration:

<ST>State List.

State list consists of a list of state statements (without separating commas). Each state statement has one of the following forms:

Sid (n): csop.
Sid (n): Be: csop.

Sid is a simple unsubscripted identifier. n is the decimal state assignment.csops include the state change operations using the state transition operator ->.

In the first form, csop is performed whenever the automaton is in the state Sid.
In the second form, csop is performed when the automaton is in Sid and also Be is satisfied. The automaton waits in the state till Be is
A 15 bit multiplier control can be described as following:

```
<ST>  S0(0):MPY:ACC<-0, CNT<-15D4,->S1.
       S1(1):->S2, DECR$ CNT$ !Q(15) ! ACC<-ACC+R.. 
       S2(2):SHR$ACC@Q$, I+/CNTI - >S1;S0 ... 
       (end of conditional, end of S2, end of ST)
```

SHR is shift right (zero fill) operator and DECR is a decrement operator assumed to be defined using <OP> declaration.

2.8. AUTOMATON and SYSTEM DECLARATIONS

Relatively independent disjoint portions of a digital system are identified as automata in DDL with syntax.

```
<AU> head body.
```

The AUtomaton declaration is the most complex type of declaration of DDL. Its head may take any of four forms, for example;

```
auid:
    auid:csop
    auid:Be:
    auid:Be:csop
```

First, an automaton identifier, auid, may be subscripted, but may not include parenthetical arguments; it names the block only. A compatible set of operations may be included in the head of an automaton. These operations are to be performed whenever the Be of the heading, if any, is satisfied. Conditional as well as unconditional operations may be included in this heading csop, so whether a specific operation is performed or not may depend on conditions throughout the automaton or system.

Be in the heading of the AUtomaton declaration is a condition on
all operations declared throughout the body of the declaration except
connection operations. Usually Be is the clock signal that synchronizes
the automaton. It is generally unnecessary and undesirable to include
such global conditions as clock signals in combinational circuits; in
fact, signal propagation in combinational networks usually precedes
clock pulses. If a clock with n phases is used to synchronize an autom-
aton, then a dimensional Be or a concatenation of n Bes appears in place
of the single Be in the AUtomaton declaration head.

The body of an AUtomaton declaration consists of other declarations.
Each of these declarations is terminated with its own period; punctuation
is not placed between them. The following declaration types may appear:

<ME>, <RE>, <LA>, <TE>
<TI>, <DE>, <OP>, <EL>, <ID>, <BO>, <ST>

ME, RE, LA, TE, TI, DE, AND EL declarations are used to declare the
existence of local facilities of the automaton. The OPerator and BOolean
declarations specify combinational blocks and interconnections of facil-
ities. The IDentifier declaration may be used to simplify or clarify the
overall AUtomaton declaration. The STate declaration is usually used to
specify the operations of the automaton. If the STate declaration is not
used, then all operations appear in the csop of the AUtomaton declaration
head.

The SYstem declaration has syntax identical to the AUtomaton decla-
ration. The system is identified in the head. Global conditions and
csop may be specified also. The body of a SYstem declaration may contain
AUtomaton declarations as well as all other types of declarations, but
STate declarations must appear within AUtomaton declarations. Public
facilities are declared with ME, RE, TE, etc., declarations outside of all
Automaton or Operator declarations.

Example:

A multiplier controller is described below to illustrate the SYstem and AUtomaton facilities. The counter is treated as a separate automaton. Perhaps other unspecified automaton of SYSTEM 1 can use the counter when automaton MC is not.

<SY> SYSTEM 1:

<RE> ACC(15), Q(15), R(15).

<TE> SET, DEC, DONE, MPY.

<TI> P(1E-7).

<AU> CPU: P:

<ST> .

Q17: DONE: Q <- Multiplier,

. R <- Multiplicand, MPY = 1.

. . (end CPU)

<AU> MC: P:

<ST> S0: MPY: ACC <- 0, SET = 1, -> S1.

S1: -> S2, DEC = 1, !Q (15)! ACC <- ACC+R..

S2: ACCQ <- SHR$ACCQ$ !DONE! -> S1 ...

<AU> K: P:


T(0): DONE = 1, !SET! -> T(15); -> T(0)... (end SY)
Automaton CPU is shown only as placing the multiplier and multiplicand in public registers and issuing command MPY to multiplier control MC. If the counter automaton K is idle, it will be issuing DONE = 1. CPU waits in its state Q17 until this condition is satisfied (perhaps K is still doing a job for some other automaton). MC clears ACC, but the counter is initialized by SET = 1. Specifically SET = 1 will cause K to go from its state T(0) to T(15) where it will remain until it is told to decrement via public terminal DEC. MC tests the multiplier, adds or not and shifts repeatedly until it is informed by K via public terminal DONE that all multiplier bits have been examined. In the example above interacting automata MC and K operate in parallel.

NOTE: The "For clause" shown in the Automaton K for the decrement operation \( i=1;15 \) \( T(i) \) : DEC: \(-\rightarrow T(i-1) \) is not allowed in the present version of the DDL software. This statement has to be broken up into;

\[
\begin{align*}
T(1) &: \text{DEC: } \rightarrow T(0) \\
T(2) &: \text{DEC: } \rightarrow T(1) \\
& \vdots \\
T(15) &: \text{DEC: } \rightarrow T(14)
\end{align*}
\]

SHR is a single argument operator (assumed to be declared earlier) that shifts the argument one bit right, and fills zero on the left.
TABLE 2.1(a) : OPERATORS

<table>
<thead>
<tr>
<th>OPERATOR</th>
<th>SYMBOL</th>
<th>TYPICAL SYNTAX</th>
<th>COMMENTS</th>
</tr>
</thead>
<tbody>
<tr>
<td>Extension</td>
<td>$</td>
<td>A$^k</td>
<td>k copies of A</td>
</tr>
<tr>
<td>Concatenation</td>
<td>$^2$</td>
<td>A$^2$B</td>
<td>Bit by bit</td>
</tr>
<tr>
<td>Complementation</td>
<td>$\Lambda$</td>
<td>$\Lambda$A</td>
<td>Bit by bit complement</td>
</tr>
<tr>
<td>Selection</td>
<td>'</td>
<td>A'kDn</td>
<td>Selective complement</td>
</tr>
<tr>
<td>Reduction</td>
<td>/</td>
<td>p/A</td>
<td>A_1pA_2p...pA_n, where p is one of these: $^<em>$, $^\Lambda$, $^\Lambda+$, $^\Lambda@$, $^\Lambda+$, $^\Lambda^</em>$.</td>
</tr>
<tr>
<td>AND</td>
<td>*</td>
<td>A*B</td>
<td>Bit by bit</td>
</tr>
<tr>
<td>NAND</td>
<td>$^\Lambda*$</td>
<td>A$^\Lambda*$B</td>
<td>Operations</td>
</tr>
<tr>
<td>NOR</td>
<td>$^\Lambda+$</td>
<td>A$^\Lambda+$B</td>
<td></td>
</tr>
<tr>
<td>XNOR</td>
<td>$^\Lambda@$</td>
<td>A$^\Lambda@$B</td>
<td></td>
</tr>
<tr>
<td>XOR</td>
<td>$@$</td>
<td>A@$^B$</td>
<td></td>
</tr>
<tr>
<td>OR</td>
<td>$+$</td>
<td>A+$^B$</td>
<td></td>
</tr>
</tbody>
</table>

TABLE 2.1(b) : SPECIAL OPERATORS

| CONNECTION |   |
| TRANSFER | <- |
| GO TO | -> |

NOTE: Refer to Chapter 3 (The Translator) for variations of these Operators.
3. THE TRANSLATOR (DDLTRN) [36]

DDLTRN is a program that translates a DDL description of a digital system into 1) a DDL description that consists of Boolean equations and register transfer statements in the heading of a system declaration only, and 2) a tabulation of facilities and subfacilities declared in the DDL description and/or defined in the translation process. Some modifications of DDL recognized by DDLTRN are listed below. The translation process is briefly discussed and illustrated in Section 3.1.

1) The following operators are changed to accommodate the SEL-32 peripherals:

<table>
<thead>
<tr>
<th>DDL Operator</th>
<th>Key Punch</th>
<th>CRT Terminal</th>
<th>Printer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Concatenation ( &amp; )</td>
<td>&amp;</td>
<td>[ ]</td>
<td>[ ]</td>
</tr>
<tr>
<td>Complement ( A )</td>
<td>( r )</td>
<td>( A )</td>
<td>( + )</td>
</tr>
<tr>
<td>IF - THEN !</td>
<td>!</td>
<td>]</td>
<td>]</td>
</tr>
<tr>
<td>IF - VALUE ( \mid )</td>
<td>( \mid )</td>
<td>]</td>
<td>]</td>
</tr>
</tbody>
</table>

The other operators of DDL are compatible with the peripherals of SEL-32 and remain the same.

2) Comment declarations end with a left angle bracket \(<\).

3) Values in "If-value" clauses are limited to a single integer values. Ranges, lists and else (;) values are not permitted.

4) Concatenation operands must be simple facilities with or without subscripts, or binary strings.

5) State assignments are specified in decimal following the state identifier of each state statement, e.g., "S1(2):..."

6) Automata names are used as state sequencing register names and thus should be dimensioned in the \(<AU>\) declaration head, e.g., "<AU> CPU (5): P:..."
7) DDLTRN accepts Flag declarations with syntax: <Flag> list. List consists of integers, and/or integers preceded by the complement symbol (A), separated by commas. Each integer specifies the setting of a flag. Each complemented integer specifies that the corresponding flag is to be reset. Table 3.1 summarizes the significance of set flags and the default states of the flags.

8) Identifiers defined in Identifier declarations must not be subscripted.

### Table 3.1 Flag Integers

<table>
<thead>
<tr>
<th>Flag</th>
<th>Significance</th>
<th>Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Print Source Card Images</td>
<td>Set</td>
</tr>
<tr>
<td>2</td>
<td>Print Declared Facilities and Operations</td>
<td>&quot;</td>
</tr>
<tr>
<td>3</td>
<td>Print DDL string after Pass 2</td>
<td>Reset</td>
</tr>
<tr>
<td>4</td>
<td>&quot; &quot; &quot; &quot; 3</td>
<td>&quot;</td>
</tr>
<tr>
<td>5</td>
<td>&quot; &quot; &quot; &quot; 4</td>
<td>&quot;</td>
</tr>
<tr>
<td>6</td>
<td>&quot; &quot; &quot; &quot; 5</td>
<td>&quot;</td>
</tr>
<tr>
<td>7</td>
<td>&quot; &quot; &quot; &quot; 6</td>
<td>Set</td>
</tr>
<tr>
<td>8</td>
<td>Print F Table after Last Pass</td>
<td>Reset</td>
</tr>
<tr>
<td>9</td>
<td>Print Encoded string after Last Pass</td>
<td>&quot;</td>
</tr>
<tr>
<td>10</td>
<td>Execute through Pass 2 only</td>
<td>&quot;</td>
</tr>
<tr>
<td>11</td>
<td>&quot; &quot; &quot; 3 &quot;</td>
<td>&quot;</td>
</tr>
<tr>
<td>12</td>
<td>&quot; &quot; &quot; 4 &quot;</td>
<td>&quot;</td>
</tr>
<tr>
<td>13</td>
<td>&quot; &quot; &quot; 5 &quot;</td>
<td>&quot;</td>
</tr>
<tr>
<td>14</td>
<td>&quot; &quot; &quot; 6 &quot;</td>
<td>Set</td>
</tr>
</tbody>
</table>

3.1. THE TRANSLATION PROCESS

DDLTRN is the result of a research effort to develop efficient language translation algorithms. As a result it emphasizes translation
efficiency rather than error detection and control. Neither the syntax of supplied DDL descriptions nor the translation process itself are checked in detail.

A DDL description is stored as a single string in a singly linked list in memory. Operator and punctuation symbols are represented by codes. As processing proceeds facility names and subscript ranges are also encoded to shorten the string and hence the time required to pass over it.

Facts about declared facilities such as name, subscript range, type, etc. are recorded in a facility table F. Translation consists of passing over the DDL string a number of times. With each pass the DDL string and F table are modified according to unique rules. Six main passes may be identified by the user: The DDL string and F table may be printed after any of these main passes.

Pass 1 -- Facilities Identified

Data cards bearing a DDL description are read and echo printed. All blank columns are ignored; all card columns 1 - 80 are examined. Declared facilities are entered in the F table. TIme, REgister, MEmory, LAtch, TErnal and DElay declarations are removed from the DDL description, as are all COMMENT declarations and parenthesized comments. Identical primary names declared in nested or parallel blocks are made unique by appending a double quote ("), and integer. Identical names declared in the same block are rejected, of course.

Pass 2 -- Syntax Reduced

Names and binary strings in connection and register transfer operations are encoded. Secondary names (names appearing on the right of an equal sign in a TErnal, REgister, etc. declaration) are replaced with
their subscripted primary name equivalents. Identifiers from IDentifier declarations are replaced in operations and expressions serving as conditions on operations with the symbol string they represent. The syntax of OPerator, B0olean and SState declarations is removed, the connection operations being transferred to the head of the enclosing AUtomaton or SYstem declaration. SState statement syntax is replaced with "if-then" conditions on operations. OPerator call arguments are transformed to connection statements. Compound Boolean expressions serving as conditions on operations are replaced with terminals of unit dimension. These new terminals are connected to the Boolean expressions via connection operations inserted in the head of the enclosing AUtomaton or SYstem declaration.

Pass 3 -- Conditions Distributed

"If-then" and "if-value" conditions on sets of operations are combined and distributed over the members of the set so that each operation appears as the body of a simple "if-then" clause. "Go-to" operations are converted to conditional transfers of a constant (the state assignment) to the state sequencing register (the enclosing automaton). AUtomaton syntax is eliminated by recognizing the global condition, if any, and distributing it as a clocking condition on all register transfer and memory operations within the AUtomaton declaration.

Pass 4 -- Concatenation Removed

All concatenation operations except those that form operands for reduction operators are eliminated by breaking operations into operations on sub-facilities formed by partitioning operand facilities according to the dimensions of the concatenation operands.
Pass 5 -- Operations Gathered

All connection and transfer operations with the same data sink (left operand) are gathered into one compound operation.

Pass 6 -- Subfacilities Disjoined

Facilities with subfacilities serving as data sinks of connection and transfer operations are broken into disjoint subfacilities and a right-hand side of a connection or transfer operation is formed for each new subfacility.

3.2 AN EXAMPLE

System EX1 illustrates the use of secondary names and IDentifier declarations. Registers A and D of automaton A1 are each broken into sub-registers via secondary names in the REGISTER declaration. Ascending and descending subscripts are illustrated. Identifiers X, Y and Z name a new concatenation of the subregisters of D, a portion of one of these sub-registers, and a NOR reduction, respectively. A register A is declared in automaton A2 also. The operations of A2 all appear in the head portion of its AUTOMATON declaration.

The listing obtained after Pass 1 summarizes the declared facilities and their relations. Since two A registers are declared in parallel blocks, the name of one is changed to A"l so that the two may be distinguished. The declared operations are listed with indentation used to indicate the nested relations of blocks. Block structure errors would be easily identified.

