Prototyping Distributed Simulation Networks

Dennis L. Doubleday
Software Engineering Institute
Abstract

Dorra is a declarative language designed to support application-level programming. In this paper we illustrate the use of Dorra to describe a simple distributed application: a simulation of a collection of networked vehicle simulators. We show how the language is used to describe the application, its components and structure, and how the runtime executive provides for the execution of the application.

1. Programming at the Application-Level

Many distributed applications consist of large-grained tasks or programs, instantiated as processes, running on possibly separate processors and communicating with each other by sending messages of different types.

Since the patterns of communication between the processes can vary over time and the speeds of the individual processors can differ widely, the developers may need explicit control over the allocation of processors to processes in order to meet performance or reliability requirements. Processors are not the only critical resource. The resources that must be allocated also include communication links and message queues. We call this network of various processor types, links, and queues a heterogeneous machine.

Currently, users of a heterogeneous machine network follow the same pattern of program development as users of conventional processors: Programmers write individual tasks as separate programs, in the different programming languages (e.g., C, Lisp, Ada) supported by the processors, and then hard code the allocation of resources to their application by explicitly assigning specific programs to run on specific processors at specific times. This coupling between the component programs and the built-in knowledge about the structure of the application and the allocation of resources often prevents the reuse of the programs in other applications or environments. Modification of the application during development is often expensive, time-consuming, and error-prone. The problem is compounded if the application must be modified while running in order to deal with faults or mode changes. We claim that developing distributed applications for a heterogeneous machine is qualitatively different from developing programs for conventional processors. It requires different kinds of languages, tools, runtime support, and methodologies. In this paper we address some of these issues by presenting a language, Dorra. We briefly describe the language and its distributed runtime support environment and then present, as an example distributed application, a simple simulation of a network of vehicle simulators.

The rest of this paper is organized as follows. Section 2 briefly describes the Dorra language and runtime environment. Section 3 discusses the problem we are attempting to address in the realm of
networked simulation devices. Section 4 describes the work we have done to date toward that end.

2. Introduction to Durra
Durra [2] is a language designed to support the development of distributed, large-grained concurrent applications running on heterogeneous machine networks. A Durra application description consists of a set of task descriptions and type declarations that prescribe a way to manage the resources of the network. The application description describes the tasks to be instantiated and executed as concurrent processes, the types of data to be exchanged by the processes, and the intermediate queues required to store the data as they move from producer to consumer processes.

2.1. The Durra Language
Task descriptions are the building blocks for applications. A task description includes the following information (Figure 1): (1) its interface to other tasks (ports); (2) its attributes; (3) its functional and timing behavior; and (4) its internal structure, thereby allowing for hierarchical task descriptions.

```
task task-name
  ports
    port-declarations
  attributes
    attribute-value-pairs
  behavior
    functional specification
    timing specification
  structure
    process-declarations
    bind-declarations
    queue-declarations
    reconfiguration-statements
end task-name
```

---

The interface information declares the ports of the processes instantiated from the task. A port declaration specifies the direction and type of data moving through the port. An in port takes input data from a queue; an out port deposits data into a queue:

```
ports
  inl: in heads;
  outl, out2: out tails;
```

The attribute information specifies miscellaneous properties of a task. Attributes are a means of indicating pragmas or hints to the compiler and/or runtime executive. In a task description, the developer of the task lists the actual value of a property; in a task selection, the user of a task lists the desired value of the property. Example attributes include author, version number, programming language, file name, and processor type:

```
attributes
  author = "jmw";
  implementation = "program_name";
  Queue_Size = 25;
```
The behavioral information specifies functional and timing properties of the task. The functional information part of a task description consists of a pre-condition on what is required to be true of the data coming through the input ports, and a post-condition on what is guaranteed to be true of the data going out through the output ports. The timing expression describes the behavior of the task in terms of the operations it performs on its input and output ports. For additional information about the syntax and semantics of the functional and timing behavior description, see the Durra reference manual [1].

The structural information defines a process-queue graph and possible dynamic reconfiguration of the graph.

A process declaration of the form

```
process_name : task task_selection
```

creates a process as an instance of the specified task. Since a given task (e.g., convolution) might have a number of different implementations that differ along different dimensions such as algorithm used, code version, performance, or processor type, the task selection in a process declaration specifies the desirable features of a suitable implementation. The presence of task selections within task descriptions provides direct linguistic support for hierarchically structured tasks.

