Application Developer's Tutorial for the CSM Testbed Architecture

Phillip Underwood and Carlos A. Felippa

Lockheed Missiles and Space Company, Inc.
Palo Alto, California

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1. Introduction

This tutorial is an extension of [1]† Appendix D to include a simple interface to GAL-DBM (Global Access Library - Database Management), the database management system for the CSM Testbed Architecture. GAL-DBM is described in [2]. The goal is to present a complete, but simple, introduction to using both CLIP (Command Language Interface Program) and GAL to write a NICE Processor. To achieve this goal the first author has extended the second author's work to include the interface to GAL. Much of the previous text describing commands and CLIP has been retained to make the tutorial stand alone.

Before beginning this tutorial, you should be familiar with the CSM Testbed Architecture (NICE). As a minimum you should: read [1] Appendices C & D — make sure you understand the use of the CLIP entry points, which may require a reading of [1]; and read [2] — at least be familiar with the ideas for nominal datasets.

The example Processor presented here is still quite simple as production Processors go, but is no longer trivial. It requires one to two weeks to put together. The Processor solves a two-dimensional elastostatic problem by a directly-formulated* Boundary Element Method (BEM), and is named, appropriately, DBEM2.

The “kernel” of the Processor is a BEM-program adapted from the book _Boundary Element Methods in Solid Mechanics_ by S. L. Crouch and A. M. Starfield [4]. The program is called TWOBI and is presented in Appendix C of the book; it is based on the boundary-integral theory covered in Section 6 therein.

The program is appropriate as an example of the use of interactive techniques because the input data are fairly simple but the commands are of multiple-item type and thus serve to illustrate things like phrases, item lists, qualifiers, and defaults. The program is somewhat weak for illustrating the true use of a database. However, it is relatively simple, so we can concentrate on the mechanics of using GAL. Ways to extend DBEM2 and use GAL for more complex problem solving are discussed as we proceed through this example Processor.

To prepare the reader for subsequent sections we need to cover some background material on the GAL-DBM [2].

The two-level conceptual model of GAL must be understood; see [2], §2.2. The two-levels of data in nominal GAL are named datasets and named records. The GAL database file is usually called a Library. In this Library are books (datasets) and in the books are chapters (records). In this tutorial example we will only consider one active Library; however, complex Processors can have several active GAL-DBM Libraries. We will have several datasets and several associated records in our one Library.

† Numbers in brackets refer to references at end of report.

* The term _direct formulation_ refers to the technique used in deriving the governing boundary-integral equations. Direct methods are formulated from the start in terms of physical quantities such as displacement and stress fluxes. On the other hand, indirect methods are formulated in terms of source strength distributions, which have no direct physical meaning and are eventually eliminated following spatial discretization.
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file is usually called a Library. In this Library are books (datasets) and in the books
are chapters (records). In this tutorial example we will only consider one active Library;
however, complex Processors can have several active GAL-DBM Libraries. We will have
several datasets and several associated records in our one Library.

Datasets are usually chosen as functional groups of data records. If you are fam-
iliar with the programming language C, a dataset is analogous to a structure and GAL
records are just like member definitions within a structure. In the boundary element
method and in other discrete element methods, such as finite element methods, typical
functional groups of data are geometry, material, boundary conditions, loadings, elements,
matrix (coefficient or stiffness), system vectors (right-hand-side(s), solution(s),
...), stress/strain/resultants, etc.

For example, we may decide to have a dataset named GEOMETRY and in this dataset
there may be records named NUMBER.NODES, NODES, COORDS (or X-COORDS, Y-COORDS,
...), etc. For this tutorial the geometry data is stored in a dataset named SEGMENT; see
§3.1.

In this tutorial, we will use fixed or "hardwired" names for the datasets and records.
This simplifies the Processor and lets us get on with how to use GAL, without getting into
the complex issues of tables and their management to relate Library dataset names and
record names to the names used internally in the Processor. In addition, the command
set would have to be expanded to include commands to bind the external database names
with the internal Processor names. In general fixed names work well with tightly
coupled Processors, because they don't interact much with other Processors. The fixed names also
make the Processor much easier for the user to run. The user doesn't have to remember
as many commands or keep track of where the data really are. Loosely coupled networks
of Processors may need the capability to use datasets and records of any given name.
However, even fixed dataset and record names can be changed by using the *rename
dataset and/or the *rename record directives; see [3], §5.1 & §5.2.

Another GAL-DBM feature is word addressability. With this feature a particular
entry in a dataset-record can be extracted or stored [2], §5.1. For example, the \textit{i}-th entry
for the geometry nodes, \texttt{X-COORDS} \& \texttt{Y-COORDS}, can easily be extracted. This feature
can be used for out-of-core techniques. In actual practice experienced Processor developers
use a local dynamic memory manager for out-of-core methods, because they are usually
more efficient than GAL. GAL was designed for efficient use of archival data. These are
advanced concepts. They are not covered in this tutorial, but the developer should be
aware of these issues - especially for large problems.
3. The Data Structures

Following sound practice, we begin by designing the data structures. The task is more complicated for DBEM2 than for the simple program in [1] Appendix C. We will retain DBEM2 as a single Processor, but add a global database that corresponds to the data structures and functions in DBEM2. Possible avenues to explore in separating DBEM2 into several Processors are presented, but not pursued in depth. The main use of the database in this example is to archive problem data for a restart, archive problem data to document what was done, and archive problem data so that some old data can be used with new input to solve a slightly different problem. Our main goal, to illustrate the mechanics of using GAL, is well served by this approach.

The task is simplified by the following considerations:

1. The Processor presented here is isolated from others. There is no need to transact business with a global database generated by other Processors.
2. DBEM2 makes use of only one matrix, which is generally unsymmetric and full. There being no need to make use of sparse storage formats, an ordinary FORTRAN array suffices.
3. Everything is assumed to fit in core at one time. Not having to deal with auxiliary storage avoids many complications.
4. The internal data structures and the GAL dataset-record structures are the same.

All data that have to be shared among many parts of DBEM2 are accommodated in labelled common blocks. The first author is not in favor of using labelled common blocks for moving data from one subroutine to another, but it is retained here for expediency. Thus in the present Processor several blocks are used to group data according to function. Furthermore, the blocks are declared in separate files whose extension is inc. These files are inserted where they are needed via INCLUDE statements. The use of INCLUDE enforces consistency (everything is declared only once) and makes maintenance and modification much easier.

Remember that in §1.0 we said we would use fixed dataset names and record names. Thus, here we will use, wherever possible, the labelled common block name for the dataset name and the variable names in the labelled common block will be the same as the record names. In a few cases we will have to break one labelled common block into two or three separate datasets to achieve a realistic functional group for the data. In these cases we create new, but hopefully meaningful, names.
3.1 The Segment Data

We begin by setting up the data for boundary segments, which is placed in file segment.inc. The maximum number of segments is parameterized to be MAXSEG, which is set to 20 in the version listed below.

The DBEM2 user will be allowed to define segments in any order and give them arbitrary numbers from 1 through MAXSEG, so we need a "marker" array that tells which segments have been defined. We also need a counter of how many boundary elements are in each defined segment. Then there are the geometric arrays: the x and y coordinates of the end points. Finally, there are the boundary condition arrays: one integer code (related to that used by reference 4) and two floating-point arrays of prescribed shear and normal values. Here is a list of the file that groups this information:

This is the file segment.inc

```plaintext
common /SEGMENT DATA/
$  segdef, numel, xbeg, ybeg, xend, yend, kode, bvs, bvn
integer MAXSEG
parameter (MAXSEG=20) ! Maximum no. of boundary segments
integer segdef(MAXSEG) ! Segment definition tag
integer numel(MAXSEG) ! Number of BE divisions of segment
real xbeg(MAXSEG) ! X-coord of starting segment point
real ybeg(MAXSEG) ! Y-coord of starting segment point
real xend(MAXSEG) ! X-coord of ending segment point
real yend(MAXSEG) ! Y-coord of ending segment point
integer kode(MAXSEG) ! Segment BC code
real bvs(MAXSEG) ! Prescribed shear value
real bvn(MAXSEG) ! Prescribed normal value
```

The style used in this INCLUDE file will be followed for all others. There is a COMMON declaration that lists the shared variables. Then each variable is declared on a separate line. The variable name is followed by an inline comment that provides a short description of the function of each variable. This brief documentation should be entered at the time you prepare or update the INCLUDE file; if you leave it for later, it'll never be done.

When you get farther into this tutorial you will see that the segment data described here is generated in at least two subroutines and possibly three, if you define non-default data for the number of BE divisions of a segment. Thus to maintain a functional breakdown of data in the database, these data are divided into two datasets: 1) SEGMENT with records named SEGDEF, NUMEL, XBEG, YBEG, XEND, and YEND that hold segdef, numel, xbeg, ybeg, xend, and yend one-dimensional array data; and 2) BCVALUES with records named KODE, BVS, and BVN that hold the kode, bvs, and bvn one-dimensional data. The default value for numel is 1, however you may enter other values by using the DEFINE ELEMENTS command. Data generated by this command are written over the default data created in the DEFINE SEGMENTS code.
Although this simple structure for the datasets and records may seem trivial, it is quite common even in a complex Processor. That is, the data in the database are often structured just as the data are used in the code. Also, note that, the two level structure of GAL, named datasets with named records, lends itself to a functional grouping of the data with names that are easy to relate to the data generated and used by the Processor. This simple structure for the datasets and the records will be used throughout the DBEM2 Processor.

In summary, we have defined two datasets with their associated records as:

Dataset - SEGMENT
  Records - SEGDEF
    NUMEL
    XBEG
    YBEG
    XEND
    YEND

and

Dataset - BCVALUES
  Records - KODE
    BVS
    BVN
3.2 The Material Data

Since we are dealing with a homogeneous elastic isotropic material and we ignore thermal effects, the material is fully characterized by two properties: the elastic modulus $E$ and the Poisson's ratio $\nu$. These two are collected in file `material.inc`:

This is the file `material.inc`

```fortran
common /MATERIAL/ em, pr
real   em     ! Elastic modulus
real   pr     ! Poisson's ratio
```

Here the entry for the database is very simple to design. We use a dataset named `MATERIAL` with two records named `EM` and `PR` to store the values for `em` and `pr`. Thus, this dataset has the following two level structure:

Dataset - MATERIAL
  Records - EM
    PR
3.3 The Symmetry Data

The program allows one or two lines parallel to the coordinate axes to be specified as axes of symmetry. For example, \( x = 2.5 \) or \( y = -1.50 \), or both. Three pieces of data accommodate this information: one symmetry tag (0=none, 1=symmetry about \( x = a \), 2=symmetry about \( y = b \), 3=double symmetry), and the values of \( a \) and \( b \) as appropriate. The necessary declarations are placed in the file `symmetry.inc`:

```plaintext
* This is the file symmetry.inc

common /SYMMETRY_DATA/
$ ksym, xsym, ysym
integer ksym ! symmetry tag
real xsym, ysym ! symmetry about x=a & y=b values
```

These data are very similar to the MATERIAL dataset above, so we choose a similar design for the SYMMETRY dataset and records. Here it is:

Dataset - SYMMETRY
Records - KSYM
XSYM
YSYM
3.4 The Prestress Data

The program allows a constant initial-stress field to exist in the undeformed medium. This prestress tensor field is defined by the three components $\sigma_{xx}^0$, $\sigma_{yy}^0$ and $\sigma_{xy}^0$. If undefined, these three values are assumed to be zero. File *prestress.inc* contains the appropriate declarations:

```
* This is the file prestress.inc

common /PRESTRESS/ sxx0, syy0, sxy0
real sxx0 ! Prestress (initial field stress) sigma_xx
real syy0 ! Ibid., for sigma_yy
real sxy0 ! Ibid., for sigma_xy
```

Prestress data are especially important for analysis of unbounded domains, for which they assume the role of conditions at infinity. For example, suppose that we want to analyze the effect of a hole in an infinite region under uniform uniaxial stress, say $\delta_{xx}$. Then we set $\sigma_{xx}^0 = \delta_{xx}$, $\sigma_{yy}^0 = \sigma_{xy}^0 = 0$ in the input data.

Again very similar to MATERIAL and SYMMETRY, so we have:

Dataset - PRESTRESS
Records - SXX0
         SYY0
         SXY0
3.5 The Element Data

The most voluminous data are that pertaining to the boundary elements, since typically there will be many elements per segment. The information is collected in file `element.inc`, which reads

```plaintext
* This is the file element.inc
* common /ELEMENT DATA/
$ numbe, xme, yme, hleng, sinbet, cosbet, kod, c, b, r, x
integer MAXELM, MAXEQS
parameter (MAXELM=100) ! Maximum no. of boundary elements
parameter (MAXEQS=2*MAXELM) ! Maximum no. of discrete equations
integer numbe ' Total number of boundary elements
real xme(MAXELM) ! X-coor of element midpoint
real yme(MAXELM) ! Y-coor of element midpoint
real hleng(MAXELM) ! Half length of element
real sinbet(MAXELM) ! Sine of (element,x) angle
real cosbet(MAXELM) ! Cosine ibid.
integer kod(MAXELM) ! Elem BC code (copies seg code)
real b(MAXEQS) ! Prescribed boundary values
real c(MAXEQS,MAXEQS) ! Influence coefficient matrix
real r(MAXEQS) ! Forcing (RHS) vector
real x(MAXEQS) ! Solution vector
```

The elements arrays such as xme, yme, etc are parametrized in terms of the maximum number of elements MAXELM.

This block also contains arrays used to set up and solve the IEM equation system, namely c, r, b and x. These are parameterized in terms of the total number of equations MAXEQS, which of course is twice MAXELM.

Now things get a little more complicated. All these data could be stored in one dataset with several records, but it is better to use a more functional design as in §3.1. So, the data in the database are organized according to where it is generated. All the actual element data are computed in the BUILD subroutine, so the first dataset of this group is named ELEMENT with records named NUMBE, XME, YME, HLENG, SINBET, COSBET, KOD, and B that hold the integer number numbe, and the one-dimensional arrays xme, yme, hleng, sinbet, cosbet, kod, and b. These data are analogous to the element stiffness data in a typical finite element code. However, for boundary elements there are no individual element stiffnesses, only a global system coefficient matrix. The element data just contain the information needed to compute the global coefficient matrix.

The data for the arrays c and r are created in the GENERATE subroutine. However, in general the system matrix, analogous to a global stiffness matrix in a finite element code, is computed in one subroutine, such as an assembler, and the right-hand-side, the forcing function, is computed in another subroutine. Thus, two datasets are added to the database for these data. The first dataset is named COEFF with a record named C that contains the data for the coefficient matrix, c. The second dataset is named RHS with a record named R that contains the one-dimensional array r.
The final dataset for this group of data is for the solution vector, \( x \). This one-dimensional array is computed by SOLVE, very similarly to a typical finite element code. So, we name the dataset SOLUTION with a named record \( X \) to archive the solution vector \( x \).

In summary, we have defined four datasets with their associated records as:

- **Dataset - ELEMENT**
  - Records - NUMBE
  - XME
  - YME
  - HLENC
  - SINBET
  - COSBET
  - KOD
  - B

- **Dataset - COEFF**
  - Record - C

- **Dataset - RHS**
  - Record - R

- **Dataset - SOLUTION**
  - Record - X
3.6 The Field Location Data

This block of data pertains to the location of field points at which stresses and displacements are to be calculated once the boundary solution is obtained. The program allows these locations to be specified as equally spaced points along straight lines defined by the user. Up to MAXLIN (= 100 in the version below) lines can be defined. The locations are specified by giving the x and y coordinates of the first and last points on the line, and the number of intermediate points (≥0) to be “collocated” between the first and last points. An isolated point may be specified by making the first and last point coincide.

All of this information is gathered in the file output.inc:

```
* This is the file output.inc
* common /OUTPUT DATA/
$ lindef, nintop, xfirst, yfirst, xlast, ylast
integer MAXLIN
parameter (MAXLIN=100)
integer lindef(MAXLIN) ! Line definition tag
integer nintop(MAXLIN) ! No. of intermediate points on line
real xfirst(MAXLIN) ! X-coor of first point on line
real yfirst(MAXLIN) ! Y-coor of first point on line
real xlast(MAXLIN) ! X-coor of last point on line
real ylast(MAXLIN) ! Y-coor of last point on line
```

This group of data are similar to other groups of one-dimensional arrays, such as the segment data in §3.1. So, we use a dataset named FIELD with named records LINDEF, NINTOP, XFIRST, YFIRST, XLAST, and YLAST to store the one-dimensional arrays lin-def, nintop, xfirst, yfirst, xlast, and ylast. In outline form this database data structure is:

Dataset - FIELD
Records - LINDEF
NINTOP
XFIRST
YFIRST
XLAST
YLAST
3.7 The Database Data

The DBEM2 Processor subroutines that communicate with the GAL-DBM need to know the logical device index (ldi) of the library (database) being used; see [2] §2.4 and the description of the DBophe subroutine in §6. So, this globally used information is kept in the labelled common block DATABASE. There is no dataset and record associated with these data because it is not archival data; it is temporary — only used for the run at hand.

This information is gathered in the file database.inc:

```plaintext
This is the file database.inc

common /DATABASE/ ldi
integer ldi ! GAL Library logical device index
```

This concludes the design of the important data structures for the internal data representation and the global database. Next we pass to the design of a command set to control logic of DBEM2.
The Commands

4. The Commands

Having described the data and the datasets for the database, we have now to design an appropriate set of commands to perform operations on the data and the database. The writers found it convenient to choose commands headed by the following action verbs:

- CLEAR
- OPEN
- DEFINE
- BUILD
- GENERATE
- SOLVE
- PRINT
- CLOSE
- STOP

Why these particular commands? Partly from a preliminary study of the problem, partly from wishes to get several command formats so that the use of many of the entry points described in [1] would be illustrated.

It turns out that the last wish (of illustrating various command formats) makes the command set a bit inconsistent, but that should not cause a great deal of concern. After all, it's only an example.

Another Processor developer faced with the same problem (even a simple problem like this one) may in fact come up with a radically different set of commands that accomplishes virtually the same thing.

We next describe briefly what the commands do.

- CLEAR: Initializes all Tables maintained by the Processor and sets some default values.
- OPEN: Opens a GAL-DBM Library to store problem data and/or load previously stored problem data.
- DEFINE: Enters data that are used in the definition of the problem to be solved. The DEFINE verb will be followed by another keyword that makes the data more specific.
- BUILD: Indicates that the problem-definition phase is complete, and calls for the generation of the discrete governing equations.
- GENERATE: Triggers the assembly of the influence coefficient matrix and forcing vector.
- SOLVE: Triggers the solution for the unknown boundary variables.
- PRINT: Prints displacements and stresses at boundary points and at specified field points.
- CLOSE: Closes the open GAL-DBM Library.
- STOP: Closes all open GAL-DBM Libraries and terminates execution of the processor.
5. Starting at the Top

We are going to build the Processor Executive “top down”. For this relatively small Processor it probably doesn’t make much difference whether we do it top-down, bottom-up or inside-out. But adhering to this approach makes life easier for bigger Processors.