Pass 2 replaces secondary names and identifiers with their primary equivalents. A careful examination of the results after Pass 2 indicates that operation A-X in state S becomes A-FWE when X is replaced. Then
secondary names are removed giving \( A \rightarrow D(4:1) \rightarrow D(8:5) \). The operations of state \( T \) require that secondary names \( F, B, C \) and \( E \) be replaced with their primary equivalents. Then \( Z \) within "if-then" punctuation is replaced with \( \neg +/Y \) is replaced with \( \neg +/F(3:2) \) is replaced with \( \neg +/D(2:1) \). Note that state statement syntax is also converted to "if-then" syntax in Pass 2.

A state decoder network on automaton register \( A1 \) is prescribed by equations in the head of the SYstem declaration at this point.

Pass 3 distributes conditions over sets of operations and removes Automaton declaration syntax. The results listed indicate that five internal signals named "double-quote-integer" have been formed in order to simplify the expression of conditions on operations. Each of the conditioned operations can be traced back to the source DDL description. "Go to" operations are converted to conditioned transfers to the automaton register.

Pass 4 eliminates the concatenation operations. As a first example observe that

\[
!P*S! A \leftarrow S*D(4:1) \rightarrow D(8:5).
\]

is broken to

\[
!P*S! A(1:4) \leftarrow S*D(4:1), \\
!P*S! A(5:8) \leftarrow S*D(8:5).
\]

Pass 5 gathers operations with the same left operand. The operations

\[
!P*S! A(1:4) \leftarrow S*D(4:1), \\
!P*S*5! A(1:4) \leftarrow "5*D(4:1).
\]

are gathered to

\[
!P*S+P*S*5! A(1:4) \leftarrow S*D(4:1) + "5*D(4:1).
\]

No logic minimization or even simplification is performed as part of the gathering process.
In Pass 6 the A and D registers of automaton A1 are partitioned and transfer statements are developed for each subfacility. Pass 5 provides the following transfers to the A register or some part of it.

\[ !P*S + P*5 ! A(1:4)<-S*D(4:1) + "5*D(4:1)., \]
\[ !P*S + P*5 ! A(5:8)<-S*D(8:5) + "5*D(8:5)., \]
\[ !P*3 ! A<-"3*D. \]

The last of these operations involves the entire A register; the others involve a part of it. Pass 6 partitions the A register to A(1:4) and A(5:8), and forms the correct transfers to each of these subfacilities.

The F table as it appears after Pass 6 is listed as the final result of this example. Facility names are followed by left and right subscripts and facility dimensions. The next column indicates the type of the facility with negative entries (-1 for SYstem, -6 for REgister, -9 for TErminal, etc.). Positive entries point to the row of the parent facility. The final columns point to the beginning and ending points of facility operation statements in the DDL string.
ASSI - FACILITIES IDENTIFIED

DECLARED FACILITIES

<SY> CAL
<TI> P(1:1)
<AL> A1
<HL> E(1:2) = E(1:2) I C(1:2)
<IL> X(1:1) = E(1:1) I (3:2)
<ST> S
T
<AL> A2
<AL> A**1(1:24)
T(1:24)

DECLARED OPERATIONS

<SY> CAL:
<AL> A1:
<ST>
S: A**A, => T,
T: E**H(I(7)), C**10(L),
U: => S,
V:

<AL> A2: P: A**A, => A,
PASS2 = SYNTAX REVEALED

<SY> EX1: S = \#/A1\'U02 , E = \#/A1\'1\'C , L = \#/A1\'2\'C

<CAV>

41: P:
   1(2:1) A = \#(4:1) L(4:5) , => 1;
   1(4:1) L(4:1) = \#(A(1:3) A(4:1) , \#(c(5:2) = \#(L(1:3) ,
   1(4:1) L(4:1) = \#(c(5:2) = \#(L(1:3) ,
   1(2:1) => \#(S) ;
   1(2:1) #=1

<CAV> 42: P: A"1<->C , =<A"1> ..

PASS3 = CONDITIONS VISIMIABLE

PASS4 = CALCULATION REVEALED

<SY> EX1: S = \#/A1\'U02 , E = \#/A1\'1\'C , L = \#/A1\'2\'C

<CAV>

1(2:1) A = \#(4:1) L(4:5) ,
1(4:1) L(4:1) = \#(A(1:3) A(4:1) ,
1(2:1) => \#(S) ;
1(2:1) #=1

PCR 41: S = \#(A(1:4) L(4:5) ,
PCR 42: S = \#(A(5:2) L(4:5) ,
PCR 43: S = \#(A(5:2) L(4:5) ,
PCR 44: S = \#(A(5:2) L(4:5) ,
PCR 45: S = \#(A(5:2) L(4:5) ,
PCR 46: S = \#(A(5:2) L(4:5) ,
PCR 47: S = \#(A(5:2) L(4:5) ,
PCR 48: S = \#(A(5:2) L(4:5) ,
PCR 49: S = \#(A(5:2) L(4:5) ,
PCR 50: S = \#(A(5:2) L(4:5) ,
PCR 51: S = \#(A(5:2) L(4:5) ,
PCR 52: S = \#(A(5:2) L(4:5) ,
PCR 53: S = \#(A(5:2) L(4:5) ,
PCR 54: S = \#(A(5:2) L(4:5) ,
PCR 55: S = \#(A(5:2) L(4:5) ,
PCR 56: S = \#(A(5:2) L(4:5) ,
PCR 57: S = \#(A(5:2) L(4:5) ,
PCR 58: S = \#(A(5:2) L(4:5) ,
PCR 59: S = \#(A(5:2) L(4:5) ,
PCR 60: S = \#(A(5:2) L(4:5) ,
PASS 5 -- PERILOUS GATHER

PASS = SUBFACILITIES DISJOINED

PASS = SUBFACILITIES DISJOINED
PASSS--OPERATIONS GATHERED

```plaintext
<SY> EX1: S=*A1'102
  f=*=a1'102
  v=*=a1'202
''1=1*+/(2:1),
  "2=1*+/(1+0(2:1)),
  "3=0*(1+0(2:1)'102 ),
  "4=0*(0(2:1)'102 ),
  "5=0*(0(2:1)'202 ),
  "P*3 + P="5) A(1:4)="5*U(4:1) + 5*U(4:1) .
  "P*8 + P="5) A(5:8)="5*U(4:5) + 5*U(4:5) .
  "P*5 + P="2 + P*U) u1="3*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="1 + P*U) u1="3*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="2 + P*U) u1="3*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="1 + P*U) u1="3*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="2 + P*U) u1="3*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
```

PASSS--SUBFACILITIES DISJOINED

```plaintext
<SY> EX1: S=*A1'102
  f=*=a1'102
  v=*=a1'202
''1=1*+/(0:c:1),
  "2=1*+/(0(2:1)),
  "3=0*(0(2:1)'102 ),
  "4=0*(0(2:1)'102 ),
  "5=0*(0(2:1)'202 ),
  "P*3 + P="S) A(1:4)="S*U(4:1) + 5*U(4:1) .
  "P*8 + P="S) A(5:8)="S*U(4:5) + 5*U(4:5) .
  "P*5 + P="2 + P*U) u1="S*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="1 + P*U) u1="S*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="2 + P*U) u1="S*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
  "P*5 + P="1 + P*U) u1="S*10c + 1*U1c + 4*U1c + 4*U1c + 4*U1c .
```

ORIGINAL PAGE IS
4. THE SIMULATOR (DDLSIM) [35]

DDLSIM is a program for simulating digital systems described using DDL. The simulator has a simple, powerful and completely free-format command language that provides the user with complete control over the simulation process without requiring that the DDL system description be modified. DDLSIM does very extensive error-checking of described systems, simulation control cards, and the simulation process itself. Self-explanatory messages that pinpoint errors are issued.

Digital systems to be simulated are first described in DDL. This description is translated by DDLTRN into a set of Boolean equations and Register Transfer expressions. These can be used for implementation or simulation of the digital system. They, together with other data structures and tables established by DDLTRN constitute the system description required by DDLSIM. This description is pre-processed by the simulator to establish data structures and tables that permit more efficient and controlled simulation.

The original and translated DDL descriptions of a system neither contain any information for controlling simulation nor do they supply any input data for simulation. These items are supplied by a second source to DDLSIM, a simulation deck. This deck consists of simulator control declarations described using a simulator command language that is not unlike DDL. Twelve different declaration types are available for selecting options and supplying control information, parameters, and data for simulation. Every simulation job consists of:

1. processing the system description,
2. processing the simulation deck, and
3. simulation of the system.
The following notational conventions are used in subsequent sections to describe the syntax of translated DDL and to define control language syntax.

Script characters - an item of the language. Item α will be defined in terms of items β and γ with notation

\[ α : β, γ \]

which is read "an α is a β or a γ."

[ ] - appearance of the enclosed syntactic structure is optional

[ ]^n - the enclosed syntactic structure may be repeated an arbitrary number of times or not at all.

Blanks have no significance in syntax descriptions just as they have no significance in DDL or the DDLSIM control language.

4.1 SIMULATION MODELS

As mentioned earlier, Boolean equations (BE) and Register Transfer Expressions (RTE) generated by DDLTRN constitute the system description required by DDLSIM. The models of combinational networks and registers used by DDLSIM is the subject of this section.

4.1.1 Terminals, Element Inputs, and State Terminals

The terminals, element inputs, and state terminals declared in a system are described using BEs. In addition, DDLTRN generates BEs for a number of intermediate signals. All such BEs constitute the combinational portion of a system. They are first sorted into an ordered list according to the level of their operands, i.e., if a terminal A is used in the BE for another terminal B, A will appear before B in the sorted list. However, if the system contains loop(s) in its combinational portion, it is not possible to sort the equations in this
manner. In such cases the BEs constituting the loop(s) (or loop equations) are separated from other BEs. The remainder of the BEs are then sorted into an ordered list as described. Loop equations are then added to the sorted list at an appropriate point.

During simulation the combinational portion of a system is simulated at the BE level. BEs can vary from a simple sum-of-products form to the most complex and compound of forms. The BEs are evaluated in the order established by sorting with the loop equations being simulated repeatedly until their output values stabilize. Failure of a loop to stabilize after a fixed number (determined by the characteristics of the loop equations) simulations, indicates instability in the loop. In such a case a warning is issued to the user and the simulation is continued with the last computed values for the loop equations taken as their final values. Thus DDLSIM also permits the simulation of systems having loops in their combinational portion.

4.1.2 Delays

The delays declared in a system (using <DE> declarations of DDL or DDLSIM) are also described using BEs. These delays are assigned their delay time periods (Δs) using <Delay> declaration of DDLSIM (see Sec. 4.2.4). All the delay facilities are assumed to be inertial delays, i.e., an output signal(s) will assume a new value(s) Δ time units after it's input prescribes that change, if and only if the input signal prescribes that value for at least Δ consecutive time units. Unlike the BEs discussed above, the BEs for delays are not sorted in any particular order.

During simulation each delay is simulated at the BE level with specified inertial delay assigned to it's output. The new computed
value(s) for each delay is compared with its present output value(s). If they are different, a future event at \( \Delta \) time units from present simulation time \( T \) is scheduled to carry out the change(s) in the output value(s). However, if the BE does not continue to prescribe the change for at least the next \( \Delta \) time units, the scheduled event is cancelled and the output(s) of the delay remains unchanged.

It is possible to assign the same delay time \( t_d/2 \) (see Sec. 4.2.2, 4.2.5) to all the BEs for the combinational portion (see Sec. 4.2.1) of the system by setting flag number 13 (see Sec. 4.2.14) In such a case all these facilities become equivalent to delays. It is important to note that the delay time assigned to these BEs is the same for all of them, irrespective of their complexity.

4.1.3 Registers

The registers declared in a system are described using RTEs. RTE consists of a Condition Expression (CE) followed by a Transfer Expression (TE). RTEs generated by DDLTRN have the following general syntax:

\[
RTE : | CE | TE.
\]

\[
CE : C [+C]^n
\]

Condition term

\[
C : C_c \cdot C_L, C_L
\]

Clock condition

\[
C_c : \text{global condition in the heading of an } <AU> \text{ declaration of DDL, a clock declared in a } <TT> \text{ declaration of DDL.}
\]

Load condition

\[
C_L : \delta \text{ with } \delta_w = 1. \text{ (see Sec. 4.2.1)}
\]

\[
TE : \delta + E
\]

Load expression

\[
E : e [+e]^n
\]
Expression term \( e \):

\[ C_\ell \ast V_e \]

Load value \( V_e \): an expression

Example: \( P*LDX + P*ORXY + P*LDY \mid ACC = LDX*X + ORXY*(X+Y) + LDY*Y \).

In the example \( P \) is a clock; \( ACC, X, \) and \( Y \) are all registers having dimensions of 24; \( LDX, ORXY, \) and \( LDY \) are terminals declared using appropriate declarations. The CE in this example has three condition terms specifying the conditions for performing different register transfers on \( ACC \). All the register transfers in this case are carried out under the control of the same clock \( P \). In the RTE for registers declared as global facilities and used in different automata, each having a separate clock or global condition, the CEs may have different clock conditions \( C_\ell \). For each condition term \( C \) in the CE, there is a corresponding expression term \( e \) in the TE. When a load condition \( C_\ell \) becomes true (logic 1) and the corresponding clock condition \( C \) performs a 0-to-1 transition, the next-output value for the register is computed using the load value \( V_e \) from corresponding expression term \( e \). On the next 0-to-1 transition of the \( C_\ell \), this next-output value becomes the present-output value.

During simulation CEs for all the registers are evaluated only at certain event-times (see Sec. 4.3). On a 0-to-1 transition in the value for a CE, the corresponding \( E \) is evaluated and the computed value is stored as the next-output value for the register. On a 1-to-0 transition of the same CE at some future evaluation, the next-output value for the register becomes its present-output value. In order to make simulation fast and efficient, CEs are evaluated only at event-times at which 0-to-1 or 1-to-0 transitions of clock conditions take place. It is not possible to have a 1-to-0 and 0-to-1 transitions for the same CE at the same
simulation time $T$. It is possible to simulate asynchronous sequential systems using DDLSIM.

The simulation model used for a register is very similar to a GL (gate and latch) flip-flop. A logical OR of load conditions $C_L$ from $CE$ constitutes the Boolean equation for the GATE of the flip-flop, $E$ from $RE$ constitutes the LATCH equation for the flip-flop, and a logical OR of the clock conditions $C_C$ from $CE$ constitute the CLOCK of the flip-flop. (See the figure below)

### 4.1.4 Memories

The memories declared in a system are also described using RTEs. A RTE for a memory is similar to that for a register with an address specified for the facility $\delta$, i.e.,

```
memory $\delta : \delta(a)
```

address expression $a : \text{an expression}$

The simulation model used for memories is also similar to that used for registers. For memory-write operations the address expression $z$ is evaluated on a 0-to-1 transition of the associated $CE$ and the computed value is stored as the address of the memory location. On the next 1-to-0 transition of the condition expression $CE$, the contents of the addressed location are changed to the supplied value. Memory-read operations are instantaneous, i.e., contents of the referenced memory location are fetched immediately.