A queue declaration of the form

```
queue_name [queue_size]: port_name_1 > data_transformation > port_name_2
```

creates a queue through which data flow from an output port of a process (port_name_1) into the input port of another process (port_name_2). Data transformations are operations applied to data coming from a source port before they are delivered to a destination port.

A port binding of the form

```
task_port = process_port
```

maps a port on an internal process to a port defining the external interface of a compound task.

A reconfiguration statement of the form

```
If condition then
    remove process-names
    process process-declarations
    queues queue-declarations
end If;
```

is a directive to the executive. It is used to specify changes in the current structure of the application (i.e., process-queue graph) and the conditions under which these changes take effect. Typically, a number of existing processes and queues are replaced by new processes and queues, which are then connected to the remainder of the original graph. The reconfiguration predicate is a Boolean expression involving time values, queue sizes, and other information available to the executive at runtime.

2.2. The Durra Runtime Environment

There are two classes of active components in the Durra runtime environment: the application processes and the Durra executives. As shown in Figure 2, an instance of the executive runs on each processor while the processes are distributed across the processors in the system.
The executives interpret the resource allocation commands produced by the Durra compiler, monitor reconfiguration conditions, and implement the necessary changes in the application structure.

The component processes making up a Durra application are instances of independent tasks (programs) that can be written in any language for which a Durra interface has been provided (currently, there are Durra interfaces for both C and Ada). The Durra interface is a collection of procedures that provide communication and control primitives. The component processes use the interface to communicate with the Durra executives and, indirectly, with other application processes. For a more detailed discussion of the Durra runtime environment, see [3].

3. Distributed Simulation Networks
The development of large networks of heterogeneous simulation and training devices often presents problems related to the performance and interconnectivity of the network components. There is a need to evaluate various design alternatives before committing to a specific implementation. Problems arise in several areas:

- Multiple protocols. Cooperating devices are often written using different communication protocols because they rely on predetermined standards or technologies. When communicating devices use different protocols, it is necessary to translate messages in a way that is transparent to the communicating agents. This message translation consumes time and reduces performance.
Multiple levels of fidelity. When developing hierarchical networks of simulation and training devices, it is often the case that the time scales (i.e., granularity), amount of data, and level of detail in the data are not compatible between levels or devices. Thus, there is a need to filter (i.e., reduce) data moving up in the hierarchy and to pad (i.e., augment) data moving down the hierarchy. This is a different type of 'translation' from the protocol translation described above. The translating programs in this case need to have a thorough understanding of the application to compensate for the mismatch in the levels of detail.

Multiple technologies. When connecting devices that use different hardware technology, the developers of the distributed application need to compensate for differences in speed, performance, and fault-tolerance requirements.

This collection of problems is just an illustration of the issues that must be addressed by the developers before implementing the network. A useful technique is to develop prototypes using emulators of the component software and hardware devices. The emulators are easier to implement than the real devices and can more easily be reconfigured into alternative structures. Experiments can be conducted under various load conditions and measurements of performance can be derived from these experiments.

4. Using Durra to Prototype Simulation Networks

We are using Durra to develop a tool for testing and evaluating various network configurations. We are implementing the tool as a distributed application consisting of clusters of emulators. These emulators are responsible for interpreting specifications of hypothetical application tasks. We use the Durra language to describe the various components of the system, their ports and message queues, and the types of messages exchanged between components. We use the Durra runtime environment to execute the application and perform dynamic reconfigurations of the application, to emulate mode changes, and to evaluate their impact on performance.

The final version of our tool will include at least four types of emulators:

1. **Generic simulation device emulators**: These programs will mimic the I/O behavior of a generic networked simulation device. Scripts specifying the behavior of the emulated device(s) will be developed. Differences in I/O behavior between different types of simulation devices can be emulated through variations in these scripts. The initial scripts consist simply of position updates and timing instructions. Eventually they should be more representative of actual networked simulation sessions; this could be accomplished by adaptation of I/O logs of an actual simulation session.