Following the top-down approach we must do the main program first. Here it is:

```fortran
* * Computer Program for the Two-Dimensional Direct
* Boundary Element Method (DBEM2)
* Adapted from program TWOBI in the book Boundary Element Methods
* Methods in Solid Mechanics by S. L. Crouch and A. M. Starfield,
* C. A. Felippa to exemplify conversion to interactive operation
* via CLIP and the use of a global database GAL-DBM.
* program DBEM2
* implicit none
character*8 CCLVAL, verb
integer ICLTYP
1000 call CLREAD ('DBEM2> ','
$ ' CLEAR, OPEN, DEFINE, BUILD[/LOAD | STORE]kk '//'
$ 'GENERATE[/LOAD | STORE], SOLVE[/LOAD | STORE]kk '//'
$ 'PRINT, CLOSE, STOP')
if (ICLTYP(1) .le. 0) then
   print *, '*** Commands must begin with keyword'
else
   verb = CCLVAL(1)
   call DO_COMMAND (verb)
end if
1000 go to
end
```

In DBEM2 the top-level command must start with an action verb, hence the error check. The prompt is the name of the Processor: this is a convention followed in the NICE system.

The top level of all Processors looks very much the same, regardless of the complexity of what lies underneath. This is not surprising if you note that all Processors fit the “do forever” model illustrated in [1] §C.2.

The next level is DO_COMMAND, which is again a “case” statement that branches on the action verb:
* Starting at the Top

Top level command interpreter for DBEM2

subroutine DO_COMMAND (verb)

implicit none
character key*8, qual*8, verb*(_)
integer nq
logical CMATCH

key = verb
if (CMATCH (key, 'B^UILD')) then
  qual = ' ';
call CLOADQ (' ', -1, qual, 0, nq)
  if ( nq .eq. 1 ) then
    call UPCASE (qual)
      if ( qual(1:1) .eq. 'L' ) then
        qual = 'LOAD'
      else if ( qual(1:1) .eq. 'S' ) then
        qual = 'STORE'
      else
        print*, ' Illegal qualifier: ', qual, ' for BUILD.'
        print*, ' BUILD not performed.'
        return
    end if
  end if
  call BUILD ( qual )
else if (CMATCH (key, 'CLE^AR')) then
  call CLEAR
else if (CMATCH (key, 'CLO^SE')) then
  call DB_CLOSE
else if (CMATCH (key, 'D^EFINE')) then
  call DEFINE
else if (CMATCH (key, 'G^E)ERATE')) then
  qual = ' ';
call CLOADQ (' ', -1, qual, 0, nq)
  if ( nq .eq. 1 ) then
    call UPCASE (qual)
      if ( qual(1:1) .eq. 'L' ) then
        qual = 'LOAD'
      else if ( qual(1:1) .eq. 'S' ) then
        qual = 'STORE'
      else
        print*, ' Illegal qualifier: ', qual, ' for GENERATE.'
        print*, ' GENERATE not performed.'
        return
    end if
  end if
end if
   call GENERATE ( qual )
else if (CMATCH (key, 'HELP')) then
   call HELP
else if (CMATCH (key, 'OPEN')) then
   call DB_OPEN
else if (CMATCH (key, 'PRINT')) then
   call PRINT
else if (CMATCH (key, 'SOLVE')) then
   qual = ' '
   call CLOADQ (', -1, qual, 0, nq)
   if ( nq .eq. 1 ) then
      call UPCASE (qual)
      if ( qual(1:1) .eq. 'L' ) then
         qual = 'LOAD'
      else if ( qual(1:1) .eq. 'S' ) then
         qual = 'STORE'
      else
         print*, ' ** Illegal qualifier: ', qual, ' for SOLVE.'
      print*, ' SOLVE not performed.'
      return
   end if
end if
   call SOLVE ( qual )
else if (CMATCH (key, 'STOP')) then
   call STOP
else
   print *, '*** Illegal or ambiguous verb: ', key
end if
return
end

Note that the tests are ordered so that keywords are alphabetically sorted. This makes it easier to insert new keywords without forgetting to expand roots of existing ones. For example, suppose you want to insert a PLOT command for your favorite graphic device; inserting it just before the test for PRINT makes it easy to spot that the root for the latter has to be expanded to PR.

For the keywords, BUILD, GENERATE, and SOLVE we are also looking for a possible qualifier. The existence of this qualifier is determined by a call to CLOADQ; see [1] §8.3. If a qualifier is found, then we check to see if the value is either LOAD or STORE. The LOAD qualifier signifies that the data are to be loaded from the database instead of being computed. The STORE qualifier signifies that the data are to be stored in the database after it has been computed.

UPCASE is a CSM Testbed (NICE) architectural utility that converts its argument to uppercase.
6. Starting and Stopping

The CLEAR subroutine is quite simple, as it only has to zero out the model definition tables:

```
Initialize tables, set default values

subroutine CLEAR

implicit none
include 'segment.inc'
include 'element.inc'
include 'material.inc'
include 'symmetry.inc'
include 'prestress.inc'
include 'output.inc'
integer i

do 1500 i = 1,MAXSEG
   segdef(i) = 0
   xbeg(i) = 0.0
   xend(i) = 0.0
   ybeg(i) = 0.0
   yend(i) = 0.0
   numel(i) = 0
   kode(i) = 0
   bvs(i) = 0.0
   bvn(i) = 0.0
1500 continue

do 2000 i = 1,MAXLIN
   lindef(i) = 0
2000 continue
numbe = 0
ksym = 0
em = 1.0
pr = 0.0
sxxo = 0.0
syyo = 0.0
syyo = 0.0
print 'Tables initialized'
return
end
```

The function of the arrays is explained in §3.
Next, you may wish to **OPEN** a GAL Library (database) to use for loading previously computed data and/or to store data computed during a run. (CLEAR does not have to be used before **OPEN** and vice versa.) The **OPEN** command has the form

\[
\text{OPEN/[/Qualifier]} \quad \text{LIB} = \text{library.name}
\]

The brackets, [], around the Qualifier signify that the Qualifier is optional. The Qualifier is to describe the characteristics of the Library to be opened. The accepted values are **NEW**, **OLD**, **ROLD**, and **SCR**. If no qualifier is given the value defaults to **COLD**. The meaning of these values is given in [2], Table 6.5. Most applications may ignore the qualifier.

The name of the file that is the GAL Library (library.name) must be entered. This is a valid file name for the computer system you are using. On a UNIX system if you use a pathname that contains '/'s the file name must be enclosed in single quotes (e.g., '/usr/king/kong/new/york'). Otherwise CLIP will try to interpret the directories and files as qualifiers.

Experienced NICE users pick file names for GAL Libraries that are descriptive of the problem and usually use the file extension **gal**. This way you can easily find the database files and the name should remind you of the problem.

The **OPEN** command produces a call to the **DB_OPEN** subroutine:

```fortran
* * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * *

Open GAL Library (Database)

subroutine DB_OPEN

implicit none
include 'database.inc'
character*80 CCLVAL
character*81 libnam
character*11 key
character*4 qual
integer ICLNIT, ICLSEK, ICLTYP, LEBETB, LMOOPEN
integer nq

if ( ICLNIT ( ) .lt. 3 ) then
  call CLREAD ( ' OPEN: Enter [/QUAL] LIB = LIB_NAME > ' ,
                 ' ' )
end if

libnam = '
key = 'COLD/GAL82'
call CLOADQ ( ' ', -1, qual, 0, nq)
if ( nq .eq. 1 ) then
  call UPCASE (qual)
```

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if ( qual(1:1) .eq. 'N') key = 'NEW/GAL82 '
if ( qual(1:1) .eq. 'O') key = 'OLD/GAL82 '
if ( qual(1:1) .eq. 'R') key = 'ROLD/GAL82 '
if ( qual(1:1) .eq. 'S') key = 'SCR/GAL82 '
end if
if ( ICLSEK (0, 'LIB') .ne. 0 ) then
  if ( ICLTYP (0) .gt. 0 ) libnam = CCLVAL(0)//'
  if ( libnam .eq. '' ) then
    print*, $
    ' Cannot find LIB_NAME. Format is LIB = LIB_NAME.'
    print*, ' No library opened.'
    return
  end if
else
  print*, ' Cannot find keyword LIB; No library opened.'
  return
end if
ldi = LMOPEII ( key, 0, libnam, 0, 500 )
if ( ldi .eq. 0 ) then
  print*, ' Unable to open library: ', libnam(l:LEHETB(libnam))
else
  call GMSIGII ('DBEM2')
end if
return
end

The call to CLREAD, [1], §2.7, at the beginning of DB OPEN prompts the user for needed data if the minimum number of items (i.e., the function ICLHIT, [1], §9.5) needed to carry out the OPEN command is not found.

In DB OPEN the call LMOPEII is used to open the GAL Library for reading (loading) and/or writing (storing) data. The subroutine GMOPEN can also be used, but for this beginning example it is a little complex. The use of LMOPEII here easily lets us use one database for one problem. If multiple databases are required then GMOPEN is the one to use. A description of GMOPEN and LMOPEII can be found in [2], §6.4.

If the Library is successfully opened (a non-zero ldi is obtained) we call GMSIGII with the name of the Processor as the argument. This places the Processor name in the Table of Contents (TOC) data for the database. This is a recommended procedure; see [2], §10.8.

The DB_CLOSE subroutine carries out the action requested by the CLOSE command. This command is used to close the currently active GAL Library. This command is used if you were finished with some problem then wished to work on another problem that required another GAL Library. So you CLOSE the old one before OPENing the new file containing the other GAL Library.

The DB_CLOSE subroutine is quite simple:
Close GAL Libraries (Databases)

subroutine DB_CLOSE

implicit none
include 'database.inc'

call GMCLOSE (ldi, 0, 100 )

return
end

The only call is to GMCLOSE [2], §6.2, which closes the active library associated with ldi. For more advanced applications with multiple libraries in use, a specific ldi may also be an input item. So a specific library could be closed after it is no longer needed.

Relatively simple is the STOP subroutine:

Terminate the run

subroutine STOP

call GMCLOSE (0, 0, 100 )

print*, 'Hope you enjoyed the ride!'
call CLPUT ('STOP ')
end

The call to GMCLOSE insures that any GAL Libraries that may be open are properly closed, so that no database data are lost. The call to CLPUT sends the *stop directive to CLIP. This is the preferred way to exit a NICE Processor, because it allows you to run the Processor in network mode. However, if the network mode (SuperCLIP) is not available on your computer, replace this line with the FORTRAN stop statement.

A description of GMCLOSE can be found in [2], §6.2. A description of CLPUT can be found in [1], §2.4. And a description of *stop can be found in [3], §59.1.

Now on to DEFINE to define the problem to solve.
7. Defining the Problem

The **DEFINE** command introduces problem-definition data. It is convenient to break up the definition into several types of data, which correspond closely to the data-structure grouping discussed in §3. Each type is identified by a keyword that immediately follows **DEFINE**. The keywords are:

- **SEGMENTS**
  Specifies the straight-line segments that make up the boundary of the problem to be solved.

- **ELEMENTS**
  Specifies into how many boundary elements each segment will be divided.

- **BOUNDARY CONDITIONS**
  Specifies the boundary conditions that apply to each boundary segment.

- **SYMmetry CONDITIONS**
  Specifies the symmetry conditions, if any, that apply to the problem to be solved.

- **MATERIAL**
  Specifies constitutive properties of the material.

- **PRESTRESS**
  Specifies prestress data in the form of initial stress components.

- **FIELD**
  Specifies the location of field points at which displacement and stresses are to be evaluated and printed later.

Subroutine **DEFINE**, unlike **CLEAR**, **OPEN**, **CLOSE** or **STOP**, branches as per the second keyword:

```fortran
subroutine DEFINE
  implicit none
  character key*8, CCLVAL*8
  integer ICLTYP
  logical CMATCH
  if (ICLTYP(2) .le. 0) then
    print *, 'No keyword after DEFINE'
    return
  end if
  key = CCLVAL(2)
  if (CMATCH (key, 'BOUND')) then
    call DEFINE_BOUNDARY_CONDITIONS
  else if (CMATCH (key, 'ELEMENTS')) then
    call DEFINE_ELEMENTS
  else if (CMATCH (key, 'FIELD')) then
    call DEFINE_FIELD_LOCATIONS
  else if (CMATCH (key, 'MATERIAL')) then
    call DEFINE_MATERIAL
  else if (CMATCH (key, 'PRESTRESS')) then
```

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call DEFINE_PRESTRESS
else if (CMATCH (key, 'SEGMENTS')) then
call DEFINE_SEGMENTS
else if (CMATCH (key, 'SYMMETRY')) then
call DEFINE_SYMMETRY_CONDITIONS
else
  print *, '*** Illegal or ambiguous keyword ', key,
$      ' after DEFINE'
end if
return
end

The program begins checking whether a keyword actually follows DEFINE. If so it compares them in the usual matter and calls appropriate input subroutines. These are described next.
7.1 A Digression: The Basics of Using GAL-DBM

Because the basic mechanics of writing (storing) and reading (loading) data are very similar throughout all the subroutines that load and store data, these basics are presented before we present the details of each subroutine. Thus, in the discussion we can focus on the important details within the basic outline presented here.

7.1.1 Storing GAL Data

The first step is to OPEN the GAL Library that is to contain the data to be stored; see §6. If the Library is already open, nothing needs to be done.

Next, each dataset needs to be installed in the Library. This is easily done with a call to GMPUNIT, [2], §7.11. In this tutorial example the dataset is always installed (even if it already exists). So, if you store the same named data more than once during the same run, the old dataset will be marked as deleted and a new dataset of the same name will appear in the Library. In general this is a safe practice, because old data are still there until the Library is packed. Thus, old data can be retrieved by enabling a deleted dataset. (See *pack [3], §8.1 and *enable [3], §24.1.)

Then for each record of data we wish to store we must construct the record name, followed by writing the data. The record name (rname is used in the code) is constructed by subroutine GMCORII [2], §10.6. The data are written with a call to GMPUTN [2], §9.9.

In this tutorial example we have chosen the record key (name) and its associated record cycles (record group), [2], §5.1, to correspond to the array name and array indices used in the dimensioned arrays in the code. For example, if we have a one-dimensional array named number and we have used number(1) through number(12), we construct the record name to be number.1:12. Thus, number(3) in the code equals number.3 in the database. Here we have constructed a record group containing 12 records with each record containing one number. If you look at the record attributes (use the *rat directive [3], §49.4), you will see the record number has a low cycle = 1, a high cycle = 12, and a logical size = 1 (1 number).

Somewhat more complex is a two-dimensional array (a matrix). In this tutorial there is only one matrix, c, the system coefficient matrix. This matrix is square, n by n, so we choose the record name to be c.1:n. We have as many records as we have columns in the matrix. Thus, each record has a logical size equal to n, i.e., each record contains the n numbers for the column it represents. For example, c.5 contains the n numbers for column 5 of the matrix, c. See §8.0 for the implementation details.

The simplest case is the record that contains only one value. For this case the record name does not have to contain any group cycle numbers; see [2], §5.1.
A Digression: The Basics of Using GAL-DBM

Finally, after the data for a dataset are stored in the GAL Library, the GAL-DBM buffers should be flushed. This insures that all the data are actually written to the GAL Library (file). Thus, if the next thing that happens causes your run to fail, you still have all the data properly stored up to that point.

In summary, to store data in an open GAL Library the following steps are needed:

1) Install the dataset — GMPUT
2) Then for each record repeat these steps
   a) Construct the record name — GMCORII
   b) Store the data — GMPUTII
3) Finish with a buffer flush — GMFLUB

7.1.2 Loading GAL Data

Again, the first step is to OPEN the GAL Library that contains the data to be read; see §6. If the Library is already open, nothing needs to be done.

Next, for each dataset to be loaded the dataset sequence number must be determined for use in subsequent calls to other GAL subroutines. Given the dataset name the integer function LMFIND returns the dataset sequence number [2], §7.5.

Then for each record name in the dataset we must determine the record group cycles, so all the data in the records can be loaded. The low and high record group cycles (ilow and ilhigh, respectively) are returned by the subroutine GMCEGY [2], §9.3.

Next, the record name is constructed so that all the records are loaded in one read operation. The record name is constructed by a call to GMCORII [2], §10.6.

Finally, the data are loaded from the GAL Library into a single variable or an array. The call to GMGETII loads numeric data [2], §9.5.

In summary, we have the following simple outline for loading data:

1) Find dataset sequence number — LMFIND
2) Then for each record repeat these steps
   a) Get record cycle information — GMCEGY
   b) Construct record name — GMCORII
   c) Load the data — GMGETII

This is about as simple as it can be done for arrays. However, if only a single value record has been stored, it can be loaded with only steps 1 and 2c. The record name does not need to contain the cycle number(s) [2], §5.1.
A sophisticated Processor should do some more error checking. The sophisticated reader should be aware of the test for an error condition, LMERCD [2], §14.5. Calls to LMERCD are typically made after all reads (GMGETx [2], §9.5) and many times after all writes (GMPUTx [2], §9.9). Also, the subroutine GMGETx returns two arguments, n and m, that contain information about how much data have been read. These arguments can also be used for error checking.

In the code that follows in §7.2, the LOAD and STORE operations are code inline in the applicable subroutine. A higher level of abstraction can be used by writing cover subroutines that load and store data. These subroutines then call the appropriate GAL subroutines to load and store. This results in cleaner code, i.e., all database I/O is done in two subroutines; so a change of databases is easily accommodated. However, you do pay the price of a modest loss of efficiency because of the additional subroutine calls. The inline code is used here for tutorial purposes. Also for a small Processor, the additional abstraction is not really needed.

Now, let's look at the CLIP and GAL calls required to define the problem, either by interactively entering the data through commands then storing it, or by loading previously stored data.
7.2 Defining Segments

The DEFINE SEGMENT command introduces a series of segment-definition commands which are expected to have the form

\[
\text{SEGMENT } = i \quad \text{BEGIN} = x_i^{\text{beg}}, y_i^{\text{beg}} \quad \text{END} = x_i^{\text{end}}, y_i^{\text{end}} \\
[ \text{LOAD [ STORE ]}]
\]

where \(x_i^{\text{beg}}, y_i^{\text{beg}}\) are the \(x, y\) coordinates of the starting point of the \(i^{th}\) segment, and \(x_i^{\text{end}}, y_i^{\text{end}}\) are the \(x, y\) coordinates of the ending point. The segment list is terminated by an END command that takes the control back to the main program. In listing the coordinates, the following boundary traversal convention must be observed: a closed contour is traversed in the counterclockwise sense if the region of interest is outside the contour (a cavity problem), and in the clockwise sense if the region of interest is inside the contour (a finite body problem); see Figure 2-1, §2.

In the CLAMP metalanguage, the \(\mid\) says that one may specify either LOAD or STORE, but not both simultaneously. The specifications are shown in brackets, meaning that they may be omitted.