![Diagram showing a Gated Latch and GLFF circuits]
4.2 SIMULATOR COMMAND LANGUAGE AND DECLARATIONS

The DDLSIM command language consists of twelve different types of declarations for supplying parameters, input data, options and other control information necessary for simulation. The language is largely free of format restrictions. Card images are scanned in turn from left to right. Any declaration may start at any point and end at any later point in the card deck. A declaration can be continued on as many cards as necessary; more than one declaration can be supplied on the same card. The start of a declaration automatically ends the previous declaration. The last declaration in a simulation deck is ended by an End-Of-File (normally a card having '$' in the first column). In general, the order in which the declarations are specified is not important. It is possible to have more than one declaration of the same type. Everything following the vertical line character (|) on a card is treated as a comment, and is not processed as a part of a declaration. Scanning continues on the next card. This provides the capability of having in-line comments in a simulation deck.

Each card from a simulation deck is processed sequentially by the simulator. First it is printed together with it's sequence number. It is possible to suppress echo printing of the simulation deck by turning the list option off, i.e., resetting Flag 1.

Each simulator declaration has the general syntax

<Declaration-id> Body

Each declaration begins with a left angle (<) followed by a Declaration-id that identifies the type of the declaration. Only the first two characters of the Declaration-id are examined by the simulator. The Declaration-id is terminated by a right angle (>). All declarations except the
<Simulate> declaration have a Body following the heading.

4.2.1 Facilities

Facilities are defined here as in DDL to be any piece of hardware declared in a digital system including terminals, registers, memories, and assemblage of hardware, clocks, delays, etc. If a facility name $\delta_n$ exceeds 8 characters, only the last 3 characters are retained. If a facility has dimension greater than one, individual elements are identified by appending a non-negative integer subscript $S_1$ enclosed in parentheses to $\delta_n$. A range of elements of a facility is identified by using a DDL subscript range, i.e., $\delta_n(S_1 : S_2)$. A script letter $\delta$ will be used to represent a facility or a part of it.

$$\delta : \delta_n(S_1 : S_2), \delta_n(S_1), \delta_n$$

where

$$\delta_n(S_1) = \delta_n(S_1 : S_1)$$

$$\delta = \delta_n(S_\ell : S_r)$$

$S_\ell$ = subscript for leftmost element of $\delta_n$.

$S_r$ = subscript for rightmost element of $\delta_n$.

Facility width $\delta_W$ of a facility $\delta$ is defined as the total number of elements in it, i.e.,

$$\delta_W = \max(S_1, S_2) - \min(S_1, S_2) + 1$$

During simulation one machine word is used to store the values of facility $\delta$. The SEL 32 machine has 32 bits per word. Hence it is necessary that the facility width $\delta_W$ for any facility $\delta$ in the system not exceed 32, i.e., $\delta_W \leq 32$. However, $S_\ell$ and $S_r$ may have larger values; only their difference is restricted.

A facility list $\ell_\delta$ is defined as a list of facilities $\delta$ separated by commas, i.e.,

$$\ell_\delta : \delta[\delta]^n$$
Whether a specific facility can be used in a facility list for a specific type of declaration is determined by both the type $\delta_F$ of the facility and the type of the declaration. The following facility types exist for DDLSIM.

$\delta_F$: System clock, Register, Memory, Terminal, System delay, Element input, Element output, State terminals, Trigger, Simulation delay, Simulation clock, List name.

Every facility $\delta$ used in a DDLSIM declaration must satisfy exactly one of the following conditions:

1. $\delta$ is declared in the DDL description of the system.

2. $\delta$ is declared during the present simulation run using a $<$CLock$>$, $<$DElay$>$, $<$TRigger$>$, or $<$LIst$>$ declaration. The type of declaration in which $\delta$ appears determines its type $\delta_F$ which cannot be changed for the remainder of the simulation job.

3. $\delta$ is declared during any previous simulation run as discussed in 2 above.

4.2.2 Numbers and Data Lists

$T_{\text{MAX}}$: a decimal integer having the value $2^{31} - 1$.

$P_{\text{MAX}}$: a decimal integer having the value $2^{16} - 1$.

$t_{i,j}$: a decimal integer $n$ in the range $i \leq n \leq j$ where $i$ and $j$ are each non-negative decimal integers. Whenever $j$ is not specified $j = T_{\text{MAX}}$, is assumed; whenever $i$ is not specified $i = 0$ is assumed.

$n_{i,j}^p = t_{i,j}$ enclosed in parentheses
\( n^P_{i,j} : (n^i_{i,j}) \)

\( R \) - Repeat factor, a positive decimal integer

\( R : n_1 \)

A repeat factor \( R \) can be used before a data value or parameter value, i.e.,

\( R \times \text{value} \),

to indicate that the same value is to be repeated \( R \) times in the list.

\( T \) - Simulation time

\( T : n \)

\( t_d \) - Default time period

\( t_d : n_1, p_{\text{MAX}} \)

Data is described with the following syntactic structures.

\( d_B \) - a binary digit

\( d_B : 0,1 \)

\( d_0 \) - an octal digit

\( d_0 : 0,1,2,3,4,5,6,7 \)

\( d_D \) - a decimal digit

\( d_D : 0,1,2,3,4,5,6,7,8,9 \)

\( d_H \) - a hexadecimal digit

\( d_H : 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F \)

\( d_G \) - a general digit excluding the hexadecimal digits B and D.

\( d_G : 0,1,2,3,4,5,6,7,8,9,A,C,E,F \)

\( B \) - a binary number

\( B \) \([+,-]Bd_B[\ldots]^n, ([+,-]Bd_B[Bd_B[\ldots]^n] \)
0 - an octal number

\[ 0 : [+, -]0d_0 [d_0]^n, \ (0d_0 [d_0]^n) \]

\( \circ 
\)

D - a decimal number

\[ D : [+, -]dd_D [d_D]^n, \ (0dd_D [d_D]^n) \]

H - a hexadecimal number

\[ H : [+, -]dd_H [d_H]^n, \ (0dd_H [d_H]^n) \]

N - a binary, octal, decimal or hexadecimal number.

\[ N : [+, -]d_G [d_G]^n, \ (0d_G [d_G]^n) \]

Optional leading minus signs (-) before any of above five types of numbers specifies the 2's complement of the number. 1's complement encoded negative numbers are obtained by setting Flag 10 (see Sec. 4.2.13).

\( N_2 \) - Data value

\[ N_2 : B, O, D, H, N, \]

\( N_1 \) - Data spec

\[ N_1 : [R^*] N_2 \]

\( \mathcal{C}_d \) - a data list consisting of data specs separated by commas.

\[ \mathcal{C}_d : N_1 [., N_1]^n \]

Whenever a data value is specified as a number \( N \) without leading radix specification, the default radix is used for computing the value of number. The default radix of 8 (octal) can be changed to 2 (binary), 8 (octal), 10 (decimal), or 16 (hexadecimal) by setting flag numbers 2 thru 5 (see Sec. 4.2.14 respectively. Resetting these flags returns the radix to the default value of 8 (octal).

4.2.3 <Clock> Declarations

This declaration provides a means for specifying or changing the time period, pulse width and phase of clock facilities. It also permits users to declare new clocks to be used to control simulation input and
output activities. Syntax for this declaration is as follows:

\[
\text{<CLock> Body}
\]

\[
\text{Body} : \ [l_1,]^{n} l_1
\]

\[
\text{List} : l : l_c/l_c
\]

\[
\text{Clock list} : l_c : l_0, \text{where}
\]

\[
\text{Facility type} : \beta_c : \text{system clock, simulation clock}
\]

\[
\text{Time list} : l_t : t[.,t]^n
\]

\[
\text{Time spec} : t : [R^*] P [W[\theta]]
\]

\[
\text{Time period} : P : n_2^P, \text{MAX}
\]

\[
\text{Pulse width} : W : n^P P-1
\]

Phase \( \theta \) : \( n^P P-W \) \( 0 \),

Example: \(<CL> \text{CLOCK1(1:5), CLOCK2/2*100(30) (50)/, CLOCK1(6:10), CLOCK3/100,100(30)/}\)

Time period \( P \) - the \( P \) field specifies the time period of a clock. In the above example each clock has a time period of 100 in some arbitrary units.

Pulse width \( W \) - This is an optional field specifying the time \( W \) for which a clock remains at logic 1 during any clock period \( P \). For the remaining time \( (P-W) \) the clock remains at logic 0. When the pulse width is not specified along with the time period, the following default value \( W \) is used.

\[
W = \lfloor P/2 \rfloor
\]

In the example a pulse width of 30 units is supplied for both \text{CLOCK(1:5)} and \text{CLOCK2}. \text{CLOCK3} is assigned a pulse width of 30 units.

No pulse width is explicitly specified for \text{CLOCK(6:10)}, hence a default value of \( \lfloor 100/2 \rfloor = 50 \) units is used as the pulse width.

Phase \( \theta \) - At the start of a simulation run, i.e., \( T = 0 \) a clock with a
period \( P \) and the pulse width \( W \) is set to start at logic 0. It remains at logic 0 for the next \((P-W)\) time units; then a 0-to-1 transition takes place. For the next \( W \) time units, it stays at logic 1; then a 1-to-0 transition takes place and the cycle is repeated. The occurrence of the first and every subsequent 0-to-1 transition can be advanced relative to the starting of simulation by specifying the phase \( \theta \).

1. For phase \( \theta < P - W \) a clock starts at logic 0 and has its first 0-to-1 transition at \((P-W-\theta)\) time units after the start of simulation.

2. For phase \( \theta = P - W \), a 0-to-1 transition takes place at \( T = 0 \).

The default value for \( \theta \) is zero. In the example a phase of 50 units is specified for \( \text{CLOCK1}(1:5) \) and \( \text{CLOCK2} \). Since no phase specification is given for \( \text{CLOCK1}(6:10) \) and \( \text{CLOCK3} \), \( \theta = 0 \) is assumed for them. Waveforms for these clocks are shown below. Note that it is necessary to specify pulse width \( W \), if it is desired to specify phase \( \theta \).

During a simulation run, none of the parameters, \( P \), \( W \), and \( \theta \) can be respecified for a clock facility. These parameters remain effective in all subsequent runs until respecified.

\[ \text{CLOCK1 (1:5)} \]
\[ \text{CLOCK2} \]
\[ \text{CLOCK1 (6:10)} \]
\[ \text{CLOCK} \]

\[ P = 100, \ W = 30, \ \theta = 50 \]
\[ P = 100, \ W = 50, \ \theta = 0 \]
\[ P = 100, \ W = 30, \ \theta = 0 \]

waveforms
As mentioned earlier this declaration allows new facilities to be declared as simulation clocks. Since these clocks cannot affect the activity within the system itself, they are a source of periodic signals which can be used to control input, reinitialization, output, etc., during simulation. They can be used in realizing signals with complex waveforms that are needed to control various activities during simulation. Simulation clocks may also be used as sources of input signals to the networks being simulated.

Each facility \( f \) from clock list \( \ell_c \) is assigned parameters \( t \) from associated time list \( \ell_t \). Insufficient or excess data in time list \( \ell_t \) will result in a non-fatal error (see Sec. 4.4 for errors). In the case of insufficient data, default parameters are assigned to facilities remaining in \( \ell_c \).

4.2.4 <DElay> Declarations

This declaration provides a means for specifying delay time \( \Delta \) for delay facilities. Syntax for this declaration is very similar to that of the <CLock> declaration.

\[
\text{<Delay>} \quad \text{Body}
\]

\[
\text{Body} \quad [\ell/,]^n \ell [/]
\]

- **list**
  \[ \ell : \ell_d/\ell_t \]
- **Delay list**
  \[ \ell_d : \ell_t \text{ where} \]
- **Facility type**
  \[ \delta_f : \text{system delay, simulation delay} \]
- **Time list**
  \[ \ell_t : t[.,t]^n \]
- **Time spec**
  \[ t : [R^*]t \]
- **Delay time**
  \[ \Delta : n_1 \]

**Example:**
\[
\text{<Delay>} \text{ DELAY1(1:2), DELAY2, DELAY1(3:5)/2*100,5u.}
\]
DELAY1(1:2) and DELAY2 are each assigned a delay time of 100 units. DELAY1(3:5) is assigned a delay time of 50 units.

All the delay facilities are assumed to be **inertial delays**, i.e., an output signal(s) will assume a new value(s) Δ time units after its input signal prescribes that change, if and only if the input signal prescribes that new value for at least Δ consecutive time units. As an example of inertial delay assume that waveform A below serves as the input signal to both DELAY1(1) and DELAY1(3). Waveforms B and C represent the actual output of DELAY1(1) and DELAY(3) respectively.

![Waveform Diagram]

Delay time period A can not be respecified within a simulation run. Once specified, A remains effective in all subsequent simulation runs until respecified.

Like the <Clock> declaration, this declaration also allows a user to declare new delay facilities that may also be used for controlling various activities during simulation.
Every delay facility from \( \ell_d \) is assigned, in turn, delay times from the associated time list \( \ell_t \). Insufficient or excess data in \( \ell_t \) will result in a non-fatal error. In the case of insufficient data, the default delay time (4.2.5) is used for remaining facilities in \( \ell_d \).

4.2.5 Default Values for Clock Parameters and Delay Times

Before any simulation can be performed, it is necessary to assign clock parameters to every clock facility and delay time to every delay facility. Values specified through \(<\text{Clock}>\) and \(<\text{Delay}>\) declarations are used for specified facilities. For the remaining clock and delay facilities, default values are used. A default time period \( t_d \) is used in determining the default values.

1. Default clock parameters
   
   Default time period \( P = t_d \)
   
   Default pulse width \( W = \lfloor t_d / 2 \rfloor \)
   
   Default Phase \( \theta = 0 \)

2. Default delay time period \( = \lfloor t_d / 2 \rfloor \)

   At the start of a simulation job \( t_d \) is set to a value of 2. If any \(<\text{Clock}>\) or \(<\text{Delay}>\) declaration is encountered in the simulation deck, the value \( t_d \) is changed to

   \[ t_d = \min(P, 2A) \] where

   \( P \) is any clock period specified, if none \( P = 2 \), and \( A \) is any delay time specified, if none \( A = 1 \).

4.2.6 <Initialize> Declaration

This declaration provides a means for initializing the output value(s) of delays, registers, memories, element outputs, primary input signals, terminals and triggers with delays. Syntax for this declaration is as follows:
<INInitialize> Body

Body
: \[ \ell \rightarrow \ell \]

List
: \[ \ell \rightarrow \ell \]

Initialize List \( \ell_i \) : \( \ell_i \rightarrow \ell \)

Facility type \( \ell \) : system delays, simulation delays, registers, memories, element outputs, primary inputs, terminals, and triggers with delays.

\( \ell_d \) : Data list (see Sec. 4.2.2.)

Every facility \( \ell \) from \( \ell_i \) is initialized to a specified value obtained from the associated \( \ell_d \). Insufficient or excess data in \( \ell_d \) will result in a non-fatal error. If data in \( \ell_d \) is insufficient, remaining facilities from \( \ell_i \) are initialized to default values.