2. **LAN emulators**: These emulators will model communications delay in the network (e.g., token ring delay). This kind of emulation can likely be accomplished via buffer tasks in the Durra runtime, which would mean that no executable version of these emulators need be developed.

3. **Intelligent gateway emulators**: These programs will model the effect of various message-filtering and protocol translation techniques on the networked simulation’s use of processor and communications resources.

4. **Console emulator**: This program will provide an interactive user interface to the simulation environment, allowing the experimenter to change emulation parameters, inject faults, and collect data.
4.1. Example: A Simple Network Specification

In this section we present a Durra specification of a simple network of simulators. In this example, we instantiate a user console and two LAN emulators, each consisting of a group of three simulators and one gateway process. The reader should note that there is nothing special about this configuration—another version consisting of some other grouping could just as easily have been constructed from the same primitive building blocks.

The following is the Durra description of the message type used for communications between the application components. The message type description is purposely a very general one. A generic description of the message type allows us in the actual implementation of the type to use a variant record to represent both simulator position updates and command messages and easily combine both types of messages in a single data stream.

```plaintext
type message is array of byte;
```

At the lowest level of the structure we have the descriptions of the primitive tasks, the `simulator`, the `gateway`, and the `console`. The `simulator` task has one output port, through which it emits its position updates, and one input port, through which it receives position updates and user commands. The `gateway` task has one input port and two output ports; port `to_wan` sends messages outside the LAN and port `to_lan` distributes remote messages to the simulators in its LAN. The `console` task is the application user's interface to the tool; it accepts a set of user commands and forwards them to the `gateway` task for each LAN in the configuration. The `gateways` may in turn forward those messages to the simulators in their respective LANs if the nature of the command requires it.

```plaintext
task simulator
    ports
        inl : in message;
        outl : out message;
    attributes
        version = "2";
        implementation = "simulator";
end simulator;

task gateway
    ports
        inl : in message;
        to_lan : out message;
        to_wan : out message;
    attributes
        version = "2";
        implementation = "gateway";
end gateway;

task console
    ports
        to_lan : out message;
    attributes
        xwindow = "-geom 80x24+0+0 -title CONSOLE";
        implementation = "console";
end console;
```

The Durra task `lan` encapsulates the internal structure of the LAN itself. This instantiation of a LAN includes one `gateway` task and three `simulator` tasks, as well as three built-in Durra buffer tasks. The
buffer tasks implement the routing of message traffic between the component tasks of the LAN. Task gate_merge merges local and remote messages intended for the local gateway. Task gate_mb merges messages from the local simulators and then distributes them to both the gate_merge task and the lan_mb task. The lan_mb task merges those local messages with the remote messages forwarded from the gateway and distributes them all to each of the local simulators. Note that, given this structure, each simulator will receive its own updates; these can either be ignored by the simulator or used as a check to ensure that its own updates are being distributed properly.

**task lan**

ports

in1 : in message;
out1 : out message;

structure

process

gate : task gateway attributes version = "2"; end gateway;
sim1, sim2, sim3 :
task simulator attributes version = "2"; end simulator;
gate_merge : task merge

ports

from_lan, from_wan : in message;
to_gate : out message;
attributes mode = fifo;
end merge;

gate_mb : task merge_broadcast

ports

froml, from2, from3 : in message;
to_gate, to_lan : out message;
attributes mode = fifo;
end merge_broadcast;

lan_mb : task merge_broadcast

ports

from_gate, from_lan : in message;
to1, to2, to3 : out message;
attributes mode = fifo;
end merge_broadcast;

queues

ggate_in[10] : gate_merge.to_gate >> gate.in1;
ggate_out[10] : gate.to_lan >> lan_mb.from_gate;
gsim1_in[10] : lan_mb.to1 >> sim1.in1;
gsim2_in[10] : lan_mb.to2 >> sim2.in1;
gsim3_in[10] : lan_mb.to3 >> sim3.in1;
gsim1_out[10] : sim1.out1 >> gate_mb.from1;
gsim2_out[10] : sim2.out1 >> gate_mb.from2;
gsim3_out[10] : sim3.out1 >> gate_mb.from3;
gmb_to_gate[10] : gate_mb.to_gate >> gate_merge.from_lan;
gmb_to_lan[10] : gate_mb.to_lan >> lan_mb.from_lan;

bind

in1 = gate_merge.from_wan;
out1 = gate.to_wan;

end lan;