For example, to define and store a 4-segment boundary that encloses a square region whose corner points are (0,0), (4,0), (4,4) and (0,4), and which constitutes the region of interest, you say

```
DEFINE SEGMENTS
SEG=1 BEGIN=0,0 END=0,4
SEG=2 BEGIN=0,4 END=4,4
SEG=3 BEGIN=4,4 END=4,0
SEG=4 BEGIN=4,0 END=0,0
STORE
END
```

(Segments may be actually defined in any order; there is also no need to number them sequentially.)

The commands that enter the segment data, plus the LOAD, STORE and END command, are called subordinate commands, because they can appear if and only if the command DEFINE SEGMENT has been entered. The DEFINE SEGMENT command, which introduces the subordinate commands, is said to be the header command (it also goes by the names master command, parent command, etc.).

The STORE command is optional; you don't have to store the data in the GAL Library, unless you wish to keep it for later use.

If you have already defined the segments and stored the data, you can use the following command sequence to load the segment data from the database.
DEFINE SEGMENTS
LOAD
END

The processing of the segment-definition commands is carried out within subroutine DEFINE SEGMENTS:

```
* Read segment-definition data *

subroutine DEFINE_SEGMENTS
*
implicit none
include 'database.inc'
include 'segment.inc'
character*8 key, CCLVAL
character*20 rname
integer iseg, n, mseg, ICLYP, ICLVAL, ICLSEK
integer idsn, LMFIN, ilow, ihigh, nrec
real xy(2)
logical CMATCH

1000 call CLREAD (' Segment data> ',
$ ' Enter SEG=iseg BEG=xbeg,ybeg END=xend,yend&'
$ ' or LOAD or STORE&'
$ 'Terminate with END')

if (ICLYP(1) .le. 0) then
   print *, '*** Command must begin with SEG or END'
go to 1000
end if
key = CCLVAL(1)
if (CMATCH (key, 'END')) then
   return
else if (CMATCH (key, 'SEGMENT')) then
   iseg = ICLYP(2)
   if (iseg .le. 0 .or. iseg .gt. MAXSEG) then
      print *, '*** Segment number', iseg, 'out of range'
go to 1000
   end if
   segdef(iseg) = 1
   if (numel(iseg) .le. 0) numel(iseg) = 1
   if (ICLSEK(3, 'BEGIN') .ne. 0) then
      call CLVALF (' ', 2, xy, n)
      if (n .ge. 1) xbeg(iseg) = xy(1)
      if (n .ge. 2) ybeg(iseg) = xy(2)
   end if
   if (ICLSEK(3, 'END') .ne. 0) then
      call CLVALF (' ', 2, xy, n)
```

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if (n .ge. 1) xend(iseg) = xy(1)
if (n .ge. 2) yend(iseg) = xy(2)
end if
else if (CMATCH (key, 'L^DAD')) then
--- find dataset
idsn = LMFIND ( idi, 'SEGMENT ', 100 )
if ( idsn .eq. 0 ) then
    print*, ' Cannot find SEGMENT dataset; nothing LOADed.'
    go to 1000
end if
--- get record name cycles, construct record name & read data
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'SEGDEF ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'SEGDEF ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'I', segdef, $
            n, 0, 0, 0, 200 )
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'NUMEL ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'NUMEL ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'I', numel, $
            n, 0, 0, 0, 300 )
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'XBEG ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'XBEG ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', xbeg, $
            n, 0, 0, 0, 400 )
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'YBEG ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'YBEG ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', ybeg, $
            n, 0, 0, 0, 500 )
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'XEND ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'XEND ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', xend, $
            n, 0, 0, 0, 600 )
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'YEND ', nrec, ilow, $
            ihigh, 220 )
call GMCORN ( rname, 'YEND ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', yend, $
            n, 0, 0, 0, 700 )
else if (CMATCH (key, 'S'TORE')) then
Defining the Problem  Defining

--- install dataset
call GMPUNT ( idi, 'SEGMENT ', idsn, 16, 1000 )

--- determine largest value of segdef
do 100 n=MAXSEG,1,-1
   if ( segdef(n) .ne. 0 ) then
      mseg = n
      go to 200
   end if
   continue
100 continue
200 continue

--- construct record name & write data
call GMCORH ( rname, 'SEGDEF ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', segdef, mseg,
              0, 0, 0, 1100 )
call GMCORH ( rname, 'NUMEL ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', numel, mseg,
              0, 0, 0, 1200 )
call GMCORH ( rname, 'XBEG ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', xbeg, mseg,
              0, 0, 0, 1300 )
call GMCORH ( rname, 'YBEG ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', ybeg, mseg,
              0, 0, 0, 1400 )
call GMCORH ( rname, 'XEND ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', xend, mseg,
              0, 0, 0, 1500 )
call GMCORH ( rname, 'YEND ', 1, mseg )
call GMPUTH ( 'W', idi, idsn, rname, 'I', yend, mseg,
              0, 0, 0, 1600 )
call GMFLUB ( idi, 0, 2000 )
else
   print *, '*** Illegal keyword ', key.' in segment data'
   end if
   go to 1000
end

The structure of this subroutine is typical of those that handle subordinate commands. A "do forever" construction is headed by a CLREAD call, and the loop is escaped only when an END command is detected. Notice the different prompt and verbose prompt input arguments.

This subroutine provides an example of the use of the "search for keyword" function ICLSEK described in [1], §5.2. A keyword match is followed by a value pair retrieval through the list-loading subroutine CLVALF described in [1], §7.2.
Defining Segments

Note the careful handling of the case in which less than two values appear after either BEGIN or END. This facilitates table editing. For example, the command

S=3 B=45.2

resets XBEG(3) to 45.2; nothing else changes.

To load the data the outline presented in §7.1.2 is followed exactly. Note that, the variable n is set to MAXSEG before every call to GMGETN. This insures that no more than MAXSEG values are read into the arrays that are dimensioned to MAXSEG. The value of n is reset after each call to GMGETN because n returns the actual number of values read. A real production Processor would perform some data checking to make sure the number of values read for each record are the same. Also, real professionals would use LMERCD [2], §14.5, to check for various errors that may have occurred during the read.

To store the data that has been entered the outline presented in §7.1.1 is followed exactly. Note that, the high cycle for the records is determined by computing the largest index of segdef that contains a non-zero value. Thus, when this data is read later the number of segments defined is known from the high cycle number for the SEGMENT dataset records.

7.2.1 Digression on Subordinate Commands

Why have we used subordinate commands rather than making the user type the segment in the DEFINE command itself? Well, contrast the above definition of the square region with the following one:

DEFINE SEGMENT=1 BEGIN=0.0 END=4.0
DEFINE SEGMENT=2 BEGIN=4.0 END=4.4
DEFINE SEGMENT=3 BEGIN=4.4 END=0.4
DEFINE SEGMENT=4 BEGIN=0.4 END=0.0

This is not too different in terms of typing effort, so the decision for adopting a one-level and a two-level structure in terms of number of keystrokes is marginal. But note that by going to a two-level scheme we have effectively separated the action of selecting what to define, namely segments, from the actual definition by entering coordinate values. This is a key aspect of object-oriented programming: first select, then operate. Let us make this a command design principle:

Try to separate selection from operation

If you are entering commands from a keyboard perhaps the advantages are not immediately apparent. But if you go to some form of interactive graphics input the advantages will be evident when you try to "cover" the commands through message-sending techniques. The user of such graphic system will then see SEGMENTS in a “model definition” menu, and by pointing to it he or she is transported to another screen or window in which the process of entering the segments is actually carried out.
7.3 Defining Elements

By default, each segment contains only one boundary element (see logic of DEFINE SEGMENT). To put more elements per segment you use the DEFINE ELEMENTS command. This introduces subordinate commands of the form

\[
\text{SEGMENT} = i \quad \text{ELEMENTS} = \{n\}
\]

[LOAD | STORE]

where \(n\) is the number of boundary elements in the \(i^{th}\) segment. The data is terminated by an END command. For the square region used as an example, let's say we want 10 BEs on segments 1 and 3, 15 BEs on segments 2 and 4, and store this data:

```
DEFINE ELEMENTS
SEG=1 EL=10 ; SEG=3 EL=10 ; SEG=2 EL=15 ; SEG=4 EL=15
STORE ; END
```

which illustrates the fact that data may be entered in any order. The implementation shown below actually allows a more general command form:

\[
\text{SEGMENTS} = i_1, \ldots, i_k \quad \text{ELEMENTS} = n_1, \ldots, n_k
\]

so that segment \(i_1\) gets \(n_1\) elements, segment \(i_2\) gets \(n_2\), and so on. The example above can be abbreviated to

```
DEFINE ELEMENTS
SEG=1:4 EL=10,15,10,15
STORE ; END
```

For this simple Processor, using a command like this is probably overkill. It is implemented in that fashion only to illustrate the processing of variable length integer lists via CLVALI

7.2: Define number of (equally spaced) boundary elements per segment

```
subroutine DEFINE_ELEMENTS

implicit none
include 'database.inc'
include 'segment.inc'

character*4 key, CCLVAL
character*20 rname
integer i, iseg, n, nseg
integer iseglist(MAXSEG), numelist(NAXSEG)
integer ICLTYP, ICLSEK
integer idsn, LNFIJID, ilow, ihigh, nrec
real FCLVAL
logical CMATCH

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```
Defining the Problem

1000 call CLREAD (' Element data> ',
$ ' Enter SEG = i1 ... ik EL = nel. ... nek&
$ ' or LOAD or STORE&
$ ' Terminate with END')

if (ICLTYP(1).le.0) then
  print *, '*** Command must begin with keyword'
goto 1000
end if
key = CCLVAL(1)
if (CMATCH (key, 'END')) then
  return
else if (CMATCH (key, 'SEG')) then
  call CLVALI (' ', -MAXSEG, iseglist, nseg)
  if (ICLSEG(0,'ELM').eq.0) then
    print *, '*** Keyword ELEMENTS is missing'
goto 1000
  end if
  call CLVALI (' ', -MAXSEG, numelist, n)
do 2500 i = 1,nseg
  iseg = iseglist(i)
  if (iseg .le. 0.or. iseg.gt. MAXSEG) then
    print *, '*** Segment number',iseg,' out of range'
  else
    numel(iseg) = max(numelist(i),1)
  end if
  continue
2500 continue
else if (CMATCH (key, 'LOAD')) then
  find dataset
  idsn = LMFIND ( idi, 'SEGMENT ', 100)
  if (ism.eq.0) then
    print *, 'Cannot find SEGMENT dataset; nothing LOADed.'
goto 1000
  end if

C --- get record name cycles, construct record name & read data
n = MAXSEG
  call GMGECY (' ', ldi, idsn, 'NUMEL ', nrec, ilow.
$. ihigh, 200 )
  call GMCORH ( rname, 'NUMEL ', ilow, ihigh )
  call GMGETH ( 'R/L', ldi, idsn, rname, 'I', numel. 
$. n, 0, 0, 0, 300 )
else if (CMATCH (key, 'STORE')) then

C --- find dataset because this is an update of SEGMENT_NUMEL
  idsn = LMFIND ( ldi, 'SEGMENT ', 500)
  if (ism.eq.0) then
    print*, 'Cannot find SEGMENT dataset; nothing STOREd.'
    print*, 'Must DEFINE SEGMENTS before DEFINing ELEMENTs.'
  end if
Defining the Problem

---

Define record name & write data

C

```c
    go to 1000
end if

---

call GMGECY ( ' ', ldi, idsn, 'NUMEL ', nrec, ilow, ihigh, 200 )

call GMCORN ( rname, 'NUMEL ', 1, nrec )
call GMPUTN ( 'W/U', ldi, idsn, rname, 'I', numel, nrec, 0, 0, 0, 1200 )
call GMFLUB ( ldi, 0, 2000 )
else
  print *, '*** Illegal keyword ', key, ' in element data'
end if

go to 1000
```

---

Here the data are loaded following the outline given in §7.1.2. However, note that, the NUMEL data are also loaded under DEFINE SEGMENTS, so it is not necessary to reload the data here. As stated above the DEFINE ELEMEI TS command is an overkill, so we end up with this strange construction. The first author recommends including the number of element definitions under DEFINE SEGMENTS as a subordinate command that calls DEFINE ELEMENTS, then all segment data operations are encapsulated in the same place. See §7.3.1 below for the second author's opinion.

Thus, to store the data here we must be sure the SEGMENTS dataset exists instead of the usual install operation (the dataset is installed in the DEFINE SEGMENTS code, §7.2). If the dataset is found we proceed, but not along the standard path. First, since the SEGMENT dataset and the NUMEL record already exist we retrieve the record cycles by calling GMGECY [2], §9.3. Then, the record name is constructed with a call to GMCORN [2], §10.6. Finally, note that, the op_code, the first argument, in GMPUTN is set to write/update. That is, we write over the existing data. See [2], §9.9 for more information on the op_codes.

If you can't follow the code, don't worry. It is more advanced than the typical input routine in DBEM2, so you can study it later.

---

7.3.1 Digression: Simplifying Commands

Why didn't we allow element data to be specified in the same commands that define the segment geometry? For example, we might have allowed commands such as

```
    SEG = 13  BEG = -1.50, 3.53  END = 14.81, 6.22  ELEM = 5
```

The answer fits within another design principle:

---

Keep commands simple

---

Simplicity is an admirable general principle, but for our case something more specific applies:
The terms "persistent" and "volatile" are used in a relative sense to denote degrees of "changeability" of the data. For example, segment data are more persistent than element data, since presumably you want to solve a problem whose geometry is dictated by external requirements; typically by engineering considerations. On the other hand, the number of elements per segment is a judgement decision: the program user attempts to get satisfactory accuracy (more elements, more accuracy) with reasonable cost (more elements, more computer time).

Frequently the number of elements is varied while keeping the segment data fixed; this is called a convergence study. So there are good reasons to separate the commands that define these two aspects.
### 7.4 Defining Boundary Conditions

Each segment may be given a different boundary condition (BC) that involves any of the following stress/displacement combinations:

<table>
<thead>
<tr>
<th>BC Code</th>
<th>Prescribed boundary values</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>Shear stress $\sigma_s$ and normal stress $\sigma_n$</td>
</tr>
<tr>
<td>1</td>
<td>Shear displacement $u_s$ and normal displacement $u_n$</td>
</tr>
<tr>
<td>2</td>
<td>Shear displacement $u_s$ and normal stress $\sigma_n$</td>
</tr>
<tr>
<td>3</td>
<td>Shear stress $\sigma_s$ and normal displacement $u_n$</td>
</tr>
</tbody>
</table>

These values are constant along the segment, so they can be read on a segment-by-segment basis. The stress values are understood to be resultants over the segment.

(The "BC codes" are related to those used by Crouch and Starfield [4]. Using integer codes is far from the best way to implement readable software, but we shall follow their convention.)

The BC data commands are introduced by a `DEFINE BOUNDARY_CONDITIONS` header command (which may be abbreviated to just D B), and have the form