EXAMPLE: <IN> INPUT, MEM(0:1023)/B1011,1024*0/

INPUT (declared as register having width 4) is initialized to the binary value 1011 and the first 1024 locations of MEM are all initialized to 0.

Before any simulation can be performed during a run, it is necessary to define output values or initialize all the facilities. For all the facilities initialized through an <INInitialize> declaration(s), specified values are used. For remaining facilities initial values are determined as follows:

1. Delays, Registers, Element outputs, Primary inputs, Terminals, or Triggers with delays are all initialized to zero.

2. Memory locations are not initialized at all. They will have the same contents as at the termination of previous simulation run. For the first simulation run their contents are unpredictable.

3. Initial values for Terminal, Triggers, and Element inputs without delays are determined by using initial values for other facilities
and simulating the system at \( T = 0 \).

### 4.2.7 <REad> Declarations

This declaration provides a means for establishing input data values for various facilities. Syntax for this declaration is as follows:

\[
\text{<REad> Body}
\]

\[
\text{Body : } [\ell_1, \ldots, \ell_n]
\]

\[
\text{List } \ell : m, \ell_{1d}
\]

\[
\text{Mode } m : X, V, Z
\]

\[
\text{Triggered or Mode, } X : \delta \text{ where } \delta_0 = 1
\]

\[
\text{Periodic or Mode, } V : P[\theta]
\]

\[
\text{Period } P : n^P_{1, P_{\text{MAX}}}
\]

\[
\text{Phase } \theta : n_{0, P}
\]

\[
\text{Specific Time or Mode } Z : n
\]

\[
\text{Read List } \ell_n : \ell_0 \text{ where}
\]

\[
\text{Facility type } \delta_{\ell} \text{: registers, system delays, simulation delays, memory locations, element outputs, terminal or triggers or element inputs with delays}
\]

\[
\text{Data List } \ell_d : \text{ same as in <INitialize>}
\]

Example: \(<TR>\)

\[
\text{TR/EXINP+EXBIN1/ (see Sec. 4.2.15)}
\]

\[
\text{<CL> P/100(30)/}
\]

\[
\text{<RE> TR/INPUT/1,2,3,4,-5/}
\]

As shown in the syntax, the READ operation may be carried out in three different modes:

1. Mode X -- Triggered Mode

   In this mode a 0-to-1 transition of the triggering signal establishes a new set of input values, obtained sequentially from the associated data.
list $\ell_d^*$, for the facilities specified in the associated read list $\ell_n^*$.

At any simulation time input values are established before any other simulation activity except for updating of clocks and delay outputs. Hence, if the triggering signal itself is not a clock or delay facility, input values will be established at a time later than the actual 0-to-1 transition time of the triggering signal. In fact they are established at the next event time.

2. **Mode V -- Periodic Mode**

   This mode provides an easy means for establishing input values periodically. $P$ specifies the time period for performing the READ operation. The first READ operation is performed at $T = P$, the next at $T = 2P$, and so on. However, the first and all subsequent READ operations may be advanced relative to the beginning of simulation i.e., $T = 0$, by optionally specifying the phase $\theta$. Thus, in the case of $P = 100$, and $\theta = 30$, the first READ operation will be performed at $T = 70$ (advanced by 30), the next at $T = 170$, and so on. When $\theta = P$, the first READ operation is performed at $T = 0$. This is equivalent to initializing the associated facilities using an `<INitialize>` declaration.

   In all cases except for $P = 1$, an identical periodic READ operation can be obtained using a clock with period $P$, any valid pulse-width $W$ and appropriate phase $\theta$ as a triggering signal in mode X.

3. **Mode Z -- Specific Time Mode:**

   In this mode the READ operation is performed only once at the specified time.

   In Mode X and Mode V READ operations, data values are supplied in sets. The first set of values are used for the first READ operation,
and the next set is used for the second READ operation. These sets are not separated by any special delimiter. Instead they are grouped in the form of a single data list $\ell_d$. In Mode 2 only one set of data values are necessary.

4.2.8 <LOad> Declarations

This declaration provides a means for establishing the same input values repeatedly on specified facilities. Syntax for this declaration is as follows:

<LOad> Body

Body : same as in the <REad> except that the Load list $\ell_d$ is used in place of the Read list $\ell_r$.

Three modes of LOAD operation function in the same way as the three modes of READ operation. The only difference between LOAD and READ operations is the input data values used during successive operations. In the case of READ operations, a new set of data values is used for each successive operation. The LOAD operation uses the same data set repeatedly, requiring only one set of data values. This peculiar property of the LOAD operation provides an easy means for establishing identical conditions in the system at desired times. If the READ operation were used to achieve the same results, the same data set would have to be repeated for every occurrence of READ operation. The Mode 2 or specific time LOAD operation is identical in all respects to the Mode 2 READ operation.

The three modes available for READ and LOAD operations give a high degree of freedom in setting up data sets in an efficient manner. Each of these modes may be used more than once. More than one mode may be used in a simulation. All the data lists $\ell_d$ specified in <REad> and
<LOAD> declarations are transferred to an incore buffer (if necessary to a disk data file) and retrieved from there whenever needed. This facility of having multiple input streams for simulation is very helpful to the user.

More than one READ and/or LOAD operation may take place at the same simulation time. Simultaneous operations may attempt to establish input values on the same facilities. As long as they do not attempt to establish conflicting values, simulation will proceed, otherwise a fatal error condition results in an immediate termination of the current simulation run. A similar situation may arise with <INItialize> declarations. In this case remaining declarations for the simulation run are processed before terminating that simulation run. This fatal error condition may be avoided by setting Flag 12 (see Sec. 4.2.14).

The following order is used in performing various input operations during simulation:

1. Periodic REad
2. Specific Time READ
3. Periodic LOAD
4. Specific Time LOAD
5. Triggered READ
6. Triggered LOAD

If more than one operations of the same type and same mode take place at the same time, they are performed according to their order in the simulation deck. This is one case in which the order of declarations may be important.

Insufficient data to complete a READ or LOAD operation during simulation will result in an immediate termination of the run. This
provides a means to terminate a simulate run without using the <STOP>
(see Sec. 4.2.11) declaration.

4.2.9 <OUTPUT> Declarations

This declaration provides a means for printing the values of various
facilities at various instants during simulation. Syntax for this de-
claration is as follows:

```
<OUTPUT> Body

Body : [\ell,] \ell /\ell

List \ell : m/\ell_o

Mode m : X, V, Z

Triggered or Mode X : \delta where \delta_w = 1

Periodic or Mode Y : P[\theta]

Period P : n_1, P_{MAX}

Phase \theta : n_1, P

Specific Time or Mode Z : \eta

Output List \ell_o : \ell \delta, where \delta = \delta_z = \text{memory}
```

Like <READ> and <LOAD> declarations, this declaration has three
modes of operation. Values are printed when a specific output operation
takes place. It is important to note that in the case of triggered or
Mode X output, 0-to-1 transition of the triggering signal causes the
output values to be printed at the same time rather than the next event
time as in case of READ or LOAD operations. This is due to the fact that
output operations are performed after all other operations in a simulation
step. More than one <OUTPUT> declarations may be specified. Any
combination of the three <OUTPUT> modes may be used.

Values are normally printed in an octal format. They may be printed
in binary, octal, decimal or hexadecimal by setting the appropriate flag number from 5 to 9 (see Sec. 4.2.14). All values are printed in the same format.

Output formatting is done by the simulator with the objective of maximizing the total no. of facilities than can be printed. If one or more output operations occur at a simulation time only a single line of output is printed. The first entry in each line printed is the simulation time in decimal. Values for each facility specified in any output lists are printed in fixed columns. Facilities are allocated columns from left to right in the following way:

1. Triggered or Mode X OUTPUT lists
2. Periodic or Mode Y OUTPUT lists
3. Specific time or Mode Z OUTPUT lists

If more than one lists of a mode are specified, they are allocated columns in the order of their specification in the simulation deck. If output values for all specified facilities cannot be printed due to lack of room, excess facilities are ignored and a message listing them is printed. Output values for two neighboring facilities are always printed. Output values for two neighboring facilities are always separated by at least one blank column. A heading of the names of the facilities along with the subscript(s), if necessary, whose values appear below is printed on alternate pages of the simulation report. If a complete facility is included in \( \ell \), no subscripts are printed in the heading. When the value of a partial facility is to be printed, a subscript range is included in the heading. The name of a facility is normally printed in a horizontal format in the heading. A vertical format (in a column) is used if doing
so saves room on output line. A subscript range is indicated by two subscripts separated by a colon (:).

Whenever an output operation occurs, only the output values for facilities from the associated output list are printed. Other columns in the line are left blank. This tends to increase the readability of results. This feature of multiple output lists with each list having its own output control may be used to make simulation reports look more informative. If the output values for one group of facilities change less frequently than those of another group, both groups can be printed with different periods. Such an output will clearly illustrate the actual activity within the system.

4.2.10 <DUMP> Declarations

This declaration provides a means for dumping the contents of specified memory locations at various instants during simulation. Syntax for this declaration is given below:

```
<DUMP> Body
```

```
Body : same as for <OUTPUT> except $l_e$ : memory
```

A DUMP operation functions in a manner identical to the OUTPUT operation. The print format is different, however. First, values for each specified memory facility are printed separately. For each facility, the first line printed indicates the facility name, locations dumped and simulation time. Following this line a heading that separates addresses and contents of locations is printed. One or more lines of DUMP output follows. In each line the first entry represents the octal address of the first location in the line. The rest of the line contains the octal contents of the next eight locations. Various DUMP operations are carried out in the following order:
1. Triggered or Mode $X$ DUMPS
2. Periodic or Mode $Y$ DUMPS
3. Specific time or Mode $Z$ DUMPS

DUMP operations are performed before any OUTPUT operations within a simulation step.

4.2.11 <STOP> Declarations

This declaration provides a means for stopping or terminating a simulation run at a specified simulation time or on a 0-to-1 transition of a triggering signal. Syntax for this declaration is as follows:

\[
\begin{align*}
\text{<STOP> Body} \\
\text{Body : } m[,m]^n \\
\text{Mode} & \quad m : X, Z \\
\text{Triggered or Mode } X & \quad \delta \text{ where } k_{w} = 1 \\
\text{Specific Time or Mode } Z & \quad n
\end{align*}
\]

It is clear from the syntax that more than one triggering signal or specific time may be specified to stop the simulation. More than one <STOP> declarations may be specified. In any case the occurrence of a first stop event will cause the simulation for that run to be terminated. At a given simulation time stop events are simulated after all other events have been simulated. If no <STOP> declaration is supplied for the current simulation run stop events, if any, from a previous simulation run are used for the current run.

Insufficient data to complete a READ or LOAD operation will result in an immediate termination of the simulation run. This condition is described as "END-OF-FILE ON INPUT." If one is sure of EOF terminations, <STOP> declarations may be omitted altogether. Whenever simulation is
stopped or terminated a message describing the reason for termination
(a stop event or EOF on input) is printed and the simulator moves to the
next simulation run. At the end of all simulation runs an "END OF
SIMULATION" message is printed.

4.2.12 <LIst> Declarations

When a list of same facilities $\ell_6$ is used in a number of different
declarations, it is convenient to identify the entire list with a single
name. This name can then replace the list of facilities in all of the
declarations. This is achieved by using a <LIst> declaration. This
declaration provides a means for assigning a unique name to a list of
facilities. Syntax for this declaration is as follows:

\[
\langle \text{LIst} \rangle \quad \text{Body}
\]

\[
\text{Body} \quad : \quad [\ell/1]^n \ell[n/1]
\]

\[
\text{List} \quad \ell \quad : \quad L/\ell_6
\]

\[
\text{List Name} \quad L \quad : \quad \ell \quad \text{where } \ell_6 = 1
\]

A list-name can be included in the facility list $\ell_6$ for a declaration
only if the list of facilities identified by it can be directly used
there. It is also possible to use list-names in defining new list-names.
Nesting of list-names can be done up to any desired level. List recursion,
i.e., using a list-name in defining itself, is not permitted. Once de-
clared for a simulation run, list-name cannot be redefined in the same
run.

For large systems, use of list-names will result in reduction of data
structure storage space. List-names are most commonly used in <REad> and
<LOad> declarations since it is necessary to respecify these declarations
for each simulation run, if a <REad> or <LOad> declaration requires a
long facility list, it is very worthwhile to assign a list-name to these facilities and use the list-name in their place.

4.2.13 <Simulate> Declarations

As discussed earlier this declaration is used to separate different simulation runs in a simulation job. Syntax for this declaration is very simple:

<Simulate>

On encountering a <Simulate> declaration, simulation is performed for current run. If this simulation is terminated normally i.e., through a <Stop> command or EOF on input condition, processing for the next run is initiated. If the termination is abnormal, the entire simulation job is terminated. Declaration- and parameters effective during one run are carried over to the next run as described below. Modifications and additions are easily made with appropriate declarations.

1. Parameters for clock and delay facilities remain effective from one simulation run to the next; any parameter may be changed by using the appropriate declaration. New clocks and delays may also be declared.

2. Trigger expressions remain unchanged from run to run unless they are respecified. New trigger facilities may be declared for a simulation run.

3. <READ> or <LOAD> declarations do not carry from one run to the next. <READ> and <LOAD> declarations must be respecified for each run or replaced with new declarations.

4. <OUTPUT>, <DUMP>, and <STOP> declarations are carried from run to run. However, supplying one of these declaration with a specified mode (X, Y, or Z) will nullify all declarations from previous run
of that type and mode.

5. All flags are carried from run to run. Flags can be changed in any way by including a new <FLag> declaration.

6. Lists are carried from run to run. If it is necessary to redefine a list the new definition must be declared before the list is used directly or indirectly in any declaration for the current run.