At the highest level of abstraction, the Durra task Internet provides the view of the application as a console process controlling two connected, but independent, local area networks. These LAN simulators may be distributed to separate physical processors. Figure 3 shows a graphical view of
the structure of the application.

```
task internet
structure
process
lan1: task lan attributes processor = net1; end lan;
lan2: task lan attributes processor = net2; end lan;
uc : task console attributes version = "xterm"; end console;
uc_b : task broadcast
    ports
    from_uc : in message;
    to_lan1, to_lan2 : out message;
end broadcast;
lan1_m, lan2_m :
    task merge
    ports
    from_uc, from_lan : in message;
    to_lan : out message;
    attributes mode = fifo;
end merge;

queues
quctob : uc.to_lan >> uc_b.from_uc;
qucbto1 : uc_b.to_lan1 >> lan1_m.from_uc;
qucbto2 : uc_b.to_lan2 >> lan2_m.from_uc;
q1tom[10] : lan1.out1 >> lan2_m.from_lan;
q2tom[10] : lan2.out1 >> lan1_m.from_lan;
qmtol[10] : lan2_m.to_lan >> lan2.in1;
qmtol2[10] : lan1_m.to_lan >> lan1.in1;
end internet;
```

Only three of the aforementioned Durra tasks, the simulator, the gateway, and the console have actual implementations associated with them. The lan task's behavior is defined constructively from the behavior of the simulator and the gateway, the three buffer tasks (whose behavior is implemented in the Durra executive), and the connections between them all. Similarly, the behavior of the internet task derives from the connections between its components, the two instantiations of the lan task and the console.

5. Conclusions

Application-level programming, as implemented by Durra, separates the structure of an application from its behavior. This separation provides developers with control over the evolution of an application during application development as well as during application execution. During development, an application evolves as the requirements of the application are better understood or as they change.

This evolution takes the form of changes in the application description, modifying task selection templates to retrieve alternative task implementations from the library, and connecting these implementations in different ways to reflect alternative designs. During execution, an application evolves through mode changes or in response to faults. This evolution takes the form of conditional,
Figure 3: Structure of the Application

dynamic reconfigurations, removal of processes and queues, and instantiation of new processes and queues without affecting the remaining components. This approach to application-level programming is similar in spirit to the constructive approach of CONIC [4]. We illustrated this method for developing distributed applications by describing the implementation of a simple prototyping tool for modelling various configurations of networked simulators. We wrote Dura task and application descriptions and used them to control the evolution of the application, both during the development and during the execution.
References

*Durra: A Task-Level Description Language.*  

*Durra: A Task-Level Description Language Reference Manual (Version 2).*  

*Developing Fault-Tolerant Distributed Systems.*  

*A Model for Change Management.*  
Session 3

Software Reuse

Chair: Robert Angier, IBM Corp.
Research Directions in Software Reuse

November 8, 1990

Will Tracz

MD 0210
IBM Federal Sector Division
Owego, NY 13827
(607) 751-2169
net: OWEGO@IBM.COM

Unclassified
Session 3

Software Reuse

Chair: Robert Angier, IBM Corp.
Research Directions in Software Reuse

November 8, 1990

Will Tracz

MD 0210
IBM Federal Sector Division
Owego, NY 13827
(607) 751-2169
net: OWEGO@IBM.COM

Unclassified
"Currently, software is put together one statement at a time. What we need is to put software together one component at a time." – Barry Boehm, at the Domain Specific Software Architecture (DSSA) Workshop, July 11-12, 1990.

Megaprogramming Motivation

"Megaprogramming is the type of thing you can go into a 3-star general's office and use to explain what DARPA is going to do for them to make their software less expensive and have better quality." – Barry Boehm, at the ISTO Software Technology Community Meeting, June 27-29, 1990.

Software productivity improvements in the past have been accidental because they allow us to "work faster". DARPA wants people to "work smarter" or to avoid work altogether." – Barry Boehm, at the Domain Specific Software Architecture (DSSA) Workshop, July 11-12, 1990.

Megaprogramming Vision

- Megaprogramming is a "giant step" toward increasing
  - "development productivity,
  - maintenance productivity,
  - reliability,
  - availability,
  - security,
  - portability,
  - interoperability and
  - operational capability ."