```
SEG = i  {SS = $\sigma_s$  | SD = $u_s$}  {NS = $\sigma_n$  | ND = $u_n$}
[ LOAD  | STORE ]
```

terminated by an `END` command. Keyword SS means shear stress, SD shear displacement, and so on.

In the CLAMP metalanguage, the | indicates that one may specify either $\sigma_s$ or $u_s$, but not both simultaneously, and similarly for $\sigma_n$ and $u_n$. The specifications are shown in braces, meaning that they may not be omitted.

If no BC is ever specified for segment $i$, that segment is assumed stress free (code 0 with $\sigma_s = \sigma_n = 0$). If only a normal value is prescribed, a zero shear stress is assumed, and so on.

The implementation of `DEFINE_BOUNDARY` follows.

---

* Read boundary condition data for segments
  
  subroutine DEFINE_BOUNDARY_CONDITIONS

  implicit none
  include 'database.inc'
  include 'segment.inc'

  character*4 key, CCLVAL, word(2)
  character*20 rname
  integer iseg, n, nw, iloc(2)
  integer ICLVAL, ICLSEK, ICLTYP
  integer idsn, LMFIND, ilow, ihigh, nrec, mseg

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logical CMATCH

1000 call CLOADK ('L', -2, word, iloc, nw)
call BCVALUES (iseg, nw, word, iloc)
else if (CMATCH (key, 'L0AD')) then

C --- find dataset
idsn = LMFIND ( idi, 'BCVALUES ', 100 )
if ( idsn .eq. 0 ) then
  print*, ' Cannot find BCVALUES dataset; nothing LOADed.'
go to 1000
end if

C --- get record name cycles, construct record name & read data
n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'KODE ', nrec, ilow,
  ihigh, 220 )
call GMCORN ( rname, 'KODE ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'I', kode,
  n, 0, 0, 0, 200 )

n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'BVN ', nrec, ilow,
  ihigh, 220 )
call GMCORN ( rname, 'BVN ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', bvn,
  n, 0, 0, 0, 300 )

n = MAXSEG
call GMGECY ( ' ', idi, idsn, 'BVS ', nrec, ilow,
  ihigh, 220 )
call GMCORN ( rname, 'BVS ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', bvs,
  n, 0, 0, 0, 400 )
else if (CMATCH (key, 'S\^TORE')) then

C --- install dataset
    call GMPUTN ( idi, 'BCVALUES ', idsn, 16, 1000 )
C --- determine largest index of data stored
    do 100 n=MAXSEG,1,-1
        if ((kode(n) .ne. 0) .or. (bvn(n) .ne. 0.0)
            .or. (bvs(n) .ne. 0.0)) then
            mseg = n
            go to 200
        end if
    end do
100 continue
200 continue
C --- construct record name & write data
    call GMCORN ( rname, 'KODE ', I, mseg )
    call GMPUTN ( 'W', idi, idsn, rname, 'I', kode, mseg,
                   0, 0, 0, 1100 )
    call GMCORN ( rname, 'BVN ', I, mseg )
    call GMPUTII ( 'W', idi, idsn, rname, 'S', bvn, mseg,
                   0, 0, 0, 1200 )
    call GMCORN ( rname, 'BVS ', I, mseg )
    call GMPUTN ( 'W', idi, idsn, rname, 'S', bvs, mseg,
                   0, 0, 0, 1300 )
    call GMFLUB ( idi, 0, 2000 )
else
    print *. '*:_* Illeagal keyword ', key, ' in BC data'
end if
end if

This subroutine follows the outline for LOAD and STORE given in §7.1.2 and §7.1.1. This also illustrates the use of the "load keyword" entry points of [1], §8.2. These calls search for keywords such as SS and move them to the subroutine work area. This simplifies keyword legality tests such as "SS and SD cannot appear in the same command." To do these chores DEFINE_BOUNDARY calls subroutine BCVALUES:

* Store boundary condition calls in tables
  subroutine BCVALUES
  $(iseg, nw, word, iloc)
  *
  implicit none
  include 'segment.inc'
  character*(*) word(2)
  real FCLVAL
  integer iseg, nw, iloc(2)
  integer code, i, isd, iloads, iloadn, ks, kd, kn
  logical CMATCH

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Defining Boundary Conditions

ks = 0
kn = 0
kd = 0
isd = 0
iloadn = 0
iloads = 0

* do 2000 i = 1,nw
   if (CMATCH (word(i), 'SS')) then
      ks = ks + 1
      iloads = iloc(i)
   else if (CMATCH (word(i), 'SD')) then
      ks = ks + 1
      kd = kd + 1
      isd = 1
      iloads = iloc(i)
   else if (CMATCH (word(i), 'NS')) then
      kn = kn + 1
      iloadn = iloc(i)
   else if (CMATCH (word(i), 'ND')) then
      kn = kn + 1
      kd = kd + 1
      iloadn = iloc(i)
   else
      print *, '*** Illegal BC keyword ', word(i), ' segment', iseg
      return
   end if
   if (kn .gt. 1 .or. ks .gt. 1) then
      print *, '*** Illegal BC combination for segment', iseg
      return
   end if
   continue

   if (iloadn .gt. 0) then
      bvn(iseg) = FCLVAL(iloadn+1)
   end if
   if (iloads .gt. 0) then
      bvs(iseg) = FCLVAL(iloads+1)
   end if
   if (kd .eq. 0) then
      code = 1
   else if (kd .eq. 1) then
      code = 3
      if (isd .eq. 0) code = 4
   else
      code = 2
   end if
   kode(iseg) = code-1
   return
end
which embodies the logic for eventually storing the user-supplied values into appropriate spots in arrays BVS and BVII.
7.5 Defining Symmetry Conditions

If the problem exhibits symmetry conditions, commands to specify symmetry axes are introduced by the header command DEFINE SYMMETRY CONDITIONS (which may be abbreviated to just D S) and have the form

\[
\begin{align*}
\text{XSYM} &= x_{sym} \\
\text{YSYM} &= y_{sym} \\
\text{[ LOAD | STORE ]}
\end{align*}
\]

terminated by an END command. The XSYM command specifies that \( x = x_{sym} \) is a line of symmetry parallel to the \( x \) axis. The YSYM command specifies that \( y = y_{sym} \) is a line of symmetry parallel to the \( y \) axis. One or two specifications may be given. The Processor logic does not allow "skew" symmetry conditions.

The implementation of the DEFINE SYMMETRY routine is straightforward:

Read symmetry condition data

subroutine DEFINE_SYMMETRY_CONDITIONS

implicit none
include 'database.inc'
include 'symmetry.inc'
character*4 key, CCLVAL
integer ixsym, iysym, FCLVAL
logical CMATCH

ixsym = mod(ksym,2)
iysym = ksym/2

1000 call CLREAD (' Symmetry data> ',
$  ' Enter XSYM=xsym or YSYM=ysym or LOAD or STORE '/
$  ' &&Terminate with END')
if (ICLTYP(1) .le. 0) then
print *, '*** Command must begin with keyword'
go to 1000
end if

key = CCLVAL(1)
if (CMATCH (key, 'E^IID')) then
return
else if (CMATCH (key, 'X^SYM')) then
ixsym = FCLVAL(2)
ixsym = 1
ksym = 2*iysym + ixsym
else if (CMATCH (key, 'Y^SYM')) then
Defining the Problem

Defining Symmetry Conditions

ysym = FCLVAL(2)
iysym = 1
ksym = 2*iysym + ixsym

else if (CMATCH (key, 'LOAD')) then
    --- find dataset
    idsn = LMFIND ( idi, 'SYMMETRY ', 100 )
    if ( idsn .eq. 0 ) then
        print*, 'Cannot find SYMMETRY dataset; nothing LOADed.'
        go to 1000
    end if

--- read data
n = 1
    call GMGETN ('R/L', idi, idsn, 'KSYM ', 'I', ksym,
        n, 0, 0, 0, 200 )
    n = 1
    call GMGETN ('R/L', idi, idsn, 'XSYM ', 'I', xsym, n, 0, 0, 0, 300)
    n = 1
    call GMGETN ('R/L', idi, idsn, 'YSYM ', 'I', ysym, n, 0, 0, 0, 400)
else if (CMATCH (key, 'STORE')) then
    --- install dataset
    call GMPUNT ( idi, 'SYMMETRY ', idsn, 16, 500 )

--- write data
    call GMPUTH ( 'W', idi, idsn, 'KSYM ', 'I', ksym, 1, 0, 0, 0, 600 )
    call GMPUTH ( 'W', idi, idsn, 'XSYM ', 'I', xsym, 1, 0, 0, 700 )
    call GMPUTH ( 'W', idi, idsn, 'YSYM ', 'I', ysym, 1, 0, 0, 800 )
    call GMFLUB ( idi, 0, 2000 )
else
    print *, '*** Illegal keyword ', key, ' in symmetry data'
end if
    go to 1000
end

(Here KSYM is an integer "symmetry flag" related to that used in the original TWOBI program.)

The LOAD and STORE commands are implemented following the outline given in §7.1.2 and §7.1.1. However, here each record is just one number, so we do not have to construct the record name. For more information see [2], §5.1 for the details of record naming.
7.6 Defining Material Properties

Material properties are introduced by a DEFINE MATERIAL header command (which can be abbreviated to just D M). The commands have a simple form:

\[
\begin{align*}
\text{EM} &= E \\
\text{PR} &= \nu \\
\text{[ LOAD | STORE ]}
\end{align*}
\]

terminated by an END command. The EM command specifies the elastic modulus and the PR command specifies Poisson's ratio. Since DBEM2 is restricted to elastic isotropic materials and does not consider thermal effects, these two material properties suffice.

The default values for \( E \) and \( \nu \) set by CLEAR are 1.0 and 0.0, respectively.

The implementation of DEFINE MATERIALS is straightforward and does not involve any fancy new construct:

* Read material property data
* subroutine DEFINE_MATERIAL
*
implicit none
include 'database.inc'
include 'material.inc'
character*4 key, CCLVAL
integer ICLTYP, idsn, LMFIND, n
real FCLVAL
logical CMATCH

1000 call CLREAD ('Material data>', 'Enter EM=em or PR=pr or LOAD or STORE&&'//
$ 'Terminate with END')

if (ICLTYP(1) .le. 0) then
print *, '*** Command must begin with keyword'
go to 1000
end if
key = CCLVAL(1)
if (CMATCH (key, 'E=ND')) then return
else if (CMATCH (key, 'EM')) then
em = FCLVAL(0)
else if (CMATCH (key, 'PR')) then
pr = FCLVAL(0)
else if (CMATCH (key, 'LOAD')) then
--- find dataset
idsn = LMFIND ( idi, 'MATERIAL ', 100 )
if ( idsn .eq. 0 ) then

print*, 'Cannot find MATERIAL dataset; nothing LOADED'
go to 1000
end if

--- read data
n = 1
call GMGETN ( 'R/L', idi, idsn, 'EM', 'S', em, n, 0, 0, 0, 200 )
$ n = 1
call GMGETN ( 'R/L', idi, idsn, 'PR', 'S', pr, n, 0, 0, 0, 300 )
else if (CMATCH (key, 'STORE')) then
--- install dataset
call GMPUNT ( idi, 'MATERIAL', idsn, 16, 500 )
--- write data
call GMPUTN ( 'W', idi, idsn, 'EM', 'S', em, 1, 0, 0, 0, 600 )
call GMPUTN ( 'W', idi, idsn, 'PR', 'S', pr, 1, 0, 0, 0, 700 )
call GMFLUB ( idi, 0, 2000 )
else
print*, '*** Illegal keyword .key.' in material data'
end if
end go to 1000
end

Here the LOAD and STORE are identical to the previous implementation.
7.7 Defining Prestress Data

If the initial stress state has nonzero components, prestress data have to be introduced by a DEFINE PRESTRESS header. The prestress-definition commands have a very simple form:

\[
\begin{align*}
S_{xx} &= \sigma_{xx} \\
S_{yy} &= \sigma_{yy} \\
S_{xy} &= \sigma_{xy}
\end{align*}
\]

As usual, these commands are terminated by an END command. Undefined prestress components are assumed zero.

The implementation of DEFINE PRESTRESS is quite similar to that of DEFINE_MATERIAL:

* Read prestress (initial field stresses) data *

```
subroutine DEFINE_PRESTRESS

implicit none
include 'database.inc'
include 'prestress.inc'
character*4 key, CCLVAL
integer ICLTYP, idsn, LMFIND, n
real FCLVAL
logical CMATCH

1000 call CLREAD (' Prestress data> ',
    $ ' Enter SXX0=sxx0, SYY0=syy0 or SXY0=sxy0&&'
    $ ' or LOAD or STORE&&'
    $ 'Terminate with END')

    if (ICLTYP(1) .le. 0) then
        print *, '*** Command must begin with keyword'
go to 1000
    end if
    key = CCLVAL(1)
    if (CMATCH (key, 'E=ND')) then
        return
    else if (CMATCH (key, 'SX=0')) then
        sxx0 = FCLVAL(0)
    else if (CMATCH (key, 'SY=0')) then
        syy0 = FCLVAL(0)
    else if (CMATCH (key, 'SX=O')) then
        sxy0 = FCLVAL(0)
    else if (CMATCH (key, 'L=OAD')) then
        idsn = LMFIND ( ldi, 'PRESTRESS ', 100 )
```

---
if ( idsn .eq. 0 ) then
  print*, ' Cannot find PRESTRESS dataset; nothing LOADed.'
  go to 1000
end if

--- read data
n = 1
  call GMGETH ( 'R/L', ldi, idsn, 'SXXO', 'S', sxx0,
                 n, 0, 0, 0, 200 )
n = 1
  call GMGETH ( 'R/L', ldi, idsn, 'SYYO', 'S', syy0,
                 n, 0, 0, 0, 300 )
n = 1
  call GMGETH ( 'R/L', ldi, idsn, 'SXYO', 'S', sxy0,
                 n, 0, 0, 0, 400 )
else if (CMATCH (key, 'STORE')) then
  --- install dataset
  call GMPUTN ( 'W', ldi, idsn, 'SXXO', 'S', sxx0, I,
                 O, O, O, 600 )
  call GMPUTN ( 'W', ldi, idsn, 'SYYO', 'S', syy0, I,
                 O, O, O, 700 )
  call GMPUTFI ( 'W', ldi, idsn, 'SXYO', 'S', sxy0, I,
                 O, O, O, 800 )
  call GMFLUB ( ldi, 0, 2000 )
else
  print *, '*** Illegal keyword ', key, ' in prestress data'
end if
  go to 1000
end
7.8 Defining Output Field Locations

The last piece of input data is not related to the problem definition, but to the specification of the field points at which the program user would like to get computed results, viz., displacements and stresses.

(This set of information is characteristic of boundary element methods, in which all basic givens and unknowns are at the boundary. If you want information at field points not on the boundary, you have to ask for it and specify where.)

For convenience the output locations are not specified point by point, but as equally spaced points on line segments. You specify the location of the first and last point on the line, and the number of points, if any, to be “collocated” between the first and last one.

The output field location specification commands are introduced by a DEFINE FIELD-LOCATIONS header command (which may be abbreviated to D F) and have a form reminiscent of the segment-definition commands:

\[
\text{LINE} = i \quad \text{FIRST} = x_{i_{\text{first}}}, y_{j_{\text{first}}} \quad \text{LAST} = x_{i_{\text{last}}}, y_{j_{\text{last}}} \quad \text{[POINTS=nn_n]}]
\]

terminated by an END command. Here \( nn_n \) is the number of intermediate points to be inserted (equally spaced) between the first and last point. If this phrase is omitted, \( nn_n = 0 \) is assumed so only the first and last points will be output points. If the first and last points coincide, output will be at only one point.

For example:

\[
\begin{align*}
\text{DEF OUT} \\
\text{LINE=1 F=200.2 L=203.8 P=9} \\
\text{LINE=2 F=3.8,0.2 L=0.2,3.8 P=9} \\
\text{STORE} \\
\text{END}
\end{align*}
\]

specifies two output lines running at 45° and 135°, respectively, with 11 output points (first + last + 9) in each, and the data are stored.

Here is the implementation of the DEFINE OUTPUT LOCATIONS routine:

```
Read location of output field points

subroutine DEFINE_FIELD_LOCATIONS

implicit none
include 'database.inc'
include 'output.inc'
character*8 key, CCLVAL
character*20 rname
real FCLVAL
integer ilin, n, mark, ICLVAL, ICLSEK, ICLTYP
```
Defining the Problem

Defining Output Field Locations

integer idsn, LMFIND, ilow, ihigh, nrec, mlin
real xy(2)
logical onepoint, CMATCH

1000 call CLREAD (' Field location data> ',
$ ' Enter LIN=ilin FIRST=xfirst,yfirst LAST=xlast,ylast'//
$ '[P=ninter]&& or LOAD or STORE&& Terminate with END')

if (ICLTYP(1) .le. 0) then
print *, '*** Command must begin with keyword'
go to 1000
end if
key = CCLVAL(1)
if (CMATCH (key, 'END')) then
return
else if (CMATCH (key, 'LINE')) then
ilin = ICLVAL(2)
if (ilin .le. 0 .or. ilin .gt. MAXLIN) then
print *, '*** Field line number',ilin,' is out of range'
go to 1000
end if
lindef(ilin) = 1
nintop(ilin) = 0
onepoint = .true.
if (ICLSEK(3, 'FIRST') .ne. 0) then
call CLVALF (' ', 2, xy, n)
if (n .ge. 1) xfirst(ilin) = xy(1)
if (n .ge. 2) yfirst(ilin) = xy(2)
end if
if (ICLSEK(3, 'LAST') .ne. 0) then
call CLVALF (' ', 2, xy, n)
if (n .ge. 1) xlast(ilin) = xy(1)
if (n .ge. 2) ylast(ilin) = xy(2)
onepoint = .false.
end if
if (onepoint) then
xlast(ilin) = xfirst(ilin)
ylast(ilin) = yfirst(ilin)
end if
if (ICLSEK(3, 'POINTS') .ne. 0) then
nintop(ilin) = max(ICLVAL(0),0)
end if
else if (CMATCH (key, 'LOAD')) then
-- find dataset
idsn = LMFIND ( idi, 'FIELD ', 100 )
if ( idsn .eq. 0 ) then
print *, ' Cannot find FIELD dataset: nothing LOADed.'
go to 1000

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end if

C --- get record name cycles, construct record name & read data

n = MAXLIN
call GMGECY ( ' ', idi, idsn, 'LINDEF ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'LINDEF ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'I', lindef,  
            n, 0, 0, 0, 200 )

n = MAXLIN
call GMGE CY ( ' ', idi, idsn, 'LINTOP ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'LINTOP ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'I', nintop,  
            n, 0, 0, 0, 300 )

n = MAXLIN
call GMGE CY ( ' ', idi, idsn, 'XFIRST ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'XFIRST ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', xfirst,  
            n, 0, 0, 0, 400 )

n = MAXLIN
call GMGE CY ( ' ', idi, idsn, 'YFIRST ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'YFIRST ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', yfirst,  
            n, 0, 0, 0, 500 )

n = MAXLIN
call GMGE CY ( ' ', idi, idsn, 'XLAST ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'XLAST ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', xlast,  
            n, 0, 0, 0, 600 )

n = MAXLIN
call GMGE CY ( ' ', idi, idsn, 'YLAST ', nrec, ilow,  
             ihigh, 220 )
call GMCORN ( rname, 'YLAST ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', ylast,  
            n, 0, 0, 0, 700 )

else if ( CMATCH (key, 'S^STORE')) then

C --- install dataset

C --- determine largest value of lindef

do 100 n=MAXLIN,1,-1
   if ( lindef(n) .ne. 0 ) then
      mlin = n
   go to 200
end if
100 continue
Defining the Problem

200     continue
C      --- construct record name & write data
     call GMCORN ( rname, 'LINDEF ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'I', lindef, mlin, 0, 0, 0, 1100 )
     call GMCORN ( rname, 'HINTOP ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'I', nintop, mlin, 0, 0, 0, 1200 )
     call GMCORN ( rname, 'XFIRST ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'S', xfirst, mlin, 0, 0, 0, 1300 )
     call GMCORN ( rname, 'YFIRST ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'S', yfirst, mlin, 0, 0, 0, 1400 )
     call GMCORN ( rname, 'XLAST ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'S', xlast, mlin, 0, 0, 0, 1500 )
     call GMCORN ( rname, 'YLAST ', 1, mlin )
     call GMPUTH ( 'W', ldi, idsn, rname, 'S', ylast, mlin, 0, 0, 0, 1600 )
else
     print *, '*** Illegal keyword ', key, ' in field loc data'
   end if
   go to 1000
end

Once again, the LOAD and STORE command implementations follow the outline given in §7.1.2 and §7.1.1.

The input data section is complete.
Defining Output Field Locations

Defining the Problem

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8. Solving the Problem

Having finished input data preparation, the three steps involved in solving the elasto-static problem are as follows.

Building the Boundary Element Model. The input data have defined the geometry of the problem in terms of segments. Segments are broken down into equally spaced boundary elements. The first step consists of building element-by-element data, and is carried out when you enter the command BUILD.

Assembling the Discrete Equations. This step generates a matrix \( C \) of “influence coefficients” and a vector \( r \) of “forcing functions.” These arrays have dimensions equal to twice the total number of boundary elements. The construction of the elements of \( C \) and \( r \) follows the direct formulation of boundary-integral methods and is not explained here. This step is triggered by the command GENERATE and is carried out by subroutine GENERATE and subordinate routines.

Solving for the unknowns. The linear equation system \( Cx - r \) is solved (by a Gauss elimination method) for vector \( x \), which contains the boundary unknowns. This step is triggered by command SOLVE and is carried out by subroutine SOLVE and a subordinate routine.

For each of these commands, BUILD, GENERATE, and SOLVE the qualifiers LOAD and STORE can be used. If the qualifier is LOAD then the data are loaded from the open GAL Library without computing the data. If the qualifier is STORE the data are computed, then stored in the GAL Library. The presence of the qualifier is determined in the subroutine DO_COMMAND; see §5. The value of the qualifier is passed to BUILD, GENERATE, and SOLVE as the argument \( op \), a blank value is the default.

All of the GAL entry points and methods for loading and storing data have been illustrated previously in §7.0, and since we are not going to explain the theory behind these tasks, the BUILD, GENERATE and SOLVE subroutines are listed next without commentary.