4.2.14 <FLag> Declarations

This declaration provides a means for selecting various options for simulation runs by setting or resetting associated flags. A flag number is associated with each option. Syntax for this declaration is as follows:

<FLag> Body

Body : \( V_0, V_0 \)^n

Option Value \( V_0 : [\sim] F \)

Flag Number \( F : \{1,14\} \)

If the flag number \( F \) is not preceded by a complement sign (\( \sim \)), the associated option is set, otherwise it is reset. An option may be set or reset as many times as desired. The Flag table provides a description of the option controlled by each flag number and the default value for the option. As shown in the table flag numbers 2 thru 5 are used to select the radix for input data. This option applies only to data values not having any explicit radix specification (see Sec. 4.2.2). Data values having explicit radix specifications are interpreted accordingly. If more than one of these options is set, only the last set option is used, i.e., <FL> 2, 4 is equivalent to <FL> 4. Moreover, resetting any of these options brings the default radix specification to its default value of 8 (octal). Similar action is taken for output format selection flags 6 thru 9.
<table>
<thead>
<tr>
<th>Flag</th>
<th>Significance</th>
<th>Default</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Print source card images</td>
<td>Set</td>
</tr>
<tr>
<td>2</td>
<td>Binary data input</td>
<td>Reset</td>
</tr>
<tr>
<td>3</td>
<td>Octal data input</td>
<td>Set</td>
</tr>
<tr>
<td>4</td>
<td>Decimal data input</td>
<td>Reset</td>
</tr>
<tr>
<td>5</td>
<td>Hexadecimal data input</td>
<td>Reset</td>
</tr>
<tr>
<td>6</td>
<td>Binary output format</td>
<td>Reset</td>
</tr>
<tr>
<td>7</td>
<td>Octal output format</td>
<td>Set</td>
</tr>
<tr>
<td>8</td>
<td>Decimal output format</td>
<td>Reset</td>
</tr>
<tr>
<td>9</td>
<td>Hexadecimal output format</td>
<td>Reset</td>
</tr>
<tr>
<td>10</td>
<td>Use 1's complement notation</td>
<td>Reset</td>
</tr>
<tr>
<td>11</td>
<td>Write processed system to file</td>
<td>Reset</td>
</tr>
<tr>
<td>12</td>
<td>Do not abort on &quot;conflicting inputs&quot; error</td>
<td>Reset</td>
</tr>
<tr>
<td>13</td>
<td>Simulate comb. portion of the system with delay</td>
<td>Reset</td>
</tr>
<tr>
<td>14</td>
<td>Not used</td>
<td></td>
</tr>
</tbody>
</table>

4.2.15 <TRIGGER> Declarations

As discussed earlier, a triggering signal is used to control triggered or mode X READ, LOAD, OUTPUT, DUMP, and STOP operations. Any element of a declared facility, except a list-name, may be used as the triggering signal for these operations. Appropriate triggering signals to control the simulation may not be available in the DDL description of a system. The <TRIGGER> declaration provides a means for declaring new facilities that can be used as triggering signals to control simulation
without requiring that the DDL system description be modified. The syntax for this declaration is as follows:

\[ <\text{Trigger}> \quad \text{Body} \]

\[ \text{Body} : [t_E, t_E] \]

Trigger expression \( t_E : t'_0 / E \)

Trigger facility \( t'_0 : \delta \) where \( \delta_w = 1 \)

Expression \( E : \) an expression

Example: \( <\text{TR}> \quad \text{TR/EXINP+EXBIN1} \)

The expression \( E \) in the above syntax is a logical expression which can vary from simple sum-of-products form to the most complex and compound of forms. It defines the associated trigger facility \( t'_0 \). A trigger facility may be used in defining other trigger facilities. Looping or trigger facilities, i.e., using a trigger facility directly or indirectly to define itself, is not permitted, however. In the example trigger TR is defined as the logical OR of EXINP and EXBIN1. Both EXINP and EXBIN1 have been declared to be states of an automaton. The automaton waits in each of these states for data from an input device. The input device can be simulated using a triggered \(<\text{READ}>\) operation with TR as the triggering signal as shown in the example in Sec. 4.2.7.

A trigger facility cannot be redefined during a simulation run. The definition of a trigger facility remains effective in all subsequent runs until respecified. A trigger facility may be assigned a delay time \( \Delta \) using a \(<\text{DELAY}>\) declaration. Similarly a delay declared during simulation may also be defined using a \(<\text{TRIGGER}>\) declaration. For such facilities, both the delay time \( \Delta \) and the expression \( E \) remain effective in subsequent runs until respecified.
4.3 SIMULATION ALGORITHM

DDLSIM is a table-driven, event-oriented simulator. Time is treated as a discrete quantity and advanced from event time to event time where the following actions are considered by the simulator to be events:

1. Zero-to-one transitions of clocks. During these events data input signals to registers and memories are sampled, and next values of register and memory contents are computed and saved.

2. One-to-zero transitions of clocks. During these events register and memory contents are updated to new values.

3. Delay lines taking new values.

4. Simulator input, output and control events.

The simulator maintains a list of events to be executed in the future. Simulation is performed only at event-times. The simulation clock is always stepped from one event-time to the next event-time, no simulation being performed for the intermediate time interval. The absence of any event during these intermediate time intervals implies that no change is taking place in the system. For each event-time tests are performed to establish the need for simulation. Simulation is performed at event-time only if needed. A periodic event causes a future event to be scheduled.

Event time simulation makes the units used for time specification unimportant. Any arbitrary units can be used. The number of events simulated and not the number of time units elapsed determine the computer time consumed by a simulation run. Since the largest integer handled by the SEL 32 machine is \(2^{32}-1\), it is necessary to keep the simulation termination time within this limit. It is suggested that smaller
time periods be used for long simulation runs to avoid overflow of the simulation clock.

At a given event-time various events, if present, are processed in the following order:

1. Zero-to-one transitions of clocks
2. One-to-zero transitions of clocks
3. Change of output values for delays
4. Periodic or Mode Y READ operations
5. Specific time or Mode Z READ operations
6. Periodic or Mode Y LOAD operations
7. Specific time or Mode Z LOAD operations
8. Periodic or Mode Y DUMP operations
9. Specific time or Mode Z DUMP operations
10. Periodic or Mode Y OUTPUT operations
11. Specific time or MODE Z OUTPUT operations
12. Specific time STOP operation

After processing these events different simulations steps, if necessary, are performed in the following order:

1. Triggered or Mode A <REad> operations are completed
2. Triggered or Mode A <LOAD> operations are completed
3. If there were any new one-to-zero transitions of clocks declared in the DDL description or combinational clocks, i.e., signals generated using combinational logic and used as clocks for registers, output values for affected Registers or memory locations are changed to their new values.
4. If necessary, the combinational portion of the system is simulated. If any new one-to-zero transitions of combinational clocks are detected, Step 3 and 4 are repeated until no more one-to-zero transitions occur.

5. If any one-to-zero transitions of system clocks or combinational clocks were registered, new output values are computed and saved for affected registers and memory locations.

6. If necessary, delays are simulated to compute new future output values.

7. Triggered or Mode X DUMP operations are completed

8. Triggered or Mode X OUTPUT operations are completed

9. Triggered or Mode X STOP operations are completed

This procedure is repeated until the simulation is terminated.
4.4 ERRORS

DDLSIM performs very extensive error checking. On detection of an error, an error message is printed. Whenever possible an attempt is made to pin-point the error. Error messages are printed in one of the two formats discussed below.

1. Error messages which can be associated with a card in the simulation deck resulting from syntax errors are printed in the following format. The card containing the error is printed (if not already printed) with a vertical bar (\|) placed under the column containing the error or the column next to the item containing error. A dotted line starting from the column next to vertical bar (\|) and terminating with the error message on right end of the page is printed.

Example: <CL> CLOCK1(185) CLOCK (6:10)/2*100/

| .................. INVALID DELIMITER

Processing of the remainder of the declaration and the simulation deck is continued by skipping to an appropriate position in the declaration.

2. Errors which cannot be easily associated with a particular card in the simulation deck are printed in this format. The error message preceded by three asterisks, i.e., '***' is printed on the left end of the line. Error messages printed in this format normally contain an error description with associated parameters, i.e., facility name with appropriate subscripts, simulation time, etc., to help in locating the error. Some of the error messages require more than one line.

Example: ***RESPECIFICATION OF DATA FOR INPUT(1:5)

***AT TIME = 200
Errors are generally classified as fatal or non-fatal depending upon the nature, position and stage of simulation during which they occur. Fatal errors normally result in an immediate termination or abort of the simulation job. However, up to 10 fatal errors are allowed during the processing of the simulation deck for a simulation run. If any fatal errors were detected during the processing of the simulation deck, the entire simulation job is aborted. Whenever a simulation job is terminated due to fatal error(s) a message identifying the action is printed, i.e.,

***TO MANY FATAL ERRORS - SIMULATION TERMINATED.

Non-fatal errors do not cause the termination of the simulation job. In this regard they are warnings rather than errors.

DDLSIM performs complete syntax checking on the BEs and RTEs describing a digital system. Any errors detected during the processing of system description are treated as fatal errors. However, the simulation job is not terminated immediately. Since the errors detected during this stage cannot be easily associated with the DDL deck, they are printed using the second format described above. During the simulation stage complete error-checking is performed on the simulation process itself looking for errors such as:

1. invalid memory addressing,
2. instability in networks containing loops, and
3. attempts to input conflicting data on a facility.
5. EXAMPLES

This chapter provides some example DDL descriptions. The examples range from small synchronous circuits to a simple, but complete computer. These examples do not illustrate all the capabilities of DDL, but provide a good introduction to the user unfamiliar with DDL.

Example 1: A Serial Twos Complementer

The serial twos complementer uses the familiar copy/complement algorithm: starting from the least significant end of the number, copy the bits as they are till and including the first non-zero bit; complement the remaining bits till the most significant end. As an example,

<table>
<thead>
<tr>
<th>0 0 1 0 1 0</th>
<th>Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>1 1 0 1 0</td>
<td>Twos complement</td>
</tr>
<tr>
<td>complement</td>
<td>copy</td>
</tr>
</tbody>
</table>

This algorithm is implemented using a shift register and right circulating its contents while copying or complementing as required. The number of shifts is equivalent to the number of bits in the register. A flip-flop can be used to store the copy or complement state.

Figure 5-1(a) shows the description of the serial twos complementer in DDL. The content of the six bit register R is to be replaced by its twos complement. Register C (3 bits) counts the number of shifts. S is a state flip-flop to indicate the copy or complement state. T is a control flip-flop to indicate RUN/STOP state for the complementer. The complementer waits for SW to be ON, to start complementing. There is a clock P. An Operator ADD is described in lines 5-8. This is a 3 bit adder to increment the contents of the argument register by 1. The Automaton COMP has two states: a waiting state I, and a processing
state S1. Setting of SW is required for the transition to S1 state. In S1, the register R is circulated right 1 bit with the least significant bit copied or complemented, depending on the state of S being 0 or 1. If register C has reached a value of 5, the complementation is stopped by setting T to 0 and returning to state I. If C ≠ 5, COMP stays in S1 state and increments C. The FLag statement (line 13) sets the flags of the translator to provide the outputs at each of its six phases. Figure 5.1(b) shows these outputs. A detailed description of Figure 5.1(a) follows:

Line 1: The name of the system is COMPLEMENTER. Only the last 8 characters of this name are retained by DDLTRN. There is no period at the end of this line, since the system description is not complete yet.

Line 2: Register R has 6 bits numbered 1 through 6, left to right; C has 3 bits numbered 2 through 0; S and T are single bit registers. C counts the shifts; S is the copy (0)/COMPLEMENT (1) state flip-flop. T is a flip-flop indicating that the complementing process is underway. It is not really required, but included to illustrate some DDL features. Period terminates the REGister description.

Line 3: A Latch by name SW

Line 4: A single phase Clock (Time) P.

Line 5: A special Operator by name ADD. The output of the operator is a 3 bit number. The input is through the argument X (X is a formal parameter). No period to terminate, since the operator description is not complete yet.

Line 6: Declares the Terminal X to be of 3 bits and a new 3 bit register C. DDLTRN changes this name to C"1.
**Line 7:** Declares a new IDentifier for the concatenation of the last two bits of C and a 1.

**Line 8:** Declares the CARRY and SUM bits of an adder consisting of 3 half adders. C has the carry bits from each half adder, CC consists of carry bits from previous stages along with a 1 for the least significant bit. ADD consists of the SUM bits output. Note that ADD is the name of the operator, which is simply an ADD 1 circuit. The circuit implied (modelled) by lines 5–8 is:

![Circuit Diagram]

Note the periods at the end of line 8. The first terminates <BO> and the second terminates <OP>.

**Line 9:** Automaton COMP is controlled by the clock P. Since COMP is not subscripted (by parenthesis) it is assumed to be having only two states (1 bit). (If there are more than 2 states, then the number of bits required for state identification must be shown)

**Line 10:** State (Step) I with identification 0. Automaton COMP waits in I till SW is 1. When SW is 1, T is set to 1, C and S are set to 0, and a transition is made to state S 1 (all in parallel). The period terminates I.
Line 11: State S1 with the designation 1. waits for T to be 1. If S is 1, R is circulated right one bit with the bit R(6) complemented; otherwise R is simply circulated. S receives R(6) if S = 0. Also in this state, the value of C is checked to be equivalent to 5\(=101_2\). If C=5, T is set to 0 and a transition \(\rightarrow\) is made; if not, C is incremented and S1 state continues. \(\rightarrow\) periods at the end of line 12 terminate the

If \(\ldots\) THEN on \(C, S1, S1, RU\) and \(SY\) declarations respectively.

Line 13: Sets the FLags of DDLTRN to output results of each of the six passes.

Figure 5.1(c) shows the input commands for DDLSIM. FLags for DDLSIM are set for decimal data input (4) and binary output (6) in Line 1. SW is initialized to 1 in line 2. Two values are read into R one each time state I is reached (line 3). An output trigger OUTTR is declared to be ON at the falling edge of clock P (line 4). The values of COMP, R, S, C, T are to be OUTPUT when OUTTR is ON and that of R when in I state (line 5). The simulation is started with \(<SI>\) in line 6. Figure 5.1(d) shows the simulation output. The TIME starts with the raising edge of clock. Each edge is a time event. At time 0, all registers are zeroed and the circuit is in state I. At 1 SW is set to 1. At 2, R receives 5. 12 more time slots (6 clock pulses) are required for R to have its twos complement (time 14). At time 16, the new value for R (20) is received and its twos complement is ready at time 28. Since all the inputs are exhausted, the simulator stops at time 29.
FIGURE 5.1(a): DDLTRN INPUT
**OFFICIAL OENYA LANGUAGE TRANSLATOR**

**OFFICIAL FACILITIES IDENTIFIED**

**DECLARED FACILITIES**

```c
<SY> declarations
<FF> "(1:6)
    "(2:6)
    "(1:1)
    "(1:1)
<1> 5*(1:1)
<11> P(1:1)

<FF> "(1:3)
<1> 5*(1:3)
<11> P(1:1)
<FF> "(1:3)
<1> 5*(1:1)
<11> P(1:1)
<FF> "(1:3)
<1> 5*(1:1)
<11> P(1:1)
```

**DECLARED OPERATIONS**

```c
<SY> operations:
<FF> all:
<FF> (=* (c, 36) <st>):
<FF> (l e p):
<FF> (1:5):

1: S: 1 <1>, J <1>, S <6>, 1.1

S1: 1;
18: 1 <1> <(-1) (c), e (c) <(-1) 5; S <t (c), S <W (h) k (1:5)>
11 (c) + (1) * (1) |<c >; 1 <0>, =>1; c <&N h> f,
```

**Figure 5.1(b) DDTRN Output**
Djital O8er LANGUAGE TRANSLATOR

PASS---SYNAP REJECT

\[
\begin{align*}
&<01> \text{ FUTILITY: } I_{=}/\text{CONP}'0, \ S_{=}/\text{CONP}'11_{=1}, \ E_{1}=E_{1}(2.4)_{11_{=1}} \quad \text{AND} \ x_{=0}C_{1}(c:3)_{11_{=1}} \quad \\
&<00> \ (\text{rep}): \ P; \\
&\quad \text{I} \\
&\quad 1_{=1} \ I_{=101} \quad I_{=0}, \ S_{=0}, \ \Rightarrow 51_{=1} \\
&\quad 1_{=11} \ I \\
&\quad 1_{=1} \ P(1)_{=1} k(1)_{=1}, \ k(2:1)_{=1} x(1:15)_{=1}; \ S_{=1}H(5), \ R_{=1}H(6)_{=1} H(1:15)_{=1} \\
&1C(3_{=1} T L(1)_{=1} C(0))_{=1} I_{=0}, \ \Rightarrow 11_{=1} (\text{==40})_{=1}, \ C_{=1}C_{=1} \\
\end{align*}
\]