- Megaprogramming will incorporate proven, well-defined components whose quality will evolve.
- Megaprogramming requires the modification of the traditional software development process.
- Domain-specific software architectures need to be defined and implemented with open interfaces according to software composition principles and open interface specifications.
- Additional environmental capabilities are needed to provide software understanding.

Megaprogramming Software Team Goal

To create an environment to:

1. "manage systems as configurations of components, interfaces, specifications, etc.,
2. increase the scale of units of software construction (to modules), and
3. increase the range of scales of units of software interchange (algorithms to subsystems) ."
Key Elements of Megaprogramming Software Team

- **Component sources** — currently, components under consideration are from reuse libraries (e.g., SIMTEL2 or RAPID) or COTS (Commercial Off-The-Shelf) software (e.g., GRACE or Booch components). Application generator technology is desirable to provide for adaptable modules. Re-engineered components (e.g., CAMP) could provide additional resources.

- **Interface definitions** — currently, there exists an ad hoc standard consisting of Ada package specifications and informal documentation. It is desirable to develop a Module Interconnect Formalism (MIF) with hidden implementations supported by formal analysis and validation tools.

- **System documentation** — currently, simple hypertext systems are supporting the textual documentation associated with software components. It is desirable to create a repository-based, hypermedia environment that provides traceability between artifacts and supports the capture, query, and navigation of domain knowledge.

**Megaprogramming Resources**

- STARS (Software Technology for Adaptable Reliable Systems) SEE (Software Engineering Environment)
- Arcadia
- CPS/CPL (Common Prototyping System/Common Prototyping Language)
- DSSA (Domain Specific Software Architectures)
- POB (Persistent Object Bases)
- SWU (Software Understanding)
- REE (Re-Engineering)

Interface and architecture codification will be supported by a Module Interconnect Formalism (MIF), which is an outgrowth of the CPS/CPL program.

**Goal of MIF**

To adequately describe a software component such that its selection and use can be accomplished without looking at its implementation.

**Component Interface**

- entry points,
- type definitions
- data formats (e.g. Ada package specification),
- a description of its functionality,
- side effects,
- performance expectations,
- degree and kind of assurance of consistency between specification and implementation (reliability), and
- appropriate test cases.
The design record will provide a "common data structure for system documentation and libraries".

The suggested data elements in a design record include:

- code,
- test cases,
- library and DSSA links,
- design structure,
- access rights,
- configuration and version data,
- hypertext paths,
- metric data,
- requirement specification fragments,
- PDL texts,
- interface and architecture specifications,
- design rationale,
- catalog information, and
- search points.

Megaprogramming Software Interchange


Megaprogramming Software Interchange Goal

To "enable wide-area commerce in software components."

Elements of Megaprogramming Software Interchange

- Conventionalization — currently, conventions are emerging. It is desirable to create a cooperative decision and consensus mechanism that supports adaptable, multi-configuration libraries, which present a standard search capability.

- Repository/Inventory— currently, repositories support code storage only. It is desirable to retain, assess, and validate other software assets such as architectures, test cases, specifications, designs, and design rationales.

- Exchange/Brokerage — current intellectual property rights and government acquisition regulations are stifling a software component industry. It is desirable to populate certain application domains (via DSSA) and to support the creation of an electronic software component commerce by
  - defining mechanisms for access control,
  - authentication/certification, and
  - establishing composition conventions.

Realities of Reuse Workshop

January 4-5 1990
Syracuse, NY

The goal of the workshop was to

"...serve as a forum for sharing practical experiences and methodologies

- for specifying and designing software for reuse,
- for defining the level and kinds of components that can be reused, and
- for incorporating reuse philosophies into organizations."
Highlights

Software Reuse: Representing a Reusable Software Collection
William Frakes, Software Productivity Consortium
- IR approach is the best way to go about organizing a library.
- Other approaches (keyword, faceted, semantic net, hypertext) require significant amounts of effort to set up and to catalog.

Realities of Language Support for Reuse: What we desire - What we have.
Larry Latour, University of Maine
- Code and type inheritance
- Parameterization
- Granularity of change
- Algorithm parameterization.