```fortran
* *
* Build detailed boundary element data
* *
subroutine BUILD ( op )
*
implicit none
include 'database.inc'
include 'segment.inc'
include 'element.inc'
include 'material.inc'
include 'prestress.inc'
character(*) op
character*20 rname
integer iseg, k, ne, num
integer idsn, LMFINID, ilow, ihigh, nrec, n
real xd, yd, side
```

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if ( op .eq. 'LOAD') go to 5000
k = 0
do 2000 iseg = 1,MAXSEG
   if (segdef(iseg) .eq. 0) go to 2000
   num = numel(iseg)
   xd = (xend(iseg)-xbeg(iseg))/num
   yd = (yend(iseg)-ybeg(iseg))/num
   side = sqrt(xd**2+yd**2)
   if (side .eq. 0.0) go to 2000
   do 1500 ne = l,num
      k= k+l
      if (k .gt. MAXELM) then
         print *, '*:_ Boundary element count exceeds ',MAXELM
         print *, ' Excess elements ignored'
         return
      end if
      xme(k) = xbeg(iseg) + 0.5*(2.*ne-l)*xd
      yme(k) = ybeg(iseg) + 0.5*(2.*ne-l)*yd
      hleng(k) = 0.5*side
      sinbet(k) = yd/side
      cosbet(k) = xd/side
      b(2*k-1) = bvs(iseg)
      b(2*k ) = bvn(iseg)
      kod(k) = kode(iseg)
   1500 continue
2000 continue
numbe = k
   print '('' Discrete model building completed:'',
      I5,''' boundary elements''/'), numbe
   if ( op .eq. 'STORE') then
      --- STORE data
      --- install dataset
      call GMPUNT ( idi, 'ELEMENT ', idsn, 16, 1000 )
      C --- construct record name & write data
      call GMPUTH ('W', idi, idsn, 'NUMBE ', 'I', numbe,
         1, 0, 0, 0, 1050 )
      call GMCORN ( rname, 'XME ', 1, numbe )
      call GMPUTH ('W', idi, idsn, rname, 'S', xme, numbe,
         0, 0, 0, 1100 )
      call GMCORN ( rname, 'YME ', 1, numbe )
      call GMPUTH ('W', idi, idsn, rname, 'S', yme, numbe,
         0, 0, 0, 1200 )
      call GMCORN ( rname, 'HLEIIG ', 1, numbe )
      call GMPUTH ('W', idi, idsn, rname, 'S', hleng, numbe,
         0, 0, 0, 1300 )
      call GMCORN ( rname, 'SINBET ', 1, numbe )
      call GMPUTH ('W', idi, idsn, rname, 'S', sinbet, numbe,
         0, 0, 0, 1400 )
Solving the Problem

call GMCORH ( rname, 'COSBET ', 1, numbe )
call GMPUTN ( 'W', idi, idsn, rname, 'S', cosbet, numbe, $
0, 0, 0, 1500 )
call GMCORH ( rname, 'KOD ', 1, numbe )
call GMPUTN ( 'W', idi, idsn, rname, 'I', kod, numbe, $
0, 0, 0, 1600 )
call GMCORH ( rname, 'B ', 1, 2*numbe )
call GMPUTN ( 'W', idi, idsn, rname, 'S', b, 2*numbe, $
0, 0, 0, 1500 )
call GMFLUB ( idi, 0, 2000 )
end if
return

5000 continue
C --- LOAD data
C --- find dataset
idsn = LMFIND ( idi, 'ELEMENT ', 100 )
if ( idsn .eq. 0 ) then
   print*, 'Cannot find ELEMENT dataset; nothing LOADed.'
   return
end if

C --- get record name cycles, construct record name & read data
n = 1
call GMGETN ( 'R/L', idi, idsn, 'NUMBE ', 'I', numbe,
$ n, 0, 0, 0, 150 )
n = MAXELM
call GMGECY ( ' ', idi, idsn, 'XME ', nrec, ilow,
$ ihigh, 180 )
call GMCORH ( rname, 'XME ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', xme,
$ n, 0, 0, 0, 200 )
n = MAXELM
call GMGECY ( ' ', idi, idsn, 'YME ', nrec, ilow,
$ ihigh, 280 )
call GMCORH ( rname, 'YME ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', yme,
$ n, 0, 0, 0, 300 )
n = MAXELM
call GMGECY ( ' ', idi, idsn, 'HLENG ', nrec, ilow,
$ ihigh, 380 )
call GMCORH ( rname, 'HLENG ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', hleng,
$ n, 0, 0, 0, 400 )
n = MAXELM
call GMGECY ( ' ', idi, idsn, 'SINBET ', nrec, ilow,
$ ihigh, 480 )
call GMCORH ( rname, 'SINBET ', ilow, ihigh )
call GMGETN ( 'R/L', idi, idsn, rname, 'S', sinbet,
Solving the Problem

$n = \text{MAXELM}$
call GMGECY ( ' ', ldi, idsn, 'COSBET ', nrec, ilow, $ihih, 500 )
call GMCORN ( rname, 'COSBET ', ilow, ihih )
call GMGETN ( 'R/L', ldi, idsn, rname, 'S', cosbet,$n, 0, 0, 0, 600 )
$n = \text{MAXELM}$
call GMGECY ( ' ', ldi, idsn, 'KOD ', nrec, ilow, $ihih, 680 )
call GMCORN ( rname, 'KOD ', ilow, ihih )
call GMGETN ( 'R/L', ldi, idsn, rname, 'I', kod, $n, 0, 0, 0, 700 )
$n = \text{MAXEQS}$
call GMGECY ( ' ', ldi, idsn, 'B ', nrec, ilow, $ihih, 780 )
call GMCORN ( rname, 'B ', ilow, ihih )
call GMGETN ( 'R/L', ldi, idsn, rname, 'S', b,$n, 0, 0, 0, 800 )

return

end

Calculate influence coefficient matrix and RHS vector

subroutine GENERATE ( op )

implicit none
include 'database.inc'
include 'material.inc'
include 'element.inc'
include 'prestress.inc'
include 'symmetry.inc'
character(*) op
character*20 rname
integer i, j, n, nrec, ilow, ihih, LMFIN, idsn, igap
real sinbi, cosbi, sinbj, cosbj, sso, sn0, g
real xi, xj, yi, yj, sj
real ass, asn, ans, ann, bss, bsn, bns, bnn

if ( op .eq. 'LOAD' ) go to 4000
$ g = 0.5*em/(1.+pr) 
do 3000 i = 1,numbe
$r(2*i-1) = 0.$
$r(2*i ) = 0.$
$xi = xme(i)$
$yi = yme(i)$
cosbi = cosbet(i)

3000 continue
Solving the Problem

\[
\begin{align*}
\text{sinbi} & = \sin\theta(i) \\
do\ 2500 & \ j = 1, \text{numbe} \\
\text{ass} & = 0.0 \\
\text{asn} & = 0.0 \\
\text{ans} & = 0.0 \\
\text{ann} & = 0.0 \\
\text{bss} & = 0.0 \\
\text{bsn} & = 0.0 \\
\text{bns} & = 0.0 \\
\text{bnn} & = 0.0 \\
xj & = x_j \ (j) \\
yj & = y_m \ (j) \\
\cosbj & = \cos\theta(j) \\
\sinbj & = \sin\theta(j) \\
sj & = hlen(j) \\
ss0 & = (syy0 - sxx0) \sinbj \cosbj + sxy0(\cosbj^2 - \sinbj^2) \\
sn0 & = sxx0 \sinbj^2 - 2.sxy0.\sinbj.\cosbj + syy0.\cosbj^2 \\
call & \text{COEFF} (x_i, y_i, x_j, y_j, sj, \\
& \quad \em, \pr, \cosbi, \sinbi, \cosbj, \sinbj, \\
& \quad \ass, \asn, \ans, \ann, \bss, \bsn, \bns, \bnn) \\
& \text{if (ksym.eq.1.or.ksym.eq.3) then} \\
call & \text{COEFF} (x_i, y_i, 2.sxsym-xi, y_j, sj, \\
& \quad \em, \pr, \cosbi, \sinbi, \cosbj, \sinbj, \\
& \quad \ass, \asn, \ans, \ann, \bss, \bsn, \bns, \bnn) \\
& \text{end if} \\
& \text{if (ksym.eq.2.or.ksym.eq.3) then} \\
call & \text{COEFF} (x_i, y_i, 2.symsym-ymj, sj, \\
& \quad \em, \pr, \cosbi, \sinbi, \cosbj, \sinbj, \\
& \quad \ass, \asn, \ans, \ann, \bss, \bsn, \bns, \bnn) \\
& \text{end if} \\
& \text{if (ksym.eq.3) then} \\
call & \text{COEFF} (x_i, y_i, 2.sxsym-xi, 2.symsym-yj, sj, \\
& \quad \em, \pr, \cosbi, \sinbi, \cosbj, \sinbj, \\
& \quad \ass, \asn, \ans, \ann, \bss, \bsn, \bns, \bnn) \\
& \text{end if} \\
call & \text{SETUP}(i, j, kod(j), g, ss0, sn0, \\
& \quad \ass, \asn, \ans, \ann, \bss, \bsn, \bns, \bnn, \\
& \quad 2 \times \text{numbe}, \text{MAXEQS}) \\
2500 & \text{continue} \\
3000 & \text{continue} \\
\text{print} *, 'Influence coefficient matrix & RHS vector generated' \\
& \text{if (op.eq. 'STORE') then} \\
C & \text{--- STORE data} \\
C & \text{--- install dataset} \\
call & \text{GMPUNT}(\ldi, 'COEFF ', \idsn, 16, 100) \\
C & \text{--- construct record name & write data} \\
& n = 2 \times \text{numbe} \\
inegap & = \text{MAXEQS} - n
Solving the Problem

call GMCORN ( rname, 'C ', 1, n )
call GMPUTH ( 'W', ldi, idsn, rname, 'S', c, -n,
$ 0, igap, 0, 200 )

C --- install dataset
call GMPUNT ( ldi, 'RHS ', idsn, 16, 500 )

C --- construct record name & write data
n = 2*numbe
    call GMCORN ( rname, 'R ', 1, n )
call GMPUTH ( 'W', ldi, idsn, rname, 'S', r, n,
$ 0, 0, 0, 600 )
call GMFLUB ( ldi, 0, 2000 )
    return
end if

C --- LOAD data
4000 continue

C --- find dataset
idsn = LMFIND ( ldi, 'COEFF ', 1000 )
if ( idsn .eq. 0 ) then
    print *, ' Cannot find COEFF dataset; nothing LOADed.'
else

C --- get record name cycles, construct record name & read data
n = MAXEQS**2
    call GMGECY ( ' ', ldi, idsn, 'C ', nrec, ilow,
$ ihigh, 1100 )
igap = MAXEQS - nrec
    call GMCORN ( rname, 'C ', ilow, ihigh )
call GMGET ( 'R/L', ldi, idsn, rname, 'S', c,
$ n, 0, igap, 0, 1200 )
end if

C --- find dataset
idsn = LMFIND ( ldi, 'RHS ', 1500 )
if ( idsn .eq. 0 ) then
    print *, ' Cannot find RHS dataset; nothing LOADed.'
else

C --- get record name cycles, construct record name & read data
n = MAXEQS
    call GMGECY ( ' ', ldi, idsn, 'R ', nrec, ilow,
$ ihigh, 1600 )
call GMCORN ( rname, 'R ', ilow, ihigh )
call GMGET ( 'R/L', ldi, idsn, rname, 'S', r,
$ n, 0, 0, 0, 1700 )
end if

return
end

* Solve for unknown boundary values

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* subroutine SOLVE ( op ) *

implicit none
include 'database.inc'
include 'element.inc'
character*(*) op
character*20 rname
integer ising
integer n, nrec, ilow, ihigh, LMFIND, idsn

if ( op .eq. 'LOAD' ) go to 1000
call GAUSSER ( c, r, x, 2*numbe, MAXEQS, ising)
if (ising .eq. 0) then
    print *, 'Discrete equations solved'
    if ( op .eq. 'STORE') then
        --- STORE data
        --- install dataset
        call GMPUNT ( idi, 'SOLUTION', idsn, 16, 500 )
        --- construct record name & write data
        n = 2*numbe
        call GMCORN ( rname, 'X', 1, n )
        call GMPUTH ( 'W', idi, idsn, rname, 'S', x, n,
                      0, 0, 0, 600 )
        call GMFLUB ( idi, 0, 2000 )
    end if
    else
        print *, 'Singularity detected at BE equation', ising
    end if
    return

C --- LOAD data
1000 continue
C --- find dataset
idsn = LMFIND ( idi, 'SOLUTION', 1500 )
if ( idsn .eq. 0) then
    print*, 'Cannot find SOLUTION dataset; nothing LOADed.'
else
    --- get record name cycles, construct record name & read data
    n = MAXEQS
    call GMGECY ( ' ', idi, idsn, 'X ', nrec, ilow,
                      ihigh, 1600 )
    call GMCORN ( rname, 'X ', ilow, ihigh )
    call GMGETN ( 'R/L', idi, idsn, rname, 'S', x,
                      n, 0, 0, 0, 1700 )
    end if
    return
Subroutine GENERATE calls COEFF (which is essentially the same as a TWOBI subroutine with the same name) and SETUP, which fills the entries of the influence coefficient matrix and right-hand-side vector:

* Calculate source/receiver coefficients

subroutine COEFF
$$ \begin{align*}
& (x_i, y_i, x_j, y_j, a_j, \\
& \text{msym, em, pr, cosbi, sinbi, cosb, sinb,} \\
& \text{ass, asn, ans, ann, bss, bsn, bns, bnn})
\end{align*}$$

implicit none
real \( x_i, y_i, x_j, y_j, a_j \)
real \( \text{em, pr, cosbi, sinbi, cosb, sinb} \)
real \( \text{ass, asn, ans, ann, bss, bsn, bns, bnn} \)
real \( \pi, \text{con, pr1, pr2, pr3} \)
integer \( \text{msym} \)
real \( \text{cma, cpa, cxb, cyb, cosg, sing} \)
real \( \text{rls, r2s, f11, f12} \)
real \( \text{tb1, tb2, tb3, tb4, tb5} \)
real \( \text{asnt, anst, annt} \)
real \( \text{bsnt, bnst, bnnt} \)

\( \pi = 4. \times \text{atan2}(1., 1.) \)
\( \text{con} = 1.0 / (4. \times \pi \times (1. - \text{pr})) \)
\( \text{pr1} = 1. - 2. \times \text{pr} \)
\( \text{pr2} = 2. \times (1. - \text{pr}) \)
\( \text{pr3} = 3. - 4. \times \text{pr} \)
\( \text{cxb} = (x_i - x_j) \times \text{cosb} + (y_i - y_j) \times \text{sinb} \)
\( \text{cyb} = -(x_i - x_j) \times \text{sinb} + (y_i - y_j) \times \text{cosb} \)
\( \text{cosg} = \text{cosbi} \times \text{cosb} + \text{sinbi} \times \text{sinb} \)
\( \text{sing} = \text{sinbi} \times \text{cosb} - \text{cosbi} \times \text{sinb} \)

\( \text{cma} = \text{cxb} - a_j \)
\( \text{cpa} = \text{cxb} + a_j \)
\( \text{rls} = \text{cma} \times 2 \times \text{cyb} \times 2 \)
\( \text{r2s} = \text{cpa} \times 2 \times \text{cyb} \times 2 \)
\( \text{f11} = 0.5 \times \log(\text{rls}) \)
\( \text{f12} = 0.5 \times \log(\text{r2s}) \)
\( \text{tb2} = -\text{con} \times (\text{f11} - \text{f12}) \)
\( \text{tb3} = -\text{cyb} \times \text{tb3} + \text{con} \times (\text{cma} \times \text{f11} - \text{cpa} \times \text{f12}) \)
\( \text{tb1} = -\text{cyb} \times \text{tb3} + \text{con} \times (\text{cyb} / \text{rls} - \text{cyb} / \text{r2s}) \)
\( \text{tb4} = \text{con} \times (\text{cma} / \text{rls} - \text{cpa} / \text{r2s}) \)
\( \text{tb5} = \text{con} \times (\text{cma} / \text{rls} - \text{cpa} / \text{r2s}) \)
\( \text{asnt} = \text{pr2} \times \text{cosg} \times \text{tb3} + \text{pr1} \times \text{sing} \times \text{tb2} + \text{cyb} \times (\text{sing} \times \text{tb4} \times \text{cosg} \times \text{tb5}) \)
Solving the Problem

\[
\begin{align*}
\text{asnt} &= -pr_1 \cos g \cdot tb_2 + pr_2 \sin g \cdot tb_3 + cyb \cdot (\cos g \cdot tb_4 - \sin g \cdot tb_5) \\
\text{anst} &= -pr_2 \sin g \cdot tb_3 + pr_1 \cos g \cdot tb_2 + cyb \cdot (\cos g \cdot tb_4 - \sin g \cdot tb_5) \\
\text{annt} &= pr_1 \sin g \cdot tb_3 + pr_2 \cos g \cdot tb_2 - cyb \cdot (\sin g \cdot tb_4 + \cos g \cdot tb_5) \\
\text{bst} &= pr_3 \cos g \cdot tb_1 + cyb \cdot (\sin g \cdot tb_2 - \cos g \cdot tb_3) \\
\text{bsnt} &= pr_3 \sin g \cdot tb_1 + cyb \cdot (\cos g \cdot tb_2 + \sin g \cdot tb_3) \\
\text{bnst} &= -pr_3 \sin g \cdot tb_1 + cyb \cdot (\cos g \cdot tb_2 + \sin g \cdot tb_3) \\
\text{bnnt} &= pr_3 \cos g \cdot tb_1 - cyb \cdot (\sin g \cdot tb_2 - \cos g \cdot tb_3) \\
\text{ass} &= \text{ass} + \text{msym} \cdot \text{asst} \\
\text{asn} &= \text{asn} + \text{asnt} \\
\text{ans} &= \text{ans} + \text{msym} \cdot \text{anst} \\
\text{ann} &= \text{ann} + \text{annt} \\
\text{bs} &= 0.5 \cdot (b(2 \cdot j - 1) - ssO)/g \\
\text{bn} &= 0.5 \cdot (b(2 \cdot j) - sn0)/g \\
\text{bs} &= \frac{r(2 \cdot i - 1) + \text{bs} \cdot \text{bs} + \text{bs} \cdot \text{bn} \cdot \text{bn}}{} \\
\text{bn} &= \frac{r(2 \cdot i)}{} \\
\text{bs} &= 0.5 \cdot (b(2 \cdot j - 1) - \text{ss}O)/g \\
\text{bn} &= 0.5 \cdot (b(2 \cdot j) - \text{sn}0)/g \\
\text{bs} &= \frac{r(2 \cdot i - 1) + \text{bs} \cdot \text{bs} + \text{bs} \cdot \text{bn} \cdot \text{bn}}{} \\
\text{bn} &= \frac{r(2 \cdot i)}{} \\
\text{bs} &= -bss \\
\text{bn} &= -bsn \\
\text{bs} &= -bns \\
\text{bn} &= -bnn \\
\end{align*}
\]

Set up influence coeff matrix and RHS of discrete system

subroutine SETUP

\$ (i, j, bckodj, g, ssO, sn0, \\
\$   ass, asn, ans, ann, \\
\$   bss, bsn, bns, bnn, \\
\$   b(!), c(nc,*), r(*) \\
\$ implicit none \\
integers i, j, n, nc, bckodj \\
reals ssO, sn0, g, bs, bn \\
reals ass, asn, ans, ann, bss, bsn, bns, bnn \\
reals b(\cdot), c(nc,\cdot), r(*) \\
if (bckodj .eq. 0) then \\
c(2\cdot i-1,2\cdot j-1) = ass \\
c(2\cdot i-1,2\cdot j ) = asn \\
c(2\cdot i ,2\cdot j-1) = ans \\
c(2\cdot i ,2\cdot j ) = ann \\
bs = 0.5*(b(2\cdot j-1)-ss0)/g \\
bn = 0.5*(b(2\cdot j )-sn0)/g \\
r(2\cdot i-1) = r(2\cdot i-1) + bs*bs + bsn*bn \\
r(2\cdot i ) = r(2\cdot i ) + bns*bs + bnn*bn \\
else if (bckodj .eq. 1) then \\
c(2\cdot i-1,2\cdot j-1) = -bss \\
c(2\cdot i-1,2\cdot j ) = -bsn \\
c(2\cdot i ,2\cdot j-1) = -bns \\
end if
Solving the Problem

c(2*i ,2*j ) = -bnn
r(2*i-1) = r(2*i-1) - ass*b(2*j-1) - asn*b(2*j)
r(2*i ) = r(2*i ) - ans*b(2*j-1) - ann*b(2*j)
else if (bckodj .eq. 2) then
  c(2*i-1,2*j-1) = -bss
c(2*i-1,2*j ) = asn
c(2*i ,2*j-1) = -bsn
c(2*i ,2*j ) = ann
bn = 0.5*(b(2*j )-sn0)/g
r(2*i-1) = r(2*i-1) - ass*b(2*j-1) + bsn*bn
r(2*i ) = r(2*i ) - ans*b(2*j-1) + bnn*bn
else
  c(2*i-1,2*j-1) = ass
c(2*i-1,2*j ) = -bsn
c(2*i ,2*j-1) = ans
c(2*i ,2*j ) = -bsn
bs = 0.5*(b(2*j-1)-ss0)/g
r(2*i-1) = r(2*i-1) + bss*bs - asn*b(2*j)
r(2*i ) = r(2*i ) + bns*bs - ann*b(2*j)
end if
return
end

Subroutine SOLVE calls GAUSSER, which is a naive implementation of unsymmetric Gauss elimination without pivoting:

Solve algebraic equation system A x = b by Gauss elimination

subroutine GAUSSER
(a, b, x, n, na, ising)

implicit none
integer n, na, ising
real a(na,+), b(+), x(+), c, sum
integer i, j, k

ising = 0
do 2000 j = 1,n-1
  if (a(j,j) .eq. 0.0) then
    ising = j
    return
  end if
  do 1500 k = j+1,n
    c = a(k,j)/a(j,j)
    do 1400 i = j,n
      a(k,i) = a(k,i) - c*a(j,i)
      1400  continue
    b(k) = b(k) - c*b(j)
  do 1500
2000 continue

Solving the Problem

1500 continue
2000 continue

* x(n) = b(n)/a(n,n)
do 3000 j = n-1,1,-1
   sum = 0.0
do 2500 i = j+1,n
   sum = sum + a(j,i)*x(i)
2500 continue
x(j) = (b(j)-sum)/a(j,j)
3000 continue
return
end

(The only redeeming quality about GAUSSER is that the code is quite short; in fact, it's about the shortest possible implementation of a linear equation solver.)
Solving the Problem

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9. Printing Data

One area in which interactive operation excels is data display. If you are using an interactive Processor for an engineering design task, you can selectively trim the otherwise voluminous output to the important essentials. Conversely, if you are debugging a new or modified implementation, you may want more output than is normally required; for example printing the influence coefficient matrix.

What goes for printed output applies with equal force to graphic output. We are not going to illustrate graphic displays here, however, since the details depend strongly on the output device and the plotting software you are using.

The PRINT command is similar to the DEFINE command in that it takes a second keyword that specifies what is to be printed:

- **SEGMENTS**
  Prints segment geometry data and number of elements per segment.
- **BOUNDARY_CONDITIONS**
  Prints boundary condition (BC) code and prescribed boundary values for each segment.
- **SYMMETRY_CONDITIONS**
  Prints symmetry conditions if any in effect.
- **MATERIAL**
  Prints material property data.
- **PRESTRESS**
  Prints prestress data.
- **FIELD_LOCATIONS**
  Prints detailed boundary-element data produced by subroutine BUILD (this is primarily for debugging).
- **ELEMENTS**
  Prints the matrix C of influence coefficients assembled by GENERATE (this is primarily for debugging).
- **COEFFICIENTS**
  Prints the right-hand side (forcing) vector \( \mathbf{r} \) assembled by GENERATE (this is primarily for debugging).
- **SOLUTION**
  Prints the solution vector \( \mathbf{x} \) calculated by SOLVE (this is primarily for debugging).
- **RESULTS**
  Print stresses and displacements at boundary-element midpoints or at output field locations, depending on a command qualifier.

In this section none of the subroutines load or store data. However, in a real Processor the results, such as displacements and stresses, which are computed under the PRINT RESULTS command would be stored. They would be stored because these results are often plotted or reordered for tabulation. The plots and tables are used to study the results or for inclusion in a report or as presentation. By now you should be able to modify PRINT RESULTS, §9.3, to LOAD and STORE data for the post-processing activities described above. Be brave give it a try.

The PRINT command is processed by subroutine PRINT, which has a “case” structure similar to that of subroutine DEFINE:
Interpret PRINT command

subroutine PRINT

implicit none
character key*8, CCLVAL*8
integer ICLTYP
logical CMATCH

if (ICLTYP(2) .le. 0) then
   call CLREAD ('PRINT what? ',
             $'
             $'BOUNDARY, ELEMENTS, COEFFICIENTS,'
             $'FIELD, MATERIAL, PRESTRESS&'
             $'RESULTS, RHS, SOLUTION, SYMMETRY')
   key = CCLVAL(1)
else
   key = CCLVAL(2)
end if
if (CMATCH (key, 'BOUNDARY')) then
   call PRINT_BOUNDARY_CONDITIONS
else if (CMATCH (key, 'COEFFICIENTS') .or.
             $CMATCH (key, 'INFLUENCE')) then
   call PRINT_INFLUENCE_COEFFICIENTS
else if (CMATCH (key, 'ELEMENTS')) then
   call PRINT_ELEMENTS
else if (CMATCH (key, 'FIELD')) then
   call PRINT_FIELD_LOCATIONS
else if (CMATCH (key, 'MATERIAL')) then
   call PRINT_MATERIAL
else if (CMATCH (key, 'PRESTRESS')) then
   call PRINT_PRESTRESS
else if (CMATCH (key, 'RESULTS')) then
   call PRINT_RESULTS
else if (CMATCH (key, 'RHS')) then
   call PRINT_RHS_VECTOR
else if (CMATCH (key, 'SEGMENT')) then
   call PRINT_SEGMENTS
else if (CMATCH (key, 'SOLUTION')) then
   call PRINT_SOLUTION_VECTOR
else if (CMATCH (key, 'SYMMETRY')) then
   call PRINT_SYMMETRY_CONDITIONS
else
   print *, '*** Illegal or ambiguous keyword ',key,' after PRINT'
end if
return
end
Subroutine PRINT provides our second example of an implementation that *prompts for missing data*. See DB OPEN, §6.0, for another example. If you type only the keyword PRINT followed by a carriage return, you will see the prompt