Figure 5.1(h) (Continued)
Figure 5.1(b): Continued
Example 2: The Serial Twos Complementer (variation 1)

Figure 5.2(a) illustrates another version of the twos complementer. Two operators are used. The six bit COM operator circulates register X. The bit fed into X(1) during circulation is either X(6) or \( \overline{X(6)} \) depending on the value of Y is 0 or 1, respectively. The CNTUP operator is the same as the ADD operator in example 1. This version just illustrates the use of operators. Figures 5.2(b) shows the DDLTRN output, 5.2(C) shows DDLSIM input and 5.2(d) shows the DDLSIM output.
FIGURE 5.2(b): DDLTRN OUTPUT
FIGURE 5.2(b) (Continued)
FIGURE 5.2(c) DDLSIM INPUT

```
1: <FL2>3,0
2: <FL2>1,0
3: <FL2>1,1,5
4: <FL2>1,1,14
5: <FL2>1,1,110
6: <FL2>1,1
```

FIGURE 5.2(d): DDLSIM OUTPUT

```
TIME | Z | X | S | C | E
---|---|---|---|---|---
0.0 | 0 | 0 | 0 | 0 | 0
2.0 | 1 | 0 | 0 | 1 | 0
4.0 | 1 | 0 | 0 | 1 | 1
6.0 | 1 | 1 | 0 | 0 | 1
8.0 | 1 | 0 | 1 | 0 | 1
10.0| 1 | 0 | 0 | 0 | 1
12.0| 1 | 1 | 0 | 1 | 1
14.0| 1 | 1 | 1 | 1 | 1
16.0| 1 | 0 | 1 | 0 | 1
18.0| 1 | 0 | 0 | 1 | 1
20.0| 1 | 0 | 0 | 1 | 1
22.0| 1 | 1 | 0 | 1 | 1
24.0| 1 | 1 | 0 | 0 | 1
26.0| 0 | 1 | 0 | 1 | 1
```
Example 3: Twos Complementer (variation 2)

Figure 5.3(a) shows a version of twos complementer description with the use of several Automata. Automaton CNT adds 1 to C, checks if it is 5 and sets DONE to 1 if C = 5. It is activated by COUNT. Automaton CMP is activated by CPT; performs the one bit circulation of R; sets COUNT to 1 to activate CNT. COMP is the controlling Automaton, activated by SW and in turn activates CMP in state S1 and waits for CCT to be 1 (for CNT completed) in S2. If DONE is 1, goes into wait state.

Figure 5.3(b) shows the DDLTRN output. Figure 5.3(c) and (d) show the DDLSIM input and output respectively. Note the effect of this version of description (Automata interaction) on simulation time.
Figure 5.3(a): DDLTRN Input
Figure 5.3(b) : DDLTRN Output
Figure 5.3(b) : Continued
Figure 5.3(b) : Continued
```c
C1 = f(1)*f(2) + f(3) + f(4)
C2 = f(5) + f(6) + f(7) + f(8) + f(9)
S1 = C1 + C2 + f(10) + f(11) + f(12) + f(13) + f(14)
S2 = C1 + C2 + f(10) + f(11) + f(12) + f(13) + f(14)
```
Figure 5.3(c): DDLSIM Input

Figure 5.3(d): DDLSIM Output
Example 4: MULTIPLIER [35]

A MULTIPLIER unit that calculates the product of two 8-bit numbers is described in Figure 5.4(a). A listing of the deck used for simulating the MULTIPLIER system along with the simulation report is given in Figure 5.4(b). The <Flag> declaration in the simulation deck specifies that all data-values specified without radix specification be interpreted in decimal (Flag 4), and that output values be printed in binary (Flag 6).

The control unit MPY of the system waits idly in state S1 until it receives a START command. A <Initialize> declaration is used to initialize the START signal to 1 and start the MULTIPLIER unit. On receiving the START command in state S1, the control unit proceeds to load the R register with the multiplicand obtained from the BUS and proceeds to state S2. In state S2 the B register is loaded with the multiplier obtained from the BUS. A triggered READ operation with state terminal S1 as the triggering signal is used to supply the BUS with the multiplicand. During simulation, whenever the control unit reaches state S1, the BUS is supplied with a new value of the multiplicand. The multiplier is supplied to the BUS in a similar manner with another triggered READ operation using state terminal S2 as the triggering signal. After loading the multiplicand and the multiplier, the control unit proceeds to state S3. In state S3 the multiplicand is added to the partial product, if the multiplier bit is a logic 1. The control proceeds to state S4 in any case. The A and B registers are shifted right together and the multiplication cycle counter MCOUNT is incremented. If the count has been completed, status line DONE is set to logic 1 and the control unit returns to its idle state S1. If not all bits of the multiplier have been tested, the control unit returns to state S3.
A triggering signal OUTTA defined using a <TRiggar> declaration is used in a triggered OUTPUT operation to control the printing of the values for MPY, HCOUNT, A, and B. These values are printed in binary on every trailing edge of the clock P signal. Another triggered OUTPUT operation using state terminal S1 as the triggering signal controls the printing of the values for the multiplicand, multiplier and the final product. Note that these values are printed only once, i.e., when the final product is available, during a given multiplication operation. The two output lists printed with different frequency make the simulation report more informative and readable. Since no <CLKock> declaration is included in the simulation deck, default values are used for P, W, and θ. Note that for a single simulation run a <SImulate> declaration is not required. Since an EOF condition is expected no explicit <STop> declaration is included in the simulation deck to terminate the simulation.
DIGITAL DESIGN LANGUAGE TRANSLATOR

1: <CO>MULTIPLIER
2: <SY>MULTIPLIER<11>.<HT=(112),(5),<10>,<COUNT>(5).
3: <NE>ZERO,UNL.
4: <TE>START,BUS(2),UNL.
5: <TE>SUM(8),COUNT(0),CSUM(3),CCLOUT(4).
6: <10>CIN=COUNT(2:3)(ZERO).
7: <10>CCIN=COUNT(2:3)(UNL).
8: <BU>COUNT=COUNT(1:8).*COUNT(1:8).SUM=COUNT(1:8).SUM.
9: CCOUNT=COUNT(1:8).CSUM=COUNT.CC1.
10: <AU>MPT(2):F:
11: <ST>S1(0);START:IN=<BUS>,COUNT=0,CREC=0,CNT=-1,->S2.
12: S2(1):A=<BUS,A=0,->S3.
13: S3(2):C(0):A=<CLOUT(1)SUM,->S4.
14: S4(3):A(1:8):B=<A(1:8),A(0)=0,->S5.
15: COUT=CSUM,1*COUNT+COUNT=1,->S1,->S5.
16: <TL>=5,4,3,0,0.

Figure 5.4(a) : DDLTRN Input
Figure 5.4(b) : DDLSIM Input
Figure 5.4(b) : JDLSIM Output
Example 5: MINICOMPUTER [52]

A description of a simple minicomputer is given in Figure 5.5. The details of the minicomputer are given in the Appendix. Lines 2-4 in Figure 5.5 describe the registers. Lines 5 declares a memory bus. Line 6 declares a START latch. Line 7 declares a four phase clock. Lines 8-11 declare a Increment (by 1) circuit. Lines 12-16 declare a 12 bit adder. Lines 18-19 are CPU initialization. Lines 20-23 show the FETCH cycle. Lines 24-25 show the DEFER state for Indirect Address calculation. Lines 26-27 show the OPCODE decoding. Lines 28-43 show the microoperations for each instruction.
Figure 5.5(a):

Minicomputer Description

1:  <SY> = [1]
2:  <PE> = AN(0:7), EN(u:11), PC(0:7), ACC (u:11), X(0:11).
3:  <HE> = IN(u:11) = H(3) # u#114b(r), u#0.
4:  <TE> = (256:12).
5:  <TE> = BUS(12).
6:  1<LS> = START.
7:  11<IP> = (u).
8:  11<LC> = IUP(4).
9:  <TE> = (H), C(r).
10:  <TE> = (C, 12) (1U1).
11:  <TE> = (C, 14) = xCC.
12:  <GU> = (12) = 12.
13:  <TE> = (12), C(12), Y(12), CUL(12).
14:  <L> = (12) = CUL(12) = 12 = 001.
15:  <E> = (12) = CUL(12) = 12.
16:  AUL = xCC.
18:  1<ST> = (0) = START, ALCC = u, HAN = u, NCC = u, X = u,
19:  1 = 1, -> FE.
20:  1<PE> = (1) = P(1) = AX = PC, 1<PE> = (2) = PC = LIUP(0:7).
21:  1<PE> = (u:11) = H(2) = u#114b(r), 1<PE> = (3) = I#0.
22:  1<PE> = (u:11) = H(2) = u#114b(r), 1<PE> = (3) = I#0.
23:  1<PE> = (11) = 1FF = 1FF = 1FF.
24:  1<PE> = (2) = (1) = AX = ALN, 1<PE> = (2) = u#114B(v#2), 1<PE> = (2) = u#114b(r).
25:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
26:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
27:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
28:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
29:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
30:  1<PE> = (3) = ALN = u#114B(v#2), 1<PE> = (4) = 1FF.
Figure 5.5(a) : (Continued)
6. CONCLUSIONS

DDLTRN and DDLSIM programs are currently being tested on SEL-32 Computer System. The output of the DDLTRN is suitable for logic generation. The output at PASS 6 and the facility table are now being analyzed to derive the algorithms for logic synthesis. With the logic synthesis programs complete, CADAT will be a true automatic design system.
APPENDIX

This is a preprint of the article to be published in the December 1979 issue of the "Proceedings of IEEE."
Computer Hardware Description Languages—A Tutorial

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Abstract—How to software designers use high level languages (HLL) so successfully in general is a mystery, except that they use hardware designers use hardware description languages (HDL) to describe the systems they are designing. HDL's are used extensively as a medium of precise yet concise description of digital hardware, they have found a variety of applications such as generating user manuals, teaching logic design, acting as an input medium for an interactive design system, etc. This tutorial paper presents HDL's as useful tools for simulation, the translation and simulation of HDL's are discussed along with the guidelines to select an HDL. The discussion for future work and an extensive bibliography are provided.

1. INTRODUCTION

A NY digital system can be described in the following at levels of complexity [1]-[4]:

1) algorithmic level which specifies only the algorithm used by the hardware for the problem solution.
2) Processor memory switch (PMS) level which describes the system in terms of processing units, memory components, peripherals, and switching networks.
3) Instructional level (programming level) where the instructions and their interpretation rules are specified.
4) Register transfer level where the registers are system elements and the data transfer between these registers are specified according to some rule.
5) Switching circuit level where the system structure consists of interconnected gates and flip-flops and the data is given by a set of boolean equations.
6) Circuit level where the gates and flip-flops are replaced by the circuit elements such as transistors, diodes, resistors, etc.

Logic diagrams and Boolean equations have been used as media of hardware description. The complexity of a given media increases rapidly as the system complexity increases and they are not convenient to suppress the details and still provide accurate descriptions as we move into the higher levels from the switching circuit level. Hardware description languages (HDL's) evolved as a solution. Although the use of computer oriented languages to describe digital system design can be traced back to Shannon's work on switching circuits in 1939. Aiken's work on switching theory in the 1940's, the logic diagrams at M.I.T. and NBS in the late 1940's and the flip-flop equations in the 1950's [5]. Ivozen's work [6] on a formal HDL probably initiated the contemporary interest in this area of research. An HDL is similar to any other high-level programming language (HLL) and provides a means of

1) precise yet concise description of the system;
2) convenient documentation to generate user manuals, service manuals, etc.;
3) input of the system description into a computer for simulation and design verification at various levels of detail;
4) software generation at the preprototype level, thus bridging the hardware/software development time gap;
5) incorporation of design changes and corresponding changes in documentation, efficiently;
6) designer/user (teacher/student) communication interface at the desired level of complexity.

HDL's are capable of describing the parallelism, nonrecursive nature, and timing issues in the hardware more naturally, and thus differ from the pure sequential nature of a general HLL. Some existing HLL's provide concurrency or simulated concurrency through their language elements, for example, PROLOG or PEPE [7]. An HDL can be classified as a procedural or a nonprocedural language [4]. Each statement in a nonprocedural HDL description would contain a label which describes the condition under which the activities described by the statements are to be performed. Thus the sequential ordering of the statements does not impose the ordering of the activities. In a procedural HLL description, the activities are performed following the sequential ordering of the statements.

HDL's are designed to describe both the structural and behavioral characteristics of a digital system. The fundamental properties of hardware systems and the art of hardware design process are the essential features of an HDL. For any HDL to be a useful tool, it has to possess the following properties:
1) It has to have a natural way of describing the parallelism, nonrecursive nature, and timing issues in digital hardware.
2) The structure and control parts of the hardware should be easily described and preferably the description of the two parts be separate (if such a division enhances the description), so that a user interested in the behavior of the system need not concern himself with the structure of the system. This division provides the flexibility to use hardware, software, or firmware for the control part, whichever is economical.
3) The language should serve as a medium at all levels of system description.
4) The design changes should easily be incorporated into the description and corresponding translation should be done preferably without a complete retranscription. This feature will be useful for the interactive environment. A translator translates the HDL description into an intermediate code from which the simulator and other programs can be driven (see Fig. 1). The intermediate code could be a set of Boolean and register transfer equations [31] or a computer executable code like polish strings [23].
5) The language should be easy to learn and remember; to accommodate the software-oriented hardware designer, although the hardware engineer cannot neglect the software aspects anymore, due to the impact of microprocessors. The design system should be portable, thus necessitating the translators and simulators of HDL be written in higher level languages.
6) Two approaches to system design are often proposed: the bottom-up approach where the elementary components are combined to form more complex ones and the top-down approach, where the system is decomposed into a collection of subsystems until the elementary components are reached. In practice, the designer may choose a combination of the two approaches. The structural detail at any design level varies from designer to designer. The HDL should allow the designer to control the amount of detail during each design phase.
7) The description of the large and medium size integrated circuit (LSI and MSI) modules as system components should be straightforward, so should be the inclusion of newer modules. If the system is partitioned by the designer to accommodate standard modules, this partitioning should be retained by the HDL translators and simulators.
Several HDL's have been reported [9]-76] since Iverson's proposal of such an HDL. Translators to convert the description into an intermediate executable code and simulators to execute this code have been written for some of these languages. No single HDL has met all the above characteristics. The tendency has been to invent a new HDL to suit a particular design environment, basically due to the difficulty in transporting the translators and simulators on to the new computing systems and extending them to accommodate the requirements of the new design environment. Table 1 [8] lists the implementation details of several HDL's reported. This list is by no means exhaustive.

Section II discusses the utility of HDL's in system design. A brief discussion of one popular HDL, the computer design language, is given in Section III along with two example descriptions. Two case studies are presented: one to select an HDL for an integrated circuit design environment (Section IV) and the other to show the utility of HDL's in concurrent hardware-software development (Section V). Future work required and current research topics are discussed in Section VI.