Highlights

Library-Based Software Design Methodology
Devid Musser, RPI
- The following are myths:
  1. Generic software is not efficient.
  2. Generic software is hard to find, and
  3. Software libraries only address the implementation level.
- Rationale:
  1. Algorithms can be more complex and efficient than any simple ones that a programmer would tend to write from scratch.
  2. Library can be organized into a semantic net that a user could easily navigate in find what was needed.
  3. 80% of the effort to build a library is writing the specifications that could be reused at high level design time.

Highlights

Reusable Specifications for Requirements Prototyping and System Construction
- Prolo system that ISSI built for RADC.
- Graphical input language for drawing data flow diagrams, then simulating them (if the contents of the nodes is real code).
- One can also watch the data flow nodes fire.

Designing for Reuse: Is Ada Class Conscious?
Sholom Cohen, Software Engineering Institute
- Feature Analysis
- Commonality Analysis to develop a generic architectures.

Highlights

THIRD ANNUAL WORKSHOP:
METHODS & TOOLS FOR REUSE

June 13-15 1990
Syracuse, NY

Highlights

- If you are not teaching software reuse, you are not teaching software engineering (Bob Cook - University of Virginia)
- The (throw everything into a) "Bag" approach was the style of software reuse in the 80's, the "Generic Architecture" approach is the style for the 90's.
- "Cloning" (a new-to-me term) is a form of unplanned reuse (salvaging) popular at HP and other companies.
- What is needed to stimulate software reuse are handbooks that describe the architectures of applications along with their design rationale.
- GOTO's were found bad in the 70's for the same reason that Top Down Decomposition will be found bad in the 90's - failure to modularize completely.
A good interface specification has enough information so the (re-) user doesn't have to look at the code to figure out what it does and how to use it.

One (large) problem that people have failed to realize is that software reuse doesn't stop at retrieval.

Data flow diagrams provide too much information to be included in the functional specification of a reusable software component.

Domain Analysis research projects are actively being addressed at TRW, Bell Labs, UNISYS, ESPRIT, Magnavox, CONTEL, MCC and SPS.

SPS (Software Productivity Solutions) speculated that in 6 years they have increased their programmer productivity an order of magnitude through

1. simple black box reuse (function libraries)
2. parameterized black box reuse (Ada generics)
3. large component reuse (modules/Ada packages)
4. inheritance (required object-oriented programming language)
5. parameterized application generators

NOTE: they indicated the switch to OOPL was the greatest facilitator of reuse.

Best malaprop: "Generics are something you use when you can't afford the name brand."

Paper Summaries

KAPTUR: KNOWLEDGE ACQUISITION FOR PRESERVATION OF TRADEOFFS AND UNDERLYING RATIONALES
Sidney C. Ballin, CTA INCORPORATED
- Roll-your-own hypertext system for capturing design decisions.
- An impressive domain analysis case study in tools to support reuse.

REUSE OF SOFTWARE KNOWLEDGE: A PROGRESS REPORT
Prem Devanbu, AT&T BELL LABORATORIES
- Knowledge Base to assist in software reuse.

HYPERBOLE: A RETRIEVAL-BY-REFORMULATION INTERFACE THAT PROMOTES SOFTWARE VISIBILITY
Patricia Carando, Schlumberger Laboratory for Computer Science
- Generic user interface and data analysis architecture to analyze well data.
- Graphical workstation tool (500-600 classes)

AN EMPIRICAL FRAMEWORK FOR SOFTWARE REUSE RESEARCH
Bill Frakes, Software Productivity Consortium
- Determine the relationships between the dependent variables in model
  1. quality,
  2. productivity, and
  3. reuse

THE 3C MODEL OF REUSABLE SOFTWARE COMPONENTS
Stephen Edwards, Institute for Defense Analyses
- Emphasis on the maintenance payback from using the 3C model.

THE THREE CONS OF SOFTWARE REUSE
Will Tracz, IBM Corporation
- The gospel according to Will
Paper Summaries

DESIGNING FOR SOFTWARE REUSE IN ADA
Sholom Cohen, SEI/Carnegie-Mellon University
- Implementation implications of using the 3C model in regards to hierarchies of parameterized models.
- Coupling inversion - where context is fixed for implementation efficiencies within the generic architecture.