```
Print what?
```

on the screen, and you are supposed to type the next keyword, e.g., SEGMENTS that you forgot. (Notice that this friendly technique was not used for the DEFINE command explained in §7.0; instead subroutine DEFINE complains about missing keywords after DEFINE.)

Next we examine the subordinate routines.
9.1 Printing Input Data

The implementation of the subroutines that print segment, boundary condition, symmetry, material, prestress, and field-location data is straightforward and so are simply listed next as a group:

```
* Print segment data
subroutine PRINT_SEGMENTS
implicit none
include 'segment.inc'
integer i, k

k = 0
do 2000 i = 1,MAXSEG
  if (segdef(i) .gt. 0) then
    if (k .eq. 0) then
      print '(/A/A6,A6,4AI2)',
        ' Boundary Segment Data',
        'Segm', 'Elements', 'Xbeg', 'Ybeg', 'Xend', 'Yend'
      end if
    k = k + 1
    print '(16,19,3X,4GI2.4)',
      i, numel(i), xbeg(i), ybeg(i), xend(i), yend(i)
  end if
continue
if (k .eq. 0) then
  print *, 'Segment tables are empty'
end if
print *, return
end ....................................
```

Print boundary data in response to a PRINT BOUNDARY command

```
* Print boundary data in response to a PRINT BOUNDARY command
subroutine PRINT_BOUNDARY_CONDITIONS
implicit none
include 'segment.inc'
integer i, k
character*9 given(0:3)
data given /'SS and NS', 'SD and ND', 'SD and NS', 'SS and ND'/

k = 0
do 2000 i = 1,MAXSEG
  if (segdef(i) .gt. 0) then
```

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if (k .eq. 0) then
    print '(/A/A6,A11,2A12)',
    $ ' Boundary Conditions Data', 'Segm',
    $ 'Given', 'Shear', 'Normal'
end if
k = k + 1
print '(/5,1X,A11,3X,1P2G12.3)',
    $ i, given(kode(i)), bvs(i), bvn(i)
end if
continue
if (k .eq. 0) then
    print _, 'Boundary tables are empty'
end if
print *,
return
end

* Print symmetry data
* subroutine PRINT_SYMMETRY_CONDITIONS
* implicit none
include 'symmetry.inc'
print '(/A)', ' Symmetry Data'
if (ksym .eq. 3) then
    print _, 'Symmetry about axis X=',xsym
    print _, ' and axis Y=',ysym
else if (ksym .eq. 1) then
    print _, 'Symmetry about axis X=',xsym
else if (ksym .eq. 2) then
    print _, 'Symmetry about axis Y=',ysym
else
    print *, 'No symmetry conditions'
end if
print *,
return
end

* Print material property data
* subroutine PRINT_MATERIAL
* implicit none
include 'material.inc'
print '(/A)', ' Material Property Data'

---

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print '('' Elastic modulus:'',.1PE12.3)', em
print '('' Poisson''s ratio:'',.F12.3)', pr
print *, ' '
return
end

Print field location data

subroutine PRINT_FIELD_LOCATIONS

implicit none
include 'output.inc'
integer i, k

k = 0
do 2000 i = 1,MAXLII
    if (lindef(i) .gt. 0) then
        if (k .eq. 0) then
            print '(IA/A6,Ag>A9,3A12)',
            $ 'Field Location Data',
            $ 'Line', 'Int.Pts', 'x-first', 'y-first',
            $ 'x-last', 'y-last'
        end if
        k = k + i
        print '(I6,I9,4G12.4)',
        $ i, nintop(i), xfirst(i), yfirst(i), xlast(i), ylast(i)
    end if
2000 continue
if (k .eq. 0) then
    print *, 'FIELD Location Tables are empty'
end if
print *, '
return
end
9.2 Debug-Oriented Print Commands

The PRINT_ELEMENTS, PRINT COEFFICIENTS, PRINT RHS and PRINT SOLUTION are
detailed print commands primarily useful in debug situations. They are implemented in
the following subroutines:

* Print detailed boundary element data
* subroutine PRINT_ELEMENTS
* implicit none
* include 'segment.inc'
include 'element.inc'
integer m

if (numbe .le. 0) then
  print *, 'Boundary element table empty'
  return
end if
print '('A/A5,A8.2A11,A12,A8,A9,A12)'
$  Boundary Element Data',
$  'Elem', 'Xmid', 'Ymid', 'Length',
$  'Orient', 'BCode', 'Shear', 'Normal'
do 2000 m = 1,numbe
  print '('I5,I1P3GII.3,0PFIO.2,I6,1P2GI2.3)
  m,xme(m),yme(m),2._:hleng(m),
  (180./3.14159265)*atan2(sinbet(m),cosbet(m)),
  kod(m), b(2*m-1).b(2*m)
2000 continue
print ' ',
return
end

* Print influence coefficient matrix
* subroutine PRINT_INFLUENCE_COEFFICIENTS
* implicit none
* include 'segment.inc'
include 'element.inc'

print '('A)', ' Influence Coefficient Matrix'
call PRINT_REAL_MATRIX (c, MAXEQS, 2*numbe, 2*numbe)
print *, ' ',
return
end

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* Print right hand side vector

subroutine PRINT_RHS_VECTOR
implicit none
include 'segment.inc'
include 'element.inc'

print '(/A)', ' Right Hand Side (Forcing) Vector'
call PRINT_REAL_MATRIX (r, 1, 1, 2+numbe)
print *, ''
return
end

* Print right hand side vector

subroutine PRINT_SOLUTION_VECTOR
implicit none
include 'segment.inc'
include 'element.inc'

print '(/A)', ' Solution Vector'
call PRINT_REAL_MATRIX (x, 1, 1, 2+numbe)
print *, ''
return
end

The last three subroutines call PRIMATRIX, which is a "no frills" array printer:

* Print real matrix (or vector) in 6-column template

subroutine PRINT_REAL_MATRIX
integer na, m, n, i, j, jref
real a(na, )
do 4000 jref = 0, n-1, 6
   print '(1X,6I12)', (j, j=jref+1, min(jref+6, n))
do 3000 i = 1, m
   print '(I4,1P6E12.4)', i, (a(i, j), j=jref+1, min(jref+6, n))
3000 continue
4000 continue
return
end
9.3 Printing Results

The PRINT RESULTS command without a qualifier lists stresses and displacements computed at boundary element midpoints. If qualifier FIELD appears, the command refers to the field points previously defined. This switch is implemented in subroutine PRINT RESULTS:

```
* Process PRINT RESULTS command
+ subroutine PRINT_RESULTS
* implicit none
integer ICLSEQ
*
if (ICLSEQ(3,'FIELD').eq. 0) then
  call PRINT_BOUNDARY_RESULTS
else
  call PRINT_FIELD_RESULTS
end if
return
end
```

The code above provides an example of the use of ICLSEQ [1], §5.3, to test for the existence of a specific qualifier, in this case FIELD.

9.3.1 Printing Boundary Results

This is done by subroutine PRINT_BOUNDARY_RESULTS, the implementation of which is straightforward:

```
* Print stresses and displacement @ boundary element midpoints
+ subroutine PRINT_BOUNDARY_RESULTS
*
implicit none
include 'segment.inc'
include 'element.inc'
include 'material.inc'
include 'prestress.inc'
integer k
real g, ss0, sn0, sinbi, cosbi
real us, un, ux, uy, sign, sigs
*
print '(/A)', ' Displacements and Stresses at'/
$ ' Boundary Element Midpoints'
print '(A5,A9,5A11)', 'Elem', 'u_s', 'u_n', 'u_x', 'u_y',
```
$ \quad g = 0.5*E/(1.+pr) \\
'sig_s', 'sig_n'

\begin{verbatim}
do 2000 k = 1, numbe
  us = x(2*k-1)
  un = x(2*k)
  sigs = b(2*k-1)
  sign = b(2*k)
  if (kod(k) .eq. 1) then
    un = b(2*k-1)
    us = b(2*k)
    sigs = x(2*k-1)
    sign = x(2*k)
  else if (kod(k) .eq. 2) then
    us = b(2*k-1)
    sigs = x(2*k-1)
  else if (kod(k) .eq. 3) then
    un = b(2*k)
    sign = b(2*k)
  end if
  sinbi = sinbet(k)
  cosbi = cosbet(k)
  ux = us*cosbi - un*sinbi
  uy = us*sinbi + un*cosbi
  print '(I5,1P6G11.3)', k, us, un, ux, uy, sigs, sign
2000 continue
print *, ' ' 
return
end
\end{verbatim}

9.3.2 Printing Field Results

Showing displacement and stresses at field points is complicated by the fact that, unlike finite element programs, such values are not readily available but must be calculated as part of the display procedure. This will become evident as one shows the coding of subroutine PRINT_FIELD_RESULTS:

* Print stresses and displacements at specified field points

* subroutine PRINT_FIELD_RESULTS

* implicit none
* include 'output.inc'
* integer m, p, points
* real xp, yp, ux, uy, sigxx, sigyy, sigxy, f
* logical skip
* print '(/A)', ' Displacements and Stresses at'//
$ ' Specified Field Points'
*

do 3000 m = 1,MAXLIN
   if (lindef(m) .eq. 0) go to 3000
   print '(AS,2AIO.A8,4AII)', ' Lin', 'x', 'y', 'u_x', 'u_y',
      'sig_xx', 'sig_yy', 'sig_xy'
   points = nintop(m) + 2
   if (xfirst(m) .eq. xlast(m).and.
      yfirst(m) .eq. ylast(m)) points = 1
   f = 0.0
   do 2000 p = 1,points
      if (points .gt. 1) f = real(p-1)/(points-1)
      xp = xfirst(m)+(1.0-f) + xlast(m)*f
      yp = yfirst(m)+(1.0-f) + ylast(m)*f
      call FIELDP (xp, yp, ux, uy, sigxx, sigyy, sigxy, skip)
      if (skip) then
         print '(IS,2F10.3,6X,A)', m, xp, yp,
            'Point is too close to boundary'
      else
         print '(I5.2FIO.3,1P5GII.3)', m, xp, yp, ux, uy,
            sigxx, sigyy, sigxy
      end if
   2000 continue
   print _, ' '
3000 continue
return
end

Subroutine FIELDP receives the location XP,YP of the field point and returns the displace-
ment components u_x and u_y, and the stress components σ_xx, σ_yy and σ_xy:

* Compute stresses and displacements at field point
*

subroutine FIELDP

implicit none
include 'segment.inc'
include 'element.inc'
include 'material.inc'
include 'symmetry.inc'
include 'prestress.inc'
real xp, yp, us, un, ux, uy, sigxx, sigyy, sigxy
logical skip
real uxus, uxun, uxss, uxsn
real uyus, uyun, uyss, uysn
real sxxus, sxxun, sxxss, sxxsn
real syyus, syyun, syyss, syysn
real sxyus, sxyun, sxyss, sxysn
real xj, yj, sj, cosbj, sinbj
real usj, unj, ssj, snj, ssng
real g, ss0, sn0
integer j

skip = .false.
ux = 0.0
uy = 0.0
sigxx = sxx0
sigyy = syy0
sigxy = sxy0
g = 0.5*em/(1.+pr)

do 2000 j = 1,numbe
uxus = 0.0
uxun = 0.0
uxss = 0.0
uxsn = 0.0
uyun = 0.0
uyun = 0.0
uyss = 0.0
uyun = 0.0
uyun = 0.0
sxxus = 0.0
sxxun = 0.0
sxxss = 0.0
sxxsn = 0.0
syyus = 0.0
syyun = 0.0
syyss = 0.0
syyun = 0.0
syyun = 0.0
sxyus = 0.0
sxyun = 0.0
sxyss = 0.0
sxyun = 0.0
sxyun = 0.0
xj = xme(j)
yj = yme(j)
sj = hleng(j)
if ( ((xp-xj)**2+(yp-yj)**2+1.01*(sj)**2) ) then
skip = .true.
return
end if
cosbj = cosbet(j)
sinbj = sinbet(j)
ss0 = (syy0-sxx0)*sinbj*cosbj + sxy0*(cosbj**2-sinbj**2)
\[ sn0 = sxx0*\sin bj**2 - 2.*sxy0*\sin bj*\cos bj + syy0*\cos bj**2 \]
call SOMIGLIANA (xp, yp, xj, yj, sj, 
\[ \begin{align*} &1, \text{ em, pr, } \cos bj, \sin bj, \\
&uxus, uxun, uxss, uxsn, \\
&uyus, uyun, uyss, uysn, \\
&sxxus, sxxun, sxxss, sxxsn, \\
&syyus, syyun, syyss, syysn, \\
&sxyus, sxyun, sxyss, sxysn \end{align*} \]
if (ksym .eq. 1 .or. ksym .eq. 3) then
  call SOMIGLIANA (xp, yp, 2.*xsym-xme(j), yj, sj, 
  \[ \begin{align*} &-1, \text{ em, pr, } \cos bj, -\sin bj, \\
&uxus, uxun, uxss, uxsn, \\
&uyus, uyun, uyss, uysn, \\
&sxxus, sxxun, sxxss, sxxsn, \\
&syyus, syyun, syyss, syysn, \\
&sxyus, sxyun, sxyss, sxysn \end{align*} \]
end if
if (ksym .eq. 2 .or. ksym .eq. 3) then
  call SOMIGLIANA (xp, yp, xj, 2.*ysym-yme(j), sj, 
  \[ \begin{align*} &1, \text{ em, pr, } -\cos bj, \sin bj, \\
&uxus, uxun, uxss, uxsn, \\
&uyus, uyun, uyss, uysn, \\
&sxxus, sxxun, sxxss, sxxsn, \\
&syyus, syyun, syyss, syysn, \\
&sxyus, sxyun, sxyss, sxysn \end{align*} \]
end if
if (ksym .eq. 3) then
  call SOMIGLIANA (xp, yp, 2.*xsym-xme(j), 2.*ysym-yme(j), sj, 
  \[ \begin{align*} &1, \text{ em, pr, } -\cos bj, -\sin bj, \\
&uxus, uxun, uxss, uxsn, \\
&uyus, uyun, uyss, uysn, \\
&sxxus, sxxun, sxxss, sxxsn, \\
&syyus, syyun, syyss, syysn, \\
&sxyus, sxyun, sxyss, sxysn \end{align*} \]
end if
usj = x(2*j-1)
unj = x(2+j )
ssj = b(2*j-1) - ss0
snj = b(2*j ) - sn0
if (kod(j) .eq. 1) then
  usj = b(2+j-1)
  unj = b(2+j )
  ssj = x(2+j-1)
  snj = x(2+j )
else if (kod(j) .eq. 2) then
  usj = b(2+j-1)
  ssj = x(2+j )
else if (kod(j) .eq. 3) then

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unj = b(2*j )
snj = x(2*j )
end if
ssg = 0.5*ssj/g
sng = 0.5*snj/g
ux = ux + uxus*usj + uxun*unj + uxss*ssg + uxsn*sng
uy = uy + uyus*usj + uyun*unj + uyss*ssg + uysn*sng
usj = 2.*g*usj
unj = 2.*g*unj
sigxx = sigxx + sxxus_usj + sxxun*unj + sxxss*ssj + sxxsn*snj
sigyy = sigyy + syyus*usj + syyun*unj + syyss*ssj + syysn*snj
sigxy = sigxy + sxysu*usj + sxyun*unj + sxysss*ssj + sxysn*snj
continue
return
end

Finally, FIELDP calls subroutine SOMIGLIAIIA to evaluate the important boundary-on-field-point influence coefficients:

Calculate field influence coefficients from Somigliana’s formula:

subroutine SOMIGLIAIIA
$ (x, y, xj, yj, aj, em, pr, cosb, sinb.$
$           uxus, uxun, uxss, uxsn,$
$           uyus, uyun, uyss, uysn,$
$           sxxus, sxxun, sxxss, sxxsn,$
$           syyus, syyun, syyss, syysn,$
$           sxysus, sxysun, sxysss, sxysnn)

implicit none
real x, y, xj, yj, aj, em, pr, cosb, sinb
real uxus, uxun, uxss, uxsn
real uyus, uyun, uyss, uysn
real sxxus, sxxun, sxxss, sxxsn
real syyus, syyun, syyss, syysn
real sxusy, sxysun, sxysss, sxysnt
integer msym
real pi, con, pr1, pr2, pr3
real cxb, cyb, cosg, sing, cpa, cma
real ris, r2s, fl1, fl2
real tbi, tbs, b3, tbs, t5, t6, t7
real uxusst, uxsst, uxsnt
real uyust, uyunt, uysst, uysnt
real sxxust, sxxunst, sxxsst, sxxssnt
real syyust, syyunst, syysst, syysnt
real sxyst, sxyst, sxysst, sxystnt
real cosb2, sin2b, cos2b, sin2b
\[
\pi = 4.0 \cdot \text{atan}_2(1.,1.)
\]
\[
\text{con} = 1.0/(4.0 \cdot \pi + (1.0 - \text{pr}))
\]
\[
\text{pr}_1 = 1.0 - 2.0 \cdot \text{pr}
\]
\[
\text{pr}_2 = 2.0 \cdot (1.0 - \text{pr})
\]
\[
\text{pr}_3 = 3.0 - 4.0 \cdot \text{pr}
\]
\[
\text{cxb} = (x - x_j) \cdot \cos_b + (y - y_j) \cdot \sin_b
\]
\[
\text{cyb} = -(x - x_j) \cdot \sin_b + (y - y_j) \cdot \cos_b
\]
\[
\text{cma} = \text{cxb} - \text{a}_j
\]
\[
\text{cpa} = \text{cxb} + \text{a}_j
\]
\[
\text{r}_1^2 = \text{cma}^2 + 2.0 \cdot \text{cyb}^2
\]
\[
\text{r}_2^2 = \text{cpa}^2 + 2.0 \cdot \text{cyb}^2
\]
\[
f_11 = 0.5 \cdot \log(\text{r}_1^2)
\]
\[
f_12 = 0.5 \cdot \log(\text{r}_2^2)
\]
\[
\text{tb}_2 = -(\text{con} \cdot (f_11 - f_12))
\]
\[
\text{tb}_3 = \text{con} \cdot (\text{atan}_2(\text{cpa}, \text{cyb}) - \text{atan}_2(\text{cma}, \text{cyb}))
\]
\[
\text{tb}_4 = \text{con} \cdot (\text{cyb} \cdot \text{r}_1^2 - \text{cyb} \cdot \text{r}_2^2)
\]
\[
\text{tb}_5 = \text{con} \cdot (\text{cma} \cdot \text{cpa} - \text{cpa} \cdot \text{r}_2^2)
\]
\[
\text{tb}_6 = \text{con} \cdot ((\text{cma}^2 - 2.0 \cdot \text{cyb}^2) / \text{r}_1^2 - (\text{cpa}^2 - 2.0 \cdot \text{cyb}^2) / \text{r}_2^2)
\]
\[
\text{tb}_7 = -(\text{con} \cdot 2.0 \cdot \text{cyb} \cdot (\text{cpa} / \text{r}_2^2 - \text{cpa} / \text{r}_1^2))
\]
\[
\text{uxust} = \text{p}_1 \cdot \sin_b \cdot \text{tb}_2 - \text{p}_2 \cdot \cos_b \cdot \text{tb}_3 + \text{cyb} \cdot (\text{sinb} \cdot \text{tb}_2 - \text{cosb} \cdot \text{tb}_3)
\]
\[
\text{uxunt} = \text{p}_1 \cdot \cos_b \cdot \text{tb}_2 + \text{p}_2 \cdot \sin_b \cdot \text{tb}_3 - \text{cyb} \cdot (\text{cosb} \cdot \text{tb}_2 + \text{sinb} \cdot \text{tb}_3)
\]
\[
\text{uxsst} = \text{p}_3 \cdot \cos_b \cdot \text{tb}_2 + \text{p}_2 \cdot \sin_b \cdot \text{tb}_3 - \text{cyb} \cdot (\text{cosb} \cdot \text{tb}_2 + \text{sinb} \cdot \text{tb}_3)
\]
\[
\text{uxsnt} = \text{p}_3 \cdot \cos_b \cdot \text{tb}_2 - \text{p}_2 \cdot \sin_b \cdot \text{tb}_3 - \text{cyb} \cdot (\text{cosb} \cdot \text{tb}_2 + \text{sinb} \cdot \text{tb}_3)
\]
\[
\text{uxust} = 2.0 \cdot \text{cosb} \cdot \text{tb}_4 + \sin_b \cdot \text{tb}_5 - \text{cyb} \cdot (\text{sinb} \cdot \text{tb}_4 - \text{cosb} \cdot \text{tb}_5)
\]
\[
\text{syyust} = 2.0 \cdot \text{sinb} \cdot \text{tb}_4 - \text{sin}_b \cdot \text{tb}_5 + \text{cyb} \cdot (\text{cosb} \cdot \text{tb}_4 - \text{sinb} \cdot \text{tb}_5)
\]
\[
\text{sxyst} = \text{sin}_b \cdot \text{tb}_4 - \text{cosb}_b \cdot \text{tb}_5 - \text{cyb} \cdot (\text{sinb} \cdot \text{tb}_4 + \text{cosb} \cdot \text{tb}_5)
\]
\[
\text{sxxust} = -\text{tb}_5 - \text{cyb} \cdot (\text{sin}_b \cdot \text{tb}_4 + \text{cosb} \cdot \text{tb}_5)
\]
\[
\text{syyyst} = -\text{tb}_5 + \text{cyb} \cdot (\text{sin}_b \cdot \text{tb}_4 - \text{cosb} \cdot \text{tb}_5)
\]
\[
\text{sxysst} = \text{sin}_b \cdot (\text{cosb} \cdot \text{tb}_4 + \text{sinb} \cdot \text{tb}_5)
\]
\[
\text{sxxsst} = \text{sin}_b \cdot (\text{cosb} \cdot \text{tb}_4 - \text{sinb} \cdot \text{tb}_5)
\]
\[
\text{sxysnt} = \text{sin}_b \cdot (\text{cosb} \cdot \text{tb}_4 - \text{sinb} \cdot \text{tb}_5)
\]
$ + cyb*(sin2b*tb4-cos2b*tb5)

sxxsnt = -tb3 + pr1*(sin2b*tb2+cos2b*tb3)

$ + cyb*(sin2b*tb4-cos2b*tb5)

syysnt = -tb3 - pr1*(sin2b*tb2+cos2b*tb3)

$ - cyb*(sin2b*tb4-cos2b*tb5)

sxysnt = - pr1*(cos2b*tb2-sin2b*tb3)

$ - cyb*(cos2b*tb4+sin2b*tb5)

uxus =  uxus + msym*uxust

uxun =  uxun + uxunt

uxss =  uxss + msym*uxsst

uxsn =  uxsn + uxsnt

uyus =  uyus + msym*uyust

uyun =  uyun + uyunt

uyss =  uyss + msym*uyss

uysn =  uysn + uysnt

sxxus = sxxus + msym*sxxust

sxxun = sxxun + sxxunt

sxxss = sxxss + msym*sxxsst

sxxsn = sxxsn + sxxsnt

syys = syys + msym*syyst

syyn = syyn + syyst

syys = syys + msym*syys

syysn = syysn + syysnt

sxyus = sxyus + msym*sxyust

sxyun = sxyun + sxyunt

sxyss = sxyss + msym*sxyss

sxyen = sxyen + sxyent

return

end

The DBEM2 Processor is complete.
10. DBEM2 Structure

After all the coding details given in §5.0 through §9.0 it is perhaps refreshing to get an overall picture of the structure of DBEM2. A hierarchical diagram of the module structure provides a portion of the picture:

```
DBEM2
  DOCOMMAND
    BUILD
    CLEAR
    CLOSE
    DEFINE
      DEFINE.BOUNDARY_CONDITIONS
      BCVALUES
      DEFINE.ELEMENTS
      DEFINE.MATERIAL
      DEFINE.FIELD_LOCATIONS
      DEFINE.PRESTRESS
      DEFINE_SEGMENTS
      DEFINE_SYMMETRY
    GENERATE
      COEFF
      SETUP
    OPEN
    PRINT
      PRINT.BOUNDARY_CONDITIONS
      PRINT.BOUNDARY_RESULTS
      PRINT.COEFFICIENTS
        PRINT_REAL_MATRIX
      PRINT_ELEMENTS
      PRINT_FIELD_LOCATIONS
      PRINT_FIELD_RESULTS
      FIELD
        SOMIGLIANA
      PRINT_MATERIAL
      PRINT_PRESTRESS
      PRINT_RHS
        PRINT_REAL_MATRIX
      PRINT_SEGMENTS
      PRINT_SOLUTION
        PRINT_REAL_MATRIX
      PRINT_SYMMETRY
    SOLVE
      GAUSSER
    STOP
```
This diagram of course excludes the NICE utilities such as the CLIP and GAL-DBM system. With this omission noted, the deepest module level is five. This is a feature symptomatic of a fairly simple Processor. (Actual production Processors in the NICE system reach module levels of order 15-20.)

Another part of the picture is provided by a diagram of the DBEM2 GAL Library structure of datasets and associated record names:

```
DBEM2 GAL LIBRARY
  BCVALUES
    BVN
    BVS
    KODE
  COEFF
    C
  ELEMENT
    B
    COSBET
    HLENG
    KOD
    NUMBE
    SINBET
    XME
    YME
  FIELD
    L1NDEF
    NINTOP
    XFIRST
    XLAST
    YFIRST
    YLAST
  MATERIAL
    EM
    PR
  PRESTRESS
    SXXO
    SXYO
    SYYO
  RHS
    R
```
Of course not all of the datasets will appear in all Libraries; only the datasets corresponding to the data that are stored appears.

A diagram that shows the interaction between the database and the Processor modules provides a connection between the first two diagrams. In the diagram below the Processor module commands are shown on the left and the corresponding GAL Library datasets are shown on the left. The \( \rightarrow \) symbol indicates that the data in the designated dataset is either loaded or stored within the designated module. The \( \leftarrow \) symbol indicates that the corresponding dataset must be loaded before the designated command should be entered. By searching the diagram for the \( \leftrightarrow \) symbol you can see which command along with a LOAD sub-command must be entered to load the desired dataset.
DEFINE
DEFINE_SEGMENTS ← Dataset: SEGMENT
DEFINE_ELEMENTS ← Updates Record NUMEL in Dataset: SEGMENT
DEFINE_BOUNDARY_CONDITIONS ← Dataset: BCVALUES
DEFINE_MATERIAL ← Dataset: MATERIAL
DEFINE_SYMMETRY ← Dataset: SYMMETRY
DEFINE_PRESTRESS ← Dataset: PRESTRESS
DEFINE_FIELD_LOCATIONS ← Dataset: FIELD
BUILD ← Dataset: SEGMENT
← Dataset: BCVALUES
← Dataset: ELEMENT
GENERATE ← Dataset: ELEMENT
← Dataset: MATERIAL
← Dataset: SYMMETRY
← Dataset: PRESTRESS
← Dataset: COEFF
← Dataset: RHS
SOLVE ← Dataset: COEFF
← Dataset: RHS
← Dataset: ELEMENT
← Dataset: SOLUTION
PRINT
PRINT_BOUNDARY_CONDITIONS ← Dataset: BCVALUES
PRINT_BOUNDARY_RESULTS ← Dataset: MATERIAL
← Dataset: ELEMENT
← Dataset: SOLUTION
PRINT_COEFFICIENTS ← Dataset: ELEMENT
← Dataset: COEFF
PRINT_ELEMENTS ← Dataset: ELEMENT
PRINT_FIELD_LOCATIONS ← Dataset: FIELD
PRINT_FIELD_RESULTS ← Dataset: FIELD
← Dataset: PRESTRESS
← Dataset: MATERIAL
← Dataset: ELEMENT
← Dataset: SYMMETRY
← Dataset: SOLUTION
PRINT_MATERIAL ← Dataset: MATERIAL
PRINT_PRESTRESS ← Dataset: PRESTRESS
PRINT_RHS ← Dataset: ELEMENT
← Dataset: RHS
PRINT_SEGMENTS ← Dataset: SEGMENT
PRINT_SOLUTION ← Dataset: ELEMENT
← Dataset: SOLUTION
PRINT_SYMMETRY ← Dataset: SYMMETRY
Now we have enough information on how the Processor and the GAL-DBM work together to discuss some potential extensions to this simple Processor.

- To make the Processor more "user gentle" you can add some checks to determine if the data required at this stage have been entered or loaded. If not you can then tell the user what is missing and what needs to be done to rectify the situation. To implement this you will need a data structure with flags to indicate whether the data have been entered or loaded. And you will need a corresponding table that contains the commands needed to enter or load the data.

  You can get very elaborate if you let the user indicate whether the data are to be loaded or entered at the keyboard. Then based on the user response you can search the open GAL Library for the proper dataset and load it, or jump to the proper subroutine for entering the data from the keyboard.

  You can even enter the realm of artificial intelligence (AI). You would develop some rules for what order the data are entered, what data are needed for what command, what commands perform what tasks, etc. Then with a simple forward chaining inference engine you can assist the user at any point in the analysis by telling them what usual comes next, what data are needed, how to get to some point in the analysis from where the user asks, etc.

- To enter a research mode or maximum flexibility mode you may wish to break the DBEM2 Processor into several independently executable Processors. A good starting point would be pre-processing -- all of the DEFINE commands, BUILD, GENERATE, SOLVE, and post-processing -- all of the PRINT commands. Now you could have other boundary elements that would be incorporated in the BUILD process. Or somewhat easier, you can replace the SOLVE Processor with a new solver, SOLVE_NEW, and compare the performance of the two solvers. Of course output data generated by the post-processing Processor can be stored on the database. Then you could develop plot Processors to display the data.

  You can be very ambitious and combine the boundary element method with a finite element code. Here you would need to develop some special matrix Processors to properly assemble the system matrix. Or be very brave and try a coupled solution procedure.

  Anyway, I hope you can see the unlimited potential of developing computational software in this mode. A common command language and a common database manager to unify the software is a very powerful paradigm.

  "From little acorns the mighty oak does grow."
11. An Example Problem

It is convenient to test DBEM2 on the same example problem used by Crouch and Starfield [4]. The problem concerns a unit-radius circular hole in an infinite body under uniaxial tension at infinity. The boundary element discretization for one-quarter of the hole is shown in Figure 11-1.

Both \( x = 0 \) and \( y = 0 \) are symmetry lines. The boundary contour is approximated by six straight-line segments, each of which consists of one boundary element. Two field point lines are chosen along portions of the \( x \) and \( y \) axes as shown in Figure 11-1(a).
An Example Problem

The input for this problem is prepared (with the text editor) in the form of a script command file:

```plaintext
open lib=circhole.gal

clear

def segments
    seg=1 b=1.0 e=.9659, .2588
    seg=2 b=.9659, .2588 e=.8660,.5000
    seg=3 b=.8660,.5000 e=.7071,.7071
    seg=4 b=.7071,.7071 e=.5000,.8660
    seg=5 b=.5000,.8660 e=.2588,.9659
    seg=6 b=.2588,.9659 e=0,1
store
end

def material
    em=7.E4 ; pr=0.2 ; store ; end


def symmetry
    xsym=0 ; ysym = 0 ; store ; end

def prestress
    sxx=100 ; store ; end

def field
    line=1 f=1.0 l=6.0 p=9
    line=2 f=0.1 l=0.6 p=9
store
end

pri seg ; pri mat ; pri bou ; pri symm ; pri pres ; pri field
build/store ; gen/store ; sol/store

pri res ; pri res/field
```

Note that there is no need for DEFINE ELEMENT input data because each segment contains only one boundary element, which is the default assumption.

Upon starting the DBEM2 processor, this file is inserted in the command stream through an ADD directive [3], §13.1. For example, under UNIX:

```plaintext
> dbem2
DBEM2> *add circhole.add
```

where circhole.add is the assumed name of the input file. The printed results should then be compared with those given on Appendix C of Crouch and Starfield [4].

Here is what you would see on your screen.

Tables initialized

<table>
<thead>
<tr>
<th>Boundary Segment Data</th>
</tr>
</thead>
<tbody>
<tr>
<td>Segm</td>
</tr>
<tr>
<td>1</td>
</tr>
<tr>
<td>2</td>
</tr>
<tr>
<td>3</td>
</tr>
<tr>
<td>4</td>
</tr>
</tbody>
</table>

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An Example Problem

Material Property Data
Elastic modulus: 7.000E+04
Poisson's ratio: 0.200

Boundary Conditions Data
<table>
<thead>
<tr>
<th>Segm</th>
<th>Given</th>
<th>Shear</th>
<th>Normal</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>2</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>3</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>4</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>5</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>6</td>
<td>SS and NS</td>
<td>0.</td>
<td>0.</td>
</tr>
</tbody>
</table>

Symmetry Data
Symmetry about axis X= 0.
and axis Y= 0.

Prestress (Initial Field Stresses) Data
Sigma_xx: 100.
Sigma_yy: 0.
Sigma_xy: 0.

Field Location Data
<table>
<thead>
<tr>
<th>Line</th>
<th>Int.Pts</th>
<th>x-first</th>
<th>y-first</th>
<th>x-last</th>
<th>y-last</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>9</td>
<td>1.000</td>
<td>0.</td>
<td>6.000</td>
<td>0.</td>
</tr>
<tr>
<td>2</td>
<td>9</td>
<td>0.</td>
<td>1.000</td>
<td>0.</td>
<td>6.000</td>
</tr>
</tbody>
</table>

Discrete model building completed: 6 boundary elements

Influence coefficient matrix & RHS vector generated
Discrete equations solved

Displacements and Stresses at Boundary Element Midpoints

<table>
<thead>
<tr>
<th>Elem</th>
<th>u_s</th>
<th>u_n</th>
<th>u_x</th>
<th>u_y</th>
<th>sig_s</th>
<th>sig_n</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>-4.768E-04</td>
<td>-2.649E-03</td>
<td>2.689E-03</td>
<td>-1.267E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>2</td>
<td>-1.302E-03</td>
<td>-2.177E-03</td>
<td>2.509E-03</td>
<td>-3.697E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>3</td>
<td>-1.778E-03</td>
<td>-1.359E-03</td>
<td>2.161E-03</td>
<td>-5.836E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>4</td>
<td>-1.778E-03</td>
<td>-4.142E-04</td>
<td>1.663E-03</td>
<td>-7.539E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>5</td>
<td>-1.302E-03</td>
<td>4.038E-04</td>
<td>1.048E-03</td>
<td>-8.712E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
<tr>
<td>6</td>
<td>-4.767E-04</td>
<td>8.761E-04</td>
<td>3.582E-04</td>
<td>-9.309E-04</td>
<td>0.</td>
<td>0.</td>
</tr>
</tbody>
</table>

Displacements and Stresses at Specified Field Points

<table>
<thead>
<tr>
<th>Lin</th>
<th>x</th>
<th>y</th>
<th>u_x</th>
<th>u_y</th>
<th>sig_xx</th>
<th>sig_yy</th>
<th>sig_xy</th>
</tr>
</thead>
</table>

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An Example Problem

<table>
<thead>
<tr>
<th>Lin</th>
<th>x</th>
<th>y</th>
<th>u_x</th>
<th>u_y</th>
<th>sig_xx</th>
<th>sig_yy</th>
<th>sig_xy</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1.000</td>
<td>0.000</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>1.500</td>
<td>0.000</td>
<td>2.116E-03</td>
<td>5.483E-11</td>
<td>64.8</td>
<td>-42.0</td>
<td>-3.078E-06</td>
</tr>
<tr>
<td>1</td>
<td>2.000</td>
<td>0.000</td>
<td>1.668E-03</td>
<td>1.026E-10</td>
<td>72.3</td>
<td>-18.1</td>
<td>-2.897E-06</td>
</tr>
<tr>
<td>1</td>
<td>2.500</td>
<td>0.000</td>
<td>1.365E-03</td>
<td>1.255E-10</td>
<td>80.2</td>
<td>-10.2</td>
<td>-3.334E-07</td>
</tr>
<tr>
<td>1</td>
<td>3.000</td>
<td>0.000</td>
<td>1.151E-03</td>
<td>2.173E-10</td>
<td>85.5</td>
<td>-6.54</td>
<td>-1.722E-09</td>
</tr>
<tr>
<td>1</td>
<td>3.500</td>
<td>0.000</td>
<td>9.935E-04</td>
<td>2.065E-10</td>
<td>89.0</td>
<td>-4.57</td>
<td>2.339E-06</td>
</tr>
<tr>
<td>1</td>
<td>4.000</td>
<td>0.000</td>
<td>8.733E-04</td>
<td>1.669E-10</td>
<td>91.4</td>
<td>-3.38</td>
<td>1.921E-06</td>
</tr>
<tr>
<td>1</td>
<td>4.500</td>
<td>0.000</td>
<td>7.787E-04</td>
<td>3.464E-11</td>
<td>93.1</td>
<td>-2.61</td>
<td>-2.666E-07</td>
</tr>
<tr>
<td>1</td>
<td>5.000</td>
<td>0.000</td>
<td>7.024E-04</td>
<td>1.132E-10</td>
<td>94.4</td>
<td>-2.08</td>
<td>1.784E-07</td>
</tr>
<tr>
<td>1</td>
<td>5.500</td>
<td>0.000</td>
<td>6.396E-04</td>
<td>4.157E-10</td>
<td>95.3</td>
<td>-1.70</td>
<td>-1.803E-07</td>
</tr>
<tr>
<td>1</td>
<td>6.000</td>
<td>0.000</td>
<td>5.870E-04</td>
<td>7.854E-11</td>
<td>96.0</td>
<td>-1.41</td>
<td>2.896E-06</td>
</tr>
</tbody>
</table>

Point is too close to boundary

Now let's take a look at the database that was generated during this run. All of the commands that begin with * are directives. The directives are described in [3].

DBEM2> *toc

Library 1 File: circhole.gal
Form: GAL82 File size: 1973 words No. of Datasets: 9

Here we see the table of contents (toc) for the GAL Library that was generated during the run. You can see that the file name is circhole.gal, the size of the file, there are 9 datasets, and other information, like the date and time the data were stored.
An Example Problem

To see the second level of the database, the records, the *print rat or *rat directive [3], §49.1, can be used. Let’s take a look at some of the record structure. You may wish to compare this output to the second diagram in §10.0 or compare with the discussion in §3.

DBEM2> *rat 1,1

Record Table of dataset SEGMENT
Key     L_cyc H_cyc Type Log_size
NUMEL   1   6   I   1
SEGDEF  1   6   I   1
XBEG    1   6   S   1
XEND    1   6   S   1
YBEG    1   6   S   1
YEND    1   6   S   1

DBEM2> *rat 1,2

Record Table of dataset MATERIAL
Key     L_cyc H_cyc Type Log_size
EM      0   0   S   1
PR      0   0   S   1

DBEM2> *rat 1,7

Record Table of dataset COEFF
Key     L_cyc H_cyc Type Log_size
C       1   12  S   12

The first number after the *rat is the GAL Library id, which is printed as the Library in the *toc [3], §49.1, output shown above. The second number is the dataset sequence number, also shown in the table of contents output. The *rat output tells us the names of the records, the Key; the number of records, from L_cyc to H_cyc; the data type, I is integer — S is floating point; and the logical size of each record. For example, the record key NUMEL in the dataset SEGMENT has six records containing one integer in each record. The record keys in the dataset MATERIAL are a bit different in that each record only has one floating point number, so there are no cycles. The record key C in the dataset COEFF has 12 records each containing 12 floating point numbers. This is the square system coefficient matrix.

We can also look at the data stored within each record key.

DBEM2> *print rec 1,1.NUMEL 1:6
Record NUMEL.1 of dataset SEGMENT
1: 1
Record NUMEL.2 of dataset SEGMENT
1: 1
Record NUMEL.3 of dataset SEGMENT
1: 1
Record NUMEL.4 of dataset SEGMENT

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An Example Problem

Here we have used the *print record directive \[3\], \$49.2 & \$49.3, to show: 1) the values in the six records in the record key NUMEL, 2) the values of the elastic modulus, EM, and Poisson's ratio, PR, in the dataset MATERIAL, and 3) the seventh column of the system coefficient matrix, C, which is stored in the dataset COEFF.

Now to exit the DBEM2 Processor we type stop to produce the following on our screen:

```
DBEM2> stop
<DM> CLOSE, Ldi: 1, File: circhole.gal
Hope you enjoyed the ride!
<CL> PIS exhausted
ENDRUN called by CLIP
```

A few days later we decide we would like to solve the same problem, but we wish to increase the number of elements to 2 for each segment to see how much this improves the results. So here is what it looks like.

```
aml.9: 17 > dbem2
DBEM2> clear
Tables initialized
DBEM2> open lib=circhole.gal
<DM> OPEN, Ldi: 1, File: circhole.gal , Attr: old, Block I/O
DBEM2> toc

*************+-----------------------------------------------+*************
+ Library | File: circhole.gal                                |
+ Form:   | GAL82 File size: 1973 words No. of Datasets: 9     +
*************+-----------------------------------------------+*************
Seq#  | Date  | Time  | Lk | Records | Processor | Dataset name  
 1    | 06:01:88 | 20:07:20 | 0   | 36      | DBEM2      | SEGMENT      
 2    | 06:01:88 | 20:07:20 | 0   | 2       | DBEM2      | MATERIAL     
 3    | 06:01:88 | 20:07:20 | 0   | 3       | DBEM2      | SYMMETRY     
```
An Example Problem

```
4 06:01:88 20:07:20 0 3 DBEM2 PRESTRESS
5 06:01:88 20:07:20 0 12 DBEM2 FIELD
6 06:01:88 20:07:20 0 49 DBEM2 ELEMENT
7 06:01:88 20:07:20 0 12 DBEM2 COEFF
8 06:01:88 20:07:20 0 12 DBEM2 RHS
9 06:01:88 20:07:20 0 12 DBEM2 SOLUTION

DBEM2> def seg ; load ; end
DBEM2> def mat ; load ; end
DBEM2> def sym ; load ; end
DBEM2> def pres ; load ; end
DBEM2> def field ; load ; end
DBEM2> pri seg

Boundary Segment Data
Segm Elements Xbeg Ybeg Xend Yend
1   1 1.000 0.000 0.9659 0.2588
2   1 0.9659 0.2588 0.8660 0.5000
3   1 0.8660 0.5000 0.7071 0.7071
4   1 0.7071 0.7071 0.5000 0.8660
5   1 0.5000 0.8660 0.2588 0.9659
6   1 0.2588 0.9659 0.000 1.000

DBEM2> pri pres

Prestress (Initial Field Stresses) Data
Sigma_xx: 100. Sigma_yy: 0. Sigma_xy: 0.

DBEM2> pri mat

Material Property Data
Elastic modulus: 7.000E+04
Poisson's ratio: 0.200

DBEM2> help def elem
<DBEM2>
DEFINE
ELEMENTS

The DEFINE ELEMENTS command introduces subordinate commands that specify into how many boundary elements segments are to be subdivided. These commands have the form:

```
SEG = iseg1, ... isegk
ELEM = ne1, ... nek
```

This specifies that segment iseg1 is to be subdivided into ne1 (ge 1) boundary elements, segment iseg2 into ne2 elements, and so on. Enter the STORE command to store this data in the NUMEL record in the SEGMENT dataset. Enter the

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LOAD command to load previously defined data from the NUMEL record in the SEGMENT dataset. Terminate these commands with an END command.

```
DBEM2> def elem
Element data> seg = 1:6 elem = 2,2,2,2,2
Element data> end
DBEM2> pri seg

Boundary Segment Data
Segm Elements Xbeg Ybeg Xend Yend
1 2  1.000 0.  0.9659 0.2588
2 2  0.9659 0.2588 0.8660 0.5000
3 2  0.8660 0.5000 0.7071 0.7071
4 2  0.7071 0.7071 0.5000 0.8660
5 2  0.5000 0.8660 0.2588 0.9659
6 2  0.2588 0.9659 0.  1.000
```

```
DBEM2> build
Discrete model building completed: 12 boundary elements

DBEM2> generate
Influence coefficient matrix & RHS vector generated

DBEM2> solve
Discrete equations solved

DBEM2> pri res
Displacements and Stresses at Boundary Element Midpoints
Elem u_s u_n u_x u_y sig_s sig_n
 1 -4.791E-04 -2.755E-03 2.794E-03 -1.152E-04 0.  0.
 2 -1.288E-03 -2.289E-03 2.438E-03 -9.782E-04 0.  0.
 3 -1.769E-03 -1.483E-03 2.047E-03 -1.067E-03 0.  0.
 4 -1.743E-03 -5.670E-04 1.191E-03 -1.394E-03 0.  0.
 5 -1.266E-03 1.977E-04 6.131E-04 -1.124E-03 0.  0.
 6 -3.329E-04 4.904E-04 -1.865E-04 -5.626E-04 0.  0.
 7 -1.523E-03 -5.829E-04 1.563E-03 -4.646E-04 0.  0.
 8 -1.513E-03 -2.601E-04 1.359E-03 -7.147E-04 0.  0.
 9 -1.306E-03 3.710E-04 1.064E-03 -8.424E-04 0.  0.
10 -1.107E-03 5.584E-04 8.093E-04 -9.396E-04 0.  0.
12 -3.178E-04 9.855E-04 1.864E-04 -1.019E-03 0.  0.
```

```
DBEM2> pri res /field
Displacements and Stresses at Specified Field Points
Lin x y u_x u_y sig_xx sig_yy sig_xy
 1 1.000 0.000 Point is too close to boundary
 1 1.500 0.000 1.935E-03 1.465E-11 55.5 -53.9 -1.108E-06
```
An Example Problem

<table>
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<th>Lin</th>
<th>x</th>
<th>y</th>
<th>u_x</th>
<th>u_y</th>
<th>sig_xx</th>
<th>sig_yy</th>
<th>sig_xy</th>
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<td>0.000</td>
<td>1.432E-03</td>
<td>-5.107E-11</td>
<td>70.7</td>
<td>-17.6</td>
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<td>1.142E-03</td>
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<td>-3.628E-06</td>
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<td>0.000</td>
<td>9.520E-04</td>
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<td>8.165E-04</td>
<td>5.708E-10</td>
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<td>6.360E-04</td>
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<td>5.208E-04</td>
<td>7.305E-10</td>
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<td>4.776E-04</td>
<td>5.825E-10</td>
<td>96.1</td>
<td>-0.85</td>
<td>2.337E-06</td>
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</table>

Lin x y u_x u_y sig_xx sig_yy sig_xy
2 0.000 1.000 Point is too close to boundary
2 0.000 1.500 -3.197E-11 -5.243E-03 160. 24.8 4.456E-07
2 0.000 2.000 7.015E-11 -9.904E-03 126. 18.4 -7.870E-08
2 0.000 2.500 -1.200E-10 -1.347E-02 114. 13.0 -1.185E-06
2 0.000 3.000 1.169E-10 -1.640E-02 109. 9.48 1.955E-06
2 0.000 3.500 9.394E-11 -1.889E-02 107. 7.17 -7.394E-07
2 0.000 4.000 -3.967E-10 -2.107E-02 105. 5.60 -2.295E-06
2 0.000 4.500 3.040E-10 -2.301E-02 104. 4.48 3.884E-07
2 0.000 5.000 -2.895E-10 -2.474E-02 103. 3.66 -1.996E-06
2 0.000 5.500 -7.609E-11 -2.632E-02 103. 3.04 -6.769E-07
2 0.000 6.000 5.487E-10 -2.777E-02 102. 2.57 1.350E-06

DBEM2> stop
<DM> CLOSE, Ldi: 1, File: circhole.gal
Hope you enjoyed the ride!
<CL> PHS exhausted
ENDRUN called by CLIP

This interactive session starts off with the usual clear and open commands. Then to make sure that I have the correct GAL Library, the $toc [3], §19.1, is used. Everything looks okay, so the problem definition commands to define segments, materials, symmetry, prestress, and field locations are issued with the subcommand load to load these data from the GAL Library. The print segments command is used to look at the old segment data that were just loaded—gives you a warm feeling to see that the data are really there. Just for insurance I check the prestress data and the material data. Now, I want to enter data to use 2 elements per segment, but I can't remember the proper syntax, so I use the help define elements command to get the on-line help (see [3], Appendix H for a discussion of help files). Now, I enter the data and check it with another print segments command. Note that, I now have 2 elements for each segment. Then the three number crunching modules, build, generate, and solve, are brought into action to obtain the new solution. Finally, the new results are printed. Note that, I did not store any of the new problem data. Only a simple change was made to compare answers. If big changes were made, I would have stored the new data.

Experiment with your own changes to this problem. Then try some new problems. Enjoy!

Revised 10/14/1988  Developer's Tutorial for the CSM Testbed Architecture  11-9
An Example Problem

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12. References


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**Application Developer's Tutorial for the CSM Testbed Architecture**

**Abstract**
This tutorial serves as an illustration of the use of the programmer interface of the CSM Testbed Architecture (NICE). It presents a complete, but simple, introduction to using both the GAL-DBM (Global Access Library-Database Manager) and CLIP (Command Language Interface Program) to write a NICE processor. Familiarity with the CSM Testbed architecture is required.

**Key Words**
- Structural analysis software
- Command language interface software
- Data management software

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