II. HDL's in System Design

Fig. 1 shows the utility of an HDL in a digital system design environment. The designer uses the HDL to describe his design. The description is translated into a computer executable data base, which serves as the source for various other operations. "The design can be refined by simulating at the description level (Loop 1), before proceeding to a more detailed simulation (Loop 2) at the logic level. The data-base also serves as a source for logic diagram generation, macrocode and test set generation. The physical construction of the system follows the simulation and refinement at the logic level.

Translation and simulation of HDL's have been well defined [9]-76]. Physical construction aspects have also been automated and are widely used in industry [77]. Test generation [78] and hardware compilation [12, 19] need further investigation. The variety of design methodologies, the artistic nature of the design process, and the ambiguity posed by the variety of components available make the hardware compilation a tedious task.

III. Computer Design Language

A hardware programming language (AHPL), computer design language (CDL), digital systems design language (DDL) and the instruction set processor (ISP) have been the most popular languages, partly due to their early introduction as general purpose HDL's. These languages were developed in university environments and are used in teaching digital logic design. New features are being added to these languages to make them more versatile. Well-tested translators and simulators are available for these languages (see Table 1 for references). Although several HDL's have been designed for an industrial use [59], [64], the design process being proprietary in nature, the use of HDL's is not widely reported [79].

This section provides a brief introduction to CDL. Example descriptions in CDL are provided. CDL was chosen over the others due to its simple structure and the author's familiarity with the language.

CDL was proposed originally by Chu [20]-[22]. A translator and simulator were written for a subset of this language [23]. Several modifications were made to this translator and simulator [26]-[29].

CDL describes the structural and functional parts of a digital system. The structural components like memory, registers, clocks, switches, etc. are declared explicitly at the beginning of the description. The functional behavior of the element is described by the community used operators and user defined operators. Valid data paths are declared implicitly whenever there is a data transfer. Both parallel and sequential operations are allowed. Synchronous operations require a conditional test for an appropriate signal. The language is easy to understand and is highly readable.

All the variables in a CDL description are global. The system description can be only at one level, and there is no subroutine facility in CDL, thus making it unsuitable for describing hardware in a modular fashion. It is not possible to include special hardware components like integrated circuits (IC's) in a description. However, its simplicity of structure and its portability resulting from the FORTRAN implementation, have made CDL a popular language. The description of CDL syntax and semantics is accepted by the present version of translator and simulator [29] as given below. Table II shows the standard operations in CDL. Facilities are declared at the beginning of the system description with declaration statements of the format

```
DEVICE list
where DEVICE can be a REGISTER, SUBREGISTER, MEMORY, DECODER, SWITCH, TERMINAL, BUS, BLOCK, and CLOCK. Some example declarations are shown below
```

```
REGISTER(A0-23, A, F16-1)
SUBREGISTER, P:OP=F(3-1), F(0)-F(6-4)
MEMORY, M@R=10(77.0-10)
DECODER, L0=15#G(2-5)
CLOCK, P(2)
SWITCH, STRT (OFF, ON)
TERMINAL, B0:1=0, B0:2=1
D@S+B
BUS, Z10=(-1)
BLOCK, SERCOM (A#4(1)-4(5-2))
```

A DO/SERCOM statement is used to invoke the set of statements declared by BLOCK, SERCOM

```
A unconditional microstatement has the form
```
```
Variable = Expression
```
```
A conditional microstatement has the form
```
```
IF (expression) THEN (microwantestant)
```
```
ELSE (microwantestant)
```
```
Examples IF (A AND A) THEN (A=A)
```
```
ELSE (A#A)
```
```
Conditional statements may be nested to any number
```
```
A labeled statement has the format
```
```
/lable/microwantestant
```
```
where
```
```
label = expression=clock
```
```
Example: X@10=F/A#B, B@4
```
```
Special operators can be established by the user through a separate subprogram. The format is
```
```
*OPERATOR Parameter Name
```
```
CDL declaration statements. RETURN
```
```
END
```
A count operator is defined below:

```pascal
OPERATOR, X: (X)->COUNT
IF (0=THEN (X+1)<1)
ELSE IF (0=THEN (X+2)<1)
ELSE (0=THEN (X+1)<1) RETURN
END
```

Examples: 4 A CNTUP, C 4ADD 8

The CDL TRANSLATOR performs a syntax check of the description and translates it into a set of tables and a polish parsing program.

The CDL SIMULATOR executes the output of the translator and can accept simulation parameters through the following command set:

- **LOAD**
  - Used to initialize registers and memory.

- **OUTPUT**
  - Provides a hexadecimal printout of the specified register and memory contents and switch positions at the desired clock or label.

- **SWITCH**
  - Enables setting switch positions.

- **RESET**
  - Restores the earlier settings of the simulation parameters.

- **SIMULATE**
  - Provides the start and stop conditions for simulation.

CDL can be used to describe simple to very complex digital systems. Two example descriptions are provided below to illustrate this feature.

**Example 1: A Serial Two's Complementer**

A circuit to replace the contents of a 6-bit register R by its two's complement will be described. The complementation is done by the well-known Copy/Complement algorithm (starting from the least significant bit of R, copy the bits as they are till the first nonzero bit, complement the bits after the first nonzero bit, till the most significant end of the register).

**Example 2: A Minicomputer**

Fig. 5 shows the structural details. Instruction set and the CDL description of a minicomputer [52]. The minicomputer has a 256 word 12-bit memory, with an 8-bit memory address register (MAR) and a 12-bit memory buffer register (MBR). There is an 8-bit program counter (PC) and an accumulator (ACC) of 12 bits. The arithmetic/logic unit (ALU) receives the operands from MBR and a 12-bit J register, and pass the results on to the 12-bit BUS. The instruction format is of a 3-bit operation code, an indirect address flag bit, and 8 address bits. The program-end description is provided by the Statements 1-3 of Fig. 5(b). The BUS's layout not explicitly described to retain the high level description nature. Fig. 5(c) shows the details of the instruction set. Statement 4 in Fig. 5(b) describes a START switch, a RUN switch to indicate the RUN/STOP state, and a th. state switch for indicating instruction fetch (F) and execution (E) phases. Statements 5 and 6 provide the instruction decoding details. There is a 4-phase clock cycle (Statements 7) which activates the synchronous control unit. Each clock cycle consists of 4 major cycles. The comments in the CDL description identify the Fetch Cycle, Decode Cycle, and the Execution cycle for each instruction. Fig. 5(d) shows a program to add the four numbers in memory locations 0-3 and place the sum in location 7. The program will be located in memory locations 10-16. Location 4 is initialized to -3 and incremented by 1 each time through the loop. Fig. 5(e) shows the program in assembly, binary, and decimal forms. Fig. 5(f) shows the memory map just before the execution of the program. This memory map is simulated by the LOAD command for the CDL simulator (Statements 43-45) in Fig. 5(b). The program is executed and the C and S are cleared. and the program counter a at to 10 (Statement 46) in Fig. 5(b). The switch is turned ON (Statement 42) and the simulator is requested for 200 label cycles (Statement 47), outputting several register contents (Statement 41) at each label cycle. The simulator results are similar to the two's complementer example and are not shown for the sake of brevity. It is evident that the CDL description of the minicomputer is concise and more precise than any natural language description could be.

IV. Selection of HDL

Due to the large number of HDL's proposed, the selection of an HDL for a particular design environment becomes a non-trivial task. Although the structure of the language, the operators available, the capabilities of the language to describe the design in a logical manner are important considerations, the implementation issues seem to override them. One such selection process is described here along with the system description.

Fig. 6 shows the details of the computer aided design and test (CADAT) system of the NASA Marshall Space Flight Center [80]. The developer inputs the details of the IC to CADAT as a set of standard cells and their interconnections. The standard cell selection is done manually from a standard cell library. That description is at the logic diagram level, detailed logic simulations and refinements are carried out on the design. The final design is input to the automatic test-vector generation and placement and routing programs. The IC mask pattern generation is done interactively and a mask analysis and performance simulations are done before fabricating the mask. The last two steps in the IC fabrication are the wafer processing and the final testing.

At present, the generation of logic diagrams and choosing the standard cells from the cell library for the design are done manually. Integration of a high-level design language would help the designer to simulate his design and refine it at a high level before entering his design into the current system. That requires an HDL with a simulator and a logic synthesizer (hardware compiler) that generates the logic netlist required of...
the CADAT System. The drawback imlementation and
testing of a complex large scale integrated circuit (LSIC) de-
ign is not feasible since it cannot be properly breadboarded
with anything but the LSIC itself. With an HDL, this bread-
boarding process can be substituted with a computer simul-
atation of the LSIC design, thus minimizing the design changes
and, hence, the cost of mask fabrication and wafer processing.
The following five criteria were used in selecting a suitable
language [11]:

1) Activity
2) Level of description
3) Software availability and portability
4) Ease of logic generation
5) Modularity.

1) Activity: It is essential to choose a language which is be-
ing used elsewhere to receive the benefits of the extensions to
the language. Most of the HDL’s proposed do not have a trans-
atc translator and simulator that is up-to-date and fairly versatile,
though the language itself is versatile. The process of improv-
ing the HDL software and capabilities would be aided by the
active interest of the other groups in the language.
2) Level of Description: The selected HDL should accom-
mmodate a description at the register transfer level and/or below
to facilitate the logic generation. A higher than register trans-
fer level description may not be needed for the IC design en-
vironment of CADAT.
3) Software Availability and Portability: The development
of a HDL is incomplete without a simulator and a translator
written for it, since this software development process refines
the language structure. The software should be portable to
accommodate the general portability of the CADAT software.
4) Ease of Logic Generation: Any HDL translator oriented
towards providing information for a simulator collects and re-
arranges the combinatorial logic and register transfer. This
intermediate translated output should be amenable to logic
generation.

Modularity: The HDL description should be modular enough
to reflect the modularity of the hardware, to enable easier
understanding and modular design verification.

A comparison of the four prominent HDL’s with respect to
the above criteria is shown in Table III. [11]. Although versa-
tile, does not lend itself to the logic generation level very well
CDL is suitable for microprogram generation. The nonmodular
description feature of CDL and the difficulty in using the
polish string output of the translator to generate logic diagram
description makes it unsuitable for the CADAT system
environment. AHPL and DDL were the strong contenders.
Both have a fairly portable software package and are suitable
for the level of description needed for CADAT. The modu-
larity is brought about by the subroutines feature in AHPL
whereas the block structure of DDL is closer to the hardware
modularity. From a traditional hardware designer’s point of
view, programming in either language is equally difficult. Al-
though a hardware compiler is available for AHPL [12], its
SNOBOL implementation makes newer implementation issues
for CADAT, which is predominantly in FORTRAN. The DDL
translator provides a set of Boolean equations and register
transfer expressions which can be used for hardware com-
plementation [9], [99] though not very easily. The block struc-
ture and the software of DDL made it a better choice over
AHPL for the CADAT system.

Note that the selection of the HDL is oriented more towards
the implementation issues, rather than a rigorous analysis of
the capabilities and the characteristics of the HDL itself, such as
the structure of the language, operators available, ease of un-
derstanding, etc. Such a rigorous analysis, although valuable, will
not aid in the selection of this language since the implemen-
tation issues override the other characteristics. Also, the selec-
tion criteria ignored the possibility of developing a new
language to exactly fit the CADAT environment. The activity
criteria also eliminated several other HDL’s like LEC [59] and
SDL [72] from consideration.

V. CONCURRENT HARDWARE AND SOFTWARE
DEVELOPMENT

The use of HDL’s in hardware development is obvious. The
recent advances in IC technology have tremendously increased
the speed of new systems announcements. But the software
development for the new system has not caught that pace.
With the ability of the HDL to describe and simulate the hard-
ware accurately, it is possible to develop the software for the
digital system concurrently to bridge the software–hardware
development gap. This section describes an experiment to
measure the performance of CEI in software development
[26].

A multiprocessor system, consisting of a Digital Equipment
Corporation PDP-8 Minicomputer and an INTEL 8080 Micro-
processor was used. The two processors were simulated indi-

dvually, followed by the simulation of the shared memory and
the input device for the system. The input device is an
on-line inspection station which interrupts the 8080 after each
part is examined to enter the measurements of the part into a
eight-word 8-bit memory. Intel 8080 handles the bookkeeping
of these measurements for use by PDP-8. Several programs
were written both for 8080 and PDP-8. The programs on
PDP-8 accept the measurement from 8080, determine if they
are within specifications, and transmit the condition of the
part to 8080. The 8080 programs handle this interrupt and
keep a record of the number of parts inspected and their con-
dition. The programs were written in assembly language of the
particular processor and were stored in the shared memory in
the machine language form. The details of the simulations can
be found in [27].

An important consideration in developing programs is the
assembly time required by the host processor running the CDL
simulations of PDP-8 and interface with the 8080. Table IV shows the CPU
time required for typical programs on an IBM 370/155.

Clearly, the cost of such simulations is prohibitive. However,
assume that the cross assemblers are available on the host
machine, developing an application program using CDL simul-
ation would not be very expensive, since these programs will
usually be shorter than an assembler or a compiler. A related
issue would be the performance comparison of such simula-
tions using both high level languages for the description of the hard-
ware, rather than an HDL. Much of the overhead of the HDL
translator/simulator software could be reduced by using an
HLL for describing and simulating the particular hardware. A
comparison of each HLL versus HDL descriptions and their
run times is needed.

VI. FUTURE WORK

Although the suitability of an HDL for hardware system
derivation is well recognized, the HDL’s are not used ex-
tensively, partly due to the variety of structures and notations
used in these HDL’s, making them harder to understand. Many
structures found in HDL’s are sample for a software pro-
fessional to understand and use. But a hardware designer not
familiar with programming finds them hard to use. This
problem will be partially solved by the popularity of the micro-
processors as design elements, requiring the hardware designer
to understand software.
The differences in notations and structures used by HDL’s
make it difficult to borrow a language developed elsewhere.
This difficulty is augmented by the nonstandard design met-

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The HDLs must be designed, described, and implemented. The following guidelines are used in designing HDLs:

1) CONLAN must support design, description, and simulation of at least the following classes of systems: gate networks, register networks, processors, memories, processor systems.

2) Any system may be described via either
   a) a network structure description or
   b) a behavior description.

3) CONLAN is a service
   a) computer architects and logic designers for purposes of trade-off exploitation and optimization, design verification, and design documentation
   b) systems, micro- and application programmers
   c) electronic system production engineers.

4) CONLAN syntax and semantics must support
   a) well-defined descriptions
   b) machine parsing, interpretation, and simulation with error detection
   c) comprehension of complex system structure and function
   d) division of design efforts
   e) control over the level of abstraction at which subsystems are described

5) CONLAN is to be evaluated in terms of benchmarks such as standard function declarations, time operator declarations, IC descriptions (including macroprocessor), and design descriptions (including a multicomputer system).

The basic aim of CONLAN is to provide a versatile uniform base language with the capabilities of augmenting the basic syntax with the specific constructs with their own semantic interpretation, as required by the environment.

The efficiency of the HDL software depends on the efficient use of the host computer on which it was developed. Hence, the software tends to be machine dependent, making it fairly nonportable. Although the efficiency suffer [28] if the software is made portable, a well-documented software package (along with a good discussion of the algorithms used) is a necessity. Several other areas of investigations could be identified.

Procedures to analyze the HDL descriptions of digital systems [44] are to be developed in order to avoid simulation or at least minimum simulation costs. HDLs are to be designed to use applications to accommodate easier description of LSI circuits and macroprocessors [83], [84]. Simplicity of HDLs as languages for macroprocessor software development [85] and architecture comparison needs investigations.

Comparison studies of HLL and HDL with respect to ease of programming, ease of understanding, description length, simulation cost and efficiency are required.

Logic synthesis from the HDL description is not well developed [12], [39], [79], [86]. Decomposition of the digital system to accommodate the LSI and MSI components and retaining this decomposition till the final stage in the design are of paramount importance. The ability of the procedure to search through a library of available IC's and the capabilities to accommodate newer modules is necessary.

The availability of macroprocessor processors has increased the popularity of distributed processing systems. The HDLs have traditionally been designed for a single processor environment and lack the facilities to describe the interprocessor communication. Such ability will make the HDL more attractive [87].

The capability of an HDL for a particular environment depends on its capabilities to accommodate the operations in the environment. Since the majority of HDLs are designed for a particular environment, they tend to be less suitable for other environments. For example, a HDL developed with a goal of efficient high-level description and simulation would hardly suit a logic synthesis environment. A classification of available HDLs according to their underlying models for behavior is needed.

VII. CONCLUSIONS

The capabilities of HDLs were discussed. A brief introduction to such language (CDL) along with example descriptions were given. Case studies for selection of an HDL and the use of HDL in hardware/software development were given. The areas for further investigations were identified.

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<td>CDC-1601</td>
<td>ALGOL</td>
</tr>
<tr>
<td>RTE II</td>
<td>69-70</td>
<td>COL RATS II</td>
<td>Sunnion 400/151</td>
<td>FORTRAN</td>
</tr>
<tr>
<td>RTE III</td>
<td>71-72</td>
<td>RTE-4</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>RIS I</td>
<td>73-74</td>
<td>COL RATS II</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>RIS II</td>
<td>75</td>
<td>RTL</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>SCL</td>
<td>76</td>
<td>APL</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>SDL</td>
<td>77</td>
<td>APL</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>SMT</td>
<td>78</td>
<td>APL</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>SMI</td>
<td>79</td>
<td>APL</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>VDL</td>
<td>80-81</td>
<td>–</td>
<td>–</td>
<td>–</td>
</tr>
<tr>
<td>V-PMS</td>
<td>82</td>
<td>–</td>
<td>–</td>
<td>–</td>
</tr>
</tbody>
</table>

Notes: 1) “-” indicates that the detail is either not available or not known.
2) No accuracy is claimed for the contents of the table.
### Table II

<table>
<thead>
<tr>
<th>Error</th>
<th>Example</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>LE</td>
<td>A.EQ.B</td>
<td>Expression is 1 ( (A) = (B) )</td>
</tr>
<tr>
<td>GE</td>
<td>A.NLE.B</td>
<td>Expression is 1 ( (A) \geq (B) )</td>
</tr>
<tr>
<td>GT</td>
<td>A.GT.B</td>
<td>Expression is 1 ( (A) &gt; (B) )</td>
</tr>
<tr>
<td>LT</td>
<td>A.LT.B</td>
<td>Expression is 1 ( (A) &lt; (B) )</td>
</tr>
<tr>
<td>LE</td>
<td>A.LE.B</td>
<td>Expression is 1 ( (A) \leq (B) )</td>
</tr>
<tr>
<td>AND</td>
<td>A.AND.B.A = B</td>
<td>Performed logical AND bit by bit</td>
</tr>
<tr>
<td>OR</td>
<td>A.OR.B.A = B</td>
<td>Performed logical inclusive OR bit by bit</td>
</tr>
<tr>
<td>XNOR</td>
<td>A.XNOR.B.A = B</td>
<td>Performed logical exclusive NOR bit by bit</td>
</tr>
<tr>
<td>ADD</td>
<td>A.ADD.B</td>
<td>Binary sum of (A) and (B) or ( (A) + (B) )</td>
</tr>
<tr>
<td>SUB</td>
<td>A.SUB.B</td>
<td>Binary difference of (A) and (B) or ( (A) - (B) )</td>
</tr>
<tr>
<td>CNTUP</td>
<td>A.CNTUP</td>
<td>Increment (A) by 1 or ( (A) = (A) + 1 )</td>
</tr>
<tr>
<td>CNTDN</td>
<td>A.CNTDN</td>
<td>Decrements (A) by 1 or ( (A) = (A) - 1 )</td>
</tr>
<tr>
<td>SHR</td>
<td>A.SHR</td>
<td>Shifter (A) right one bit position, output 0 at left</td>
</tr>
<tr>
<td>SRL</td>
<td>A.SRL</td>
<td>Shifter (A) left one bit position, output 0 at right</td>
</tr>
<tr>
<td>SHR</td>
<td>A.CIR</td>
<td>Circular (closed) right shift of (A) 1 bit</td>
</tr>
<tr>
<td>CIL</td>
<td>A.CIL</td>
<td>Circular (closed) left shift of (A) 1 bit</td>
</tr>
<tr>
<td>SHR</td>
<td>A.ASHRA</td>
<td>Arithmetic right shift of (A) 1 bit, no change in left most bit (zero bit)</td>
</tr>
</tbody>
</table>

* \( A \neq B \) Consists of 4 registers replaced by contents of \( B \)

### Table III

<table>
<thead>
<tr>
<th>ISP</th>
<th>CDL</th>
<th>AMPL</th>
<th>DDL</th>
</tr>
</thead>
<tbody>
<tr>
<td>1) Software translator</td>
<td>PDP-10 BILLES</td>
<td>Main-Fortran/Assembly</td>
<td>CDC 6400 Fortran</td>
</tr>
<tr>
<td>emulator</td>
<td>PDP-10 BILLES</td>
<td>Main-Fortran/Assembly</td>
<td>CDC 6400 Assembler</td>
</tr>
<tr>
<td>hardware compiler</td>
<td>no</td>
<td>no</td>
<td>CDC 6400 Assembler</td>
</tr>
<tr>
<td>memory</td>
<td>no</td>
<td>no</td>
<td>failed</td>
</tr>
<tr>
<td>2) Level of Description</td>
<td>instruction set level</td>
<td>register transfer level</td>
<td>register transfer and below</td>
</tr>
<tr>
<td>3) Modularity</td>
<td>yes</td>
<td>no</td>
<td>yes</td>
</tr>
<tr>
<td>4) Logic Connection</td>
<td>no</td>
<td>not very well</td>
<td>yes</td>
</tr>
<tr>
<td>5) Programming Ease</td>
<td>difficult</td>
<td>easy</td>
<td>difficult</td>
</tr>
</tbody>
</table>

### Table IV

<table>
<thead>
<tr>
<th>CPU</th>
<th>IBM 370 155 CPU Times for CDL Simulation</th>
</tr>
</thead>
<tbody>
<tr>
<td>ISP</td>
<td>INTEL 8080</td>
</tr>
<tr>
<td>1) Average time to simulate an instruction</td>
<td>0.22 s</td>
</tr>
<tr>
<td>2) Number of instructions to assemble an instruction</td>
<td>0.8 k</td>
</tr>
<tr>
<td>3) CPU time for assembling a 20 instruction program</td>
<td>96 000 s</td>
</tr>
</tbody>
</table>

*Note: \( x = 1024 \)*
Figure 1: Digital System Design Automation Process
Figure 2(a) : Serial Twos Complementer Hardware Structure
**Simulation**

**Translate**

1. **STORNGL***
2. **REGISER((1=0))
3. **SWITCH(SW(0/1))
4. **CUTROL***
5. **REGISER(C=0)
6. **CLOCK**
7. **CUTROLSSUR***
8. **/5a((u1)/1{l}0=0/52=)
9. **/1P/IF(S(b)=S)(1)=(S=K(u), N=K(u)-K(1-8)) ELSE
   (K(R(u))=K(1=3))-IF (C.E=E) THEN(1=C)
   ELSE(C=C.CUTUP)
10. **CUTROL***

**Simulation**

11. **SOUTP0**
12. **SWITCH**
13. **CUTROL**
14. **=K**
15. **=3**

**Figure 2(b) : CDL Description**
Figure 3: Controller for the Two's Complementer
OUTPUT OF SIMULATION - OCTAL

SWITCH TRANSITION AT LABEL CYCLE 1
Sw  ->  ON

R = 05   C = 0   S = 0   T = 1
LABEL CYCLE 1   TRUE LABELS   CLOCK TIME 1

R = 42   C = 1   S = 1   T = 1
LABEL CYCLE 2   TRUE LABELS   CLOCK TIME 2

R = 61   C = 2   S = 1   T = 1
LABEL CYCLE 3   TRUE LABELS   CLOCK TIME 3

R = 30   C = 0   S = 1   T = 1
LABEL CYCLE 4   TRUE LABELS   CLOCK TIME 4

R = 54   C = 4   S = 1   T = 1
LABEL CYCLE 5   TRUE LABELS   CLOCK TIME 5

R = 00   C = 0   S = 1   T = 1
LABEL CYCLE 6   TRUE LABELS   CLOCK TIME 6

R = 70   C = 5   S = 1   T = 0

SIMULATION ENDS AFTER 6 REPETITIONS
FINAL LABEL CYCLE IS:

*RESET   CYCLE + CLOCK

*LOAD   R = 21
*SW  300

Figure 4(a) : Two Complementer Simulation Results for R = (05) _8
OUTPUT OF SIMULATION - OJTH

SWITCH TRANSITION AT LABEL CYCLE 1

5w -> 0w:

<table>
<thead>
<tr>
<th>R</th>
<th>C</th>
<th>S</th>
<th>I</th>
</tr>
</thead>
<tbody>
<tr>
<td>21</td>
<td>0</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>50</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>64</td>
<td>2</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>72</td>
<td>3</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>75</td>
<td>4</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>36</td>
<td>5</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>50</td>
<td>6</td>
<td>1</td>
<td>0</td>
</tr>
</tbody>
</table>

SIMULATION ENDS AFTER 9 REPLICATIONS
FINAL LABEL CYCLE IS:

Figure 4(b) : Twos Complementer Simulation Results for R = (21)g
Figure 5(a) : Minicomputer Bus Structure
Figure 5(b): CDL Description
**$**

```c
23 /RUN(U1) *STATE(L) *P(L) *STATE(R1) (RUN)
24 /RUN(U1) *STATE(L) *P(L) *STATE(R1) (RUN)
25 /RUN(U1) *STATE(L) *P(L) *STATE(R1) (RUN)
26 /RUN(U1) *STATE(L) *P(L) *STATE(R1) (RUN)

Then

(P = PC + UNIT) * STATE = F

**UCA CALCULATION

```

27 /RUN(U1) *STATE(L) *P(U) *UCA/UNIT = F
28 /RUN(U1) *STATE(L) *P(U) *UCA/UNIT = F
29 /RUN(U1) *STATE(L) *P(U) *UCA/UNIT = F
30 /RUN(U1) *STATE(L) *P(U) *UCA/UNIT = F

**USR CALCULATION

31 /RUN(U1) *STATE(L) *P(U) *USR/UNIT = F
32 /RUN(U1) *STATE(L) *P(U) *USR/UNIT = F
33 /RUN(U1) *STATE(L) *P(U) *USR/UNIT = F
34 /RUN(U1) *STATE(L) *P(U) *USR/UNIT = F

**RETURN CALCULATION

35 /RUN(U1) *STATE(L) *P(U) *RET/UNIT = F
36 /RUN(U1) *STATE(L) *P(U) *RET/UNIT = F
37 /RUN(U1) *STATE(L) *P(U) *RET/UNIT = F

**JUMP CALCULATION

38 /RUN(U1) *STATE(L) *P(U) *JUMP/UNIT = F
39 /RUN(U1) *STATE(L) *P(U) *JUMP/UNIT = F

END

SIMULATE

**$**

```

**$**

```c
41 *OUTPUT
42 *IN
43 *IN
44 *IN
45 *IN
46 *IN
47 *IN

Figure 5(b) (Continued)
```
<table>
<thead>
<tr>
<th>Operation Code</th>
<th>Mnemonic</th>
<th>Operation Code</th>
<th>Mnemonic</th>
<th>Comments</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>AND</td>
<td>(ACC) * (Mem) - ACC</td>
<td>AND Memory</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>TAD</td>
<td>(ACC) + (Mem) - ACC</td>
<td>ADD</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>ISZ</td>
<td>Increment memory and skip next instruction, if zero.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>DCA</td>
<td>Deposit and clear ACC.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>JSR</td>
<td>Jump to Subroutine, (PC) - MP(0)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>JMP</td>
<td>Jump</td>
<td></td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>RET</td>
<td>Return</td>
<td></td>
<td></td>
</tr>
<tr>
<td>7</td>
<td>HLT</td>
<td>Halt</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**NOTE:** ( ) indicates "Contents of"

---

*Figure 5(c) : Instruction Set*
## PROGRAM

<table>
<thead>
<tr>
<th>Memory Location</th>
<th>Assembly</th>
<th>Binary</th>
<th>Decimal</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td>AND 5</td>
<td>000 0 00000101</td>
<td>5</td>
</tr>
<tr>
<td>11</td>
<td>L1 TAD 6</td>
<td>000 1 00000110</td>
<td>774</td>
</tr>
<tr>
<td>12</td>
<td>ISZ 6</td>
<td>010 0 00000110</td>
<td>1030</td>
</tr>
<tr>
<td>13</td>
<td>ISZ 4</td>
<td>010 0 00000100</td>
<td>1028</td>
</tr>
<tr>
<td>14</td>
<td>JMP L1</td>
<td>101 0 00001011</td>
<td>2571</td>
</tr>
<tr>
<td>15</td>
<td>DCA 7</td>
<td>011 0 00000111</td>
<td>1543</td>
</tr>
<tr>
<td>16</td>
<td>HLT</td>
<td>111 0 00000000</td>
<td>3584</td>
</tr>
</tbody>
</table>

*Figure 5(d) : Program to Add Four Integers*
<table>
<thead>
<tr>
<th>Address</th>
<th>Contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>5</td>
</tr>
<tr>
<td>1</td>
<td>6</td>
</tr>
<tr>
<td>2</td>
<td>7</td>
</tr>
<tr>
<td>3</td>
<td>8</td>
</tr>
<tr>
<td>4</td>
<td>-3</td>
</tr>
<tr>
<td>5</td>
<td>0</td>
</tr>
<tr>
<td>6</td>
<td>0</td>
</tr>
<tr>
<td>7</td>
<td>-</td>
</tr>
<tr>
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<td>10</td>
<td>5</td>
</tr>
<tr>
<td>11</td>
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<td>1030</td>
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<td>1028</td>
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<td>15</td>
<td>1543</td>
</tr>
<tr>
<td>16</td>
<td>3584</td>
</tr>
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</table>

**Figure 5(e) : Memory Map**
Figure 6: CADAT System