THE PRACTITIONER REUSE SUPPORT SYSTEM (PRESS):
A TOOL SUPPORTING SOFTWARE REUSE
Cornelia Boldyreff, Brunel University
- ESPRIT 1064 Practitioner Project (one of many reuse projects funded by ESPRIT).
- Capture and reuse software concepts from designs through code.
- Questionnaire was passed out to the inam company to assist in domain analysis
- “Canonical” form for describing software components developed

REUSE AT HEWLETT-PACKARD LABORATORIES
Martin L. Griss, Hewlett-Packard Laboratories
- Hypertext tools
- Object-Oriented Design.

BEYOND RETRIEVAL: UNDERSTANDING AND ADAPTATION IN SOFTWARE REUSE
Karen Huff & Ronnie Thomson, GTE Laboratories Inc.
- SATURN (Software Adaptation Through Understandable Reuse Notation)

THE STARLITE INTELLECTUAL REUSE PROJECT
Robert P. Cook, University of Virginia
- Reusable operating system and system modeling components

Conceptual Model
Reusable Software Components

- Context
- Concepts
- Content
  - Context
  - Concepts
  - Content

Context

- “Language shapes thought”
  - Inheritance
  - Genericity/Parameterization
  - Importation
- Binding time
  - Compile time
  - Load/Bind time
  - Run Time
Conceptual Model

Concepts

• Concept: What
• Content: How
• Context:
  1. Conceptual – relationship
  2. Operational – with/to what
  3. Implementation – trade-offs

Context: what is needed to complete the definition of a concept or content within an environment. (Lévi-Strauss)

Software Components

Formal Foundations

• Horizontal Structure
  1. type inheritance
  2. code inheritance
• Vertical Structure
  – implementation dependencies
  – virtual interfaces
• Generic Structure
  – variations/adaptations

Conceptual Model

Example

• Concept: Stack
  – Operational Context: Element/Type
  – Conceptual Context: Deque
  – Implementation Context: Sequence

Conceptual Model

Example

• Stack Implementation
  1. Inherit Deque
  2. Use an array
  3. Use a linked-list
     • memory management
     • no memory management
     • concurrent access
Megaprogramming Example

Stack -> Deque

make Deque [ Triv ] is
    Stack [ Triv ]
    * ( rename ( Push = > Push_Right )
        ( Pop = > Pop_Right )
        ( Stack = > Deque )
    * ( add Push_Left, Push_Right )
end;

Hyperprogramming Example

Make with View

make Integer_Set is
    LIL_Set [ Integer_View ]
end;

view Integer_View :: Triv = > Standard is
    types ( Element = > Integer );
end;

Megaprogramming Example

Make with Vertical Composition

make Short_Stack is
    LIL_Stack
    -- horizontal composition
    needs ( List_Theory = > List_Array )
    -- vertical composition
end;

LILEANNA Example

Package Expressions

make This Ald. Logic_Interface is
    Identifier_Package <
    Domain_Package < Include Copy >
    Substituting_Package <
    Database_Package <
    Query_Package < add function Query_fail ( : clause;
    L list_of_clauses )>
    return Boolean
    return Boolean
end;

*(rename ( Query_Answer = Query_Results ));
Ada Net
John McBride
Planned Solutions

Paper not available at time of printing.
Session 4

Software Engineering: Issues for Ada's Future

Chair: Rod L. Bown, University of Houston-Clear Lake

Assessment of Formal Methods for Trustworthy Computer Systems

Susan Gerhart
Microelectronics and Computer Technology Corp. (MCC)

Paper not available at time of printing.
Issues Related to Ada 9X

John McHugh

*Computational Logic, Inc.*

Paper not available at time of printing.
This paper discusses the integration of real-time POSIX and real-time, multiprogramming Ada in the Space Station Freedom Data Management System. Use of POSIX as well as use of Ada has been mandated for Space Station Freedom flight software. However, POSIX and Ada assume execution models that are not always compatible. This becomes particularly true once Ada has been extended to support multiprogramming. This paper points out the conflicts between POSIX and Ada multiprogramming execution models and describes the approach taken in the Data Management System to resolve those conflicts.
Session 5

Ada Run-Time Issues

Chair: Alan Burns, University of York (U. K